

1 To do this, we introduce the following preconditions and postconditions on the fragment, C

$$\overline{\{x = a, y = b\} C \{x = b, y = a\}}$$

We now apply the assignment, consequence and composition axiom schemas several times to verify the expression:

$$\frac{\overline{\{x = a, y = b\} x = x \oplus y; y = x \oplus y; \{x = a \oplus b, y = a\}} \quad \overline{\{x = a \oplus b, y = a\} x = y \oplus x; \{x = b, y = a\}}}{\overline{\{x = a, y = b\} x = x \oplus y; y = x \oplus y; x = x \oplus y; \{x = b, y = a\}}}$$

$$\frac{\overline{\{x = a, y = b\} x = x \oplus b \{x = a \oplus b, y = b\}} \quad \overline{\{x = a \oplus b, y = b\} y = x \oplus y \{x = a \oplus b, y = a\}}}{\overline{\{x = a, y = b\} x = x \oplus y; y = x \oplus y; \{x = a \oplus b, y = a\}}}$$

(Note that \oplus is used in place of \wedge to denote exclusive-or to remove ambiguity.)

2 Using the same pre/post conditions from problem 1, we proceed with the modified assignment axiom schema:

$$\overline{\{x = a, y = b\} t = x; x = y; y = t; \{x = b, y = a\}}$$

Working from the left, we get the following:

$$\frac{\overline{\{x = a, y = b\} t = x; \{t = a, x = a, y = b\}, \{t = a, x = a, y = b\} x = y; y = t; \{x = b, y = a\}}}{\overline{\{x = a, y = b\} t = x; x = y; y = t; \{x = b, y = a\}}}$$

Taking the toprightmost expression, we further refine our derivation:

$$\frac{\overline{\{t = a, x = a, y = b\} x = y; \{t = a, x = b, y = t\}, \{t = a, x = b, y = t\} y = t; \{x = b, y = a\}}}{\overline{\{t = a, x = a, y = b\} x = y; y = t; \{x = b, y = a\}}}$$

And to complete the derivation, we observe that:

$$\frac{\overline{\{t = a, x = b, y = t\} y = t; \{t = a, x = b, y = a\}, (t = a, y = a) \implies y = a}}{\overline{\{t = a, x = b, y = t\} y = t; \{x = b, y = a\}}}$$

And so the segment correctly implements swap.

3 To prove the correctness of this segment, we introduce the following loop invariant:

$$\{x = y * q + r, 0 \leq r\} \text{ while } r \geq y \text{ do } r = r - y; q = q + 1; \text{ od } \{x = yq + r, 0 \leq r, r < y\}$$

Applying the while axiom

$$\frac{\overline{\{x = yq + r, 0 \leq r, r \geq y\} r = r - y; \{x = y, (q + 1) +, 0 \leq r\}} \quad \overline{\{x = y(q + 1) + r, 0 \leq r, r < y\} q = q + 1; \{x = yq + r, 0 \leq r\}}}{\overline{\{x = yq + r, 0 \leq r, r \geq y\} r = r - y; q = q + 1; \{x = yq + r, 0 \leq r\}}}$$

$$\frac{\overline{\{x = yq + r, 0 \leq r\} \text{ while } r \geq y \text{ do } r = r - y; q = q + 1; \text{ od } \{x = yq + r, 0 \leq r, r < y\}}}{\overline{\{x = yq + r, 0 \leq r\} \text{ while } r \geq y \text{ do } r = r - y; q = q + 1; \text{ od } \{x = yq + r, 0 \leq r, r < y\}}}$$

Which we complete using the following:

$$\frac{\overline{(x = yq + r \implies x = y(q - 1) + r - y), \{x = y(q - 1) + (r - y), 0 \leq r, r \geq y\} r = r - y \{x = y(q - 1) + r, 0 \leq r, r < y\}}}{\overline{\{x = yq + r, 0 \leq r, r \geq y\} r = r - y; \{x = y, (q + 1) +, 0 \leq r\}}}$$

To finish the proof, we must check the conditions on the intro fragment:

$$\frac{\overline{0 \leq x, x = r \implies 0 \leq r} \quad \overline{\{0 \leq x, 0 < y, r = x\} r = x; \{x = r, 0 \leq r\}} \quad \overline{\{x = r, 0 \leq r\} q = 0; \{x = yq + r, 0 \leq r\}}}{\overline{\{0 \leq x, 0 < y\} r = x; q = 0; \{x = yq + r, 0 \leq r\}}}$$

And so the code correctly implements swap.