

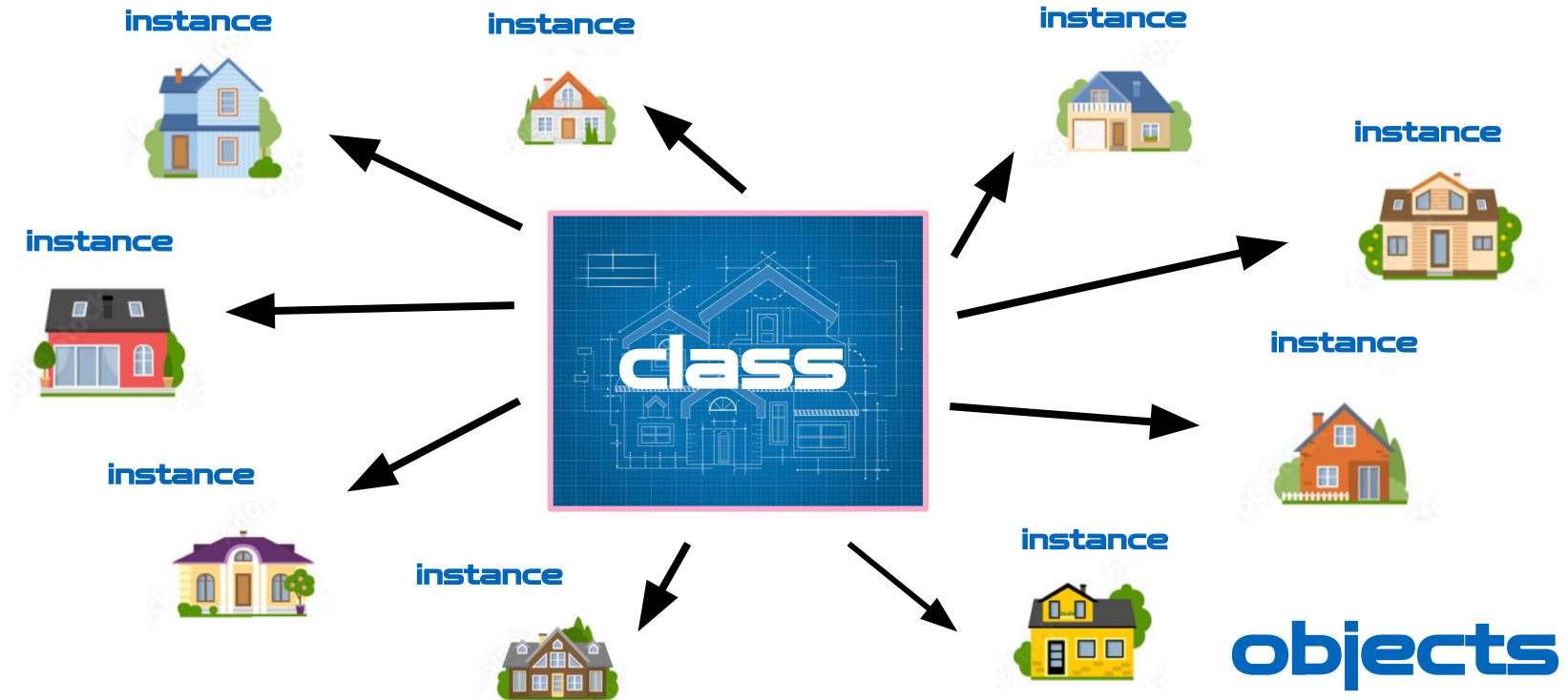
# Lecture 16: Mutable Trees

CS 61A - Summer 2024

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# Review: Objects + Inheritance

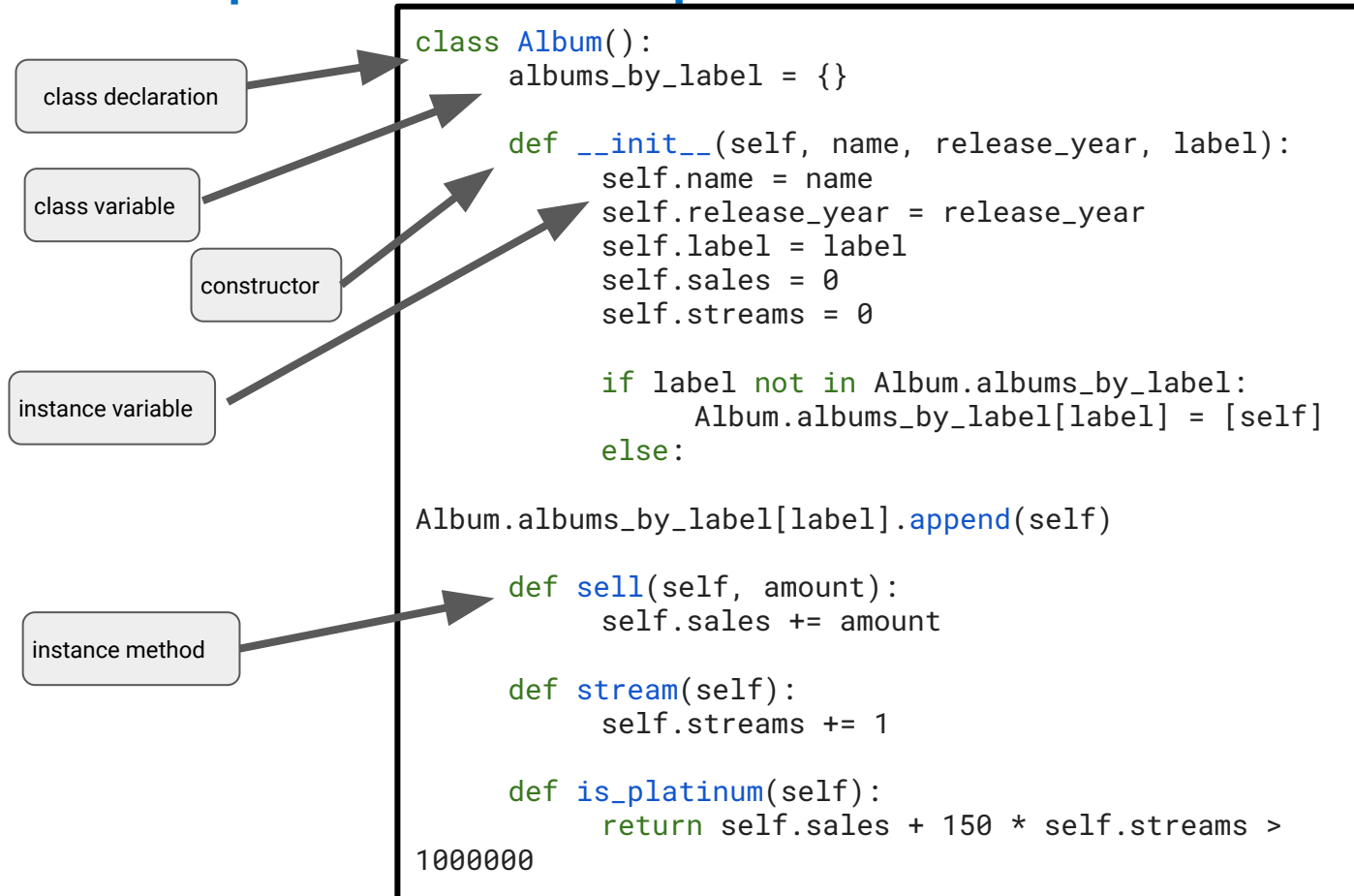
# Objects



# Definitions

- Class
  - A template for creating objects
- Instance
  - A single object created from a class
- Instance Variable
  - A data attribute of an object, specific to an instance
- Class Variable
  - A data attribute of an object, shared by all instances of a class
- Instance method
  - A function that operates on individual instances of a class
- Constructor
  - A method that specifies how to initialize an individual instance

# Example of A Complete Class Definition



# Inheritance vs Composition

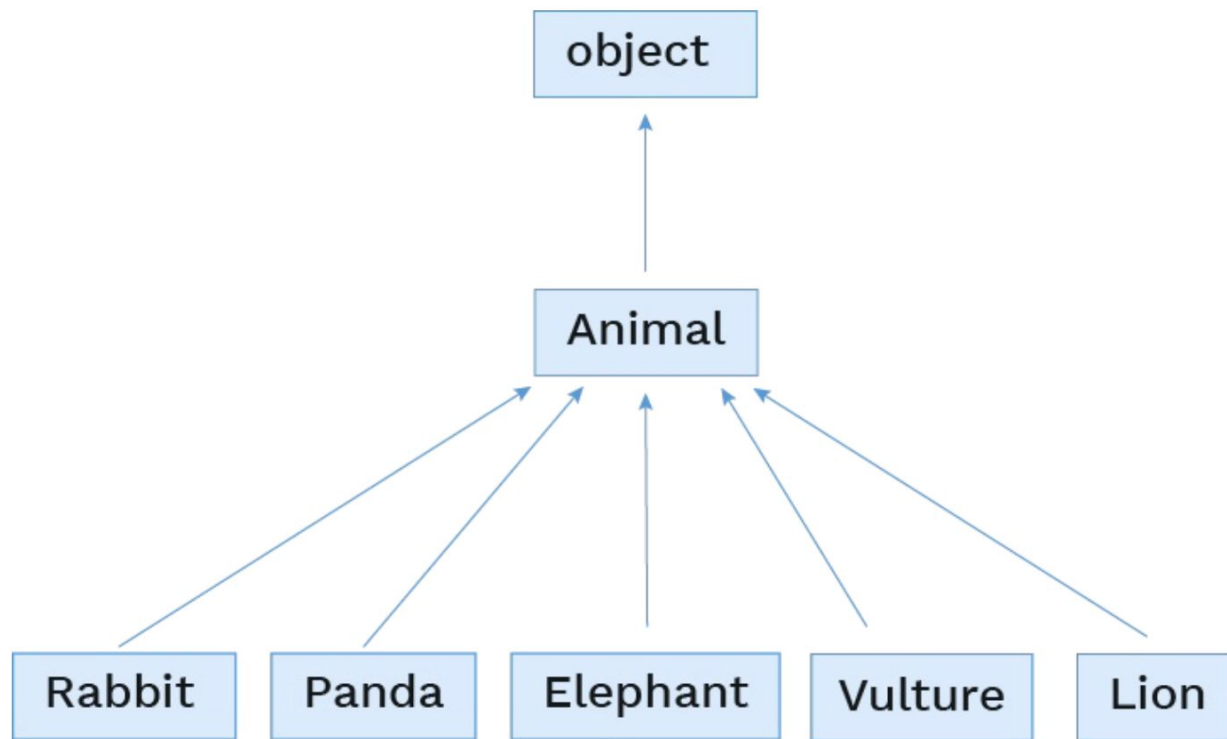
OOP is best when we think about the following metaphor:

Inheritance is best for representing **is-a** relationships

- A panda *is* an animal so Panda inherits from Animal

Composition is best for representing **has-a** relationships

- An album *has* listeners
- A zoo *has* animals



# The Album Class + The Listener Class

```
class Album():
    albums_by_label = {}

    def __init__(self, name, release_year, label):
        self.name = name
        self.release_year = release_year
        self.label = label
        self.sales = 0
        self.streams = 0

        if label not in Album.albums_by_label:
            Album.albums_by_label[label] = [self]
        else:

Album.albums_by_label[label].append(self)

    def sell(self, amount):
        self.sales += amount

    def stream(self):
        self.streams += 1

    def is_platinum(self):
        return self.sales + 150 * self.streams >
1000000
```

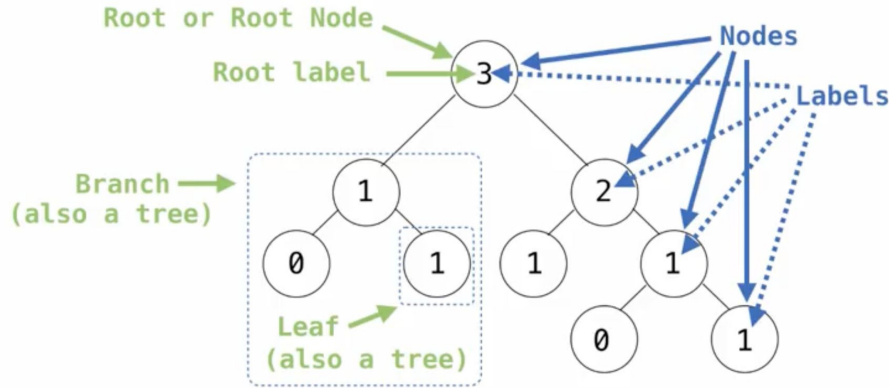
```
class Listener():
    def __init__(self, name,
favorite_album):
        self.name = name
        self.favorite_album = favorite_album
        self.albums = []
        self.buy(favorite_album)

    def buy(self, album):
        album.sell(1)
        self.albums.append(album)
```



# Mutable Trees

# Review: ADT Trees



## Recursive description (wooden trees):

A **tree** has a **root label** and a list of **branches**

Each **branch** is a **tree**

A **tree** with zero **branches** is called a **leaf**

A **tree** starts at the **root**

## Relative description (family trees):

Each location in a tree is called a **node**

Each **node** has a **label** that can be any value

One node can be the **parent/child** of another

The top node is the **root node**

# Mutable Trees

- Now that we've learned about the idea of using classes to create custom objects, let's revisit the Tree ADT and see how we can convert that to a formal Python class
- Why?
  - Using formal classes translate more directly to how we would create custom objects in the real world
    - Introducing ADTs was a delayed introduction to OOP
  - These new trees are mutable, which unlock different methods of problem solving

# Mutable Trees vs Tree ADT

```
class Tree:
    def __init__(self, label, branches=[]):
        self.label = label
        for branch in branches:
            assert isinstance(branch, Tree)
        self.branches = list(branches)
```

Tree class

```
def tree(label, branches=[]):
    for branch in branches:
        assert is_tree(branch)
    return [label] + list(branches)
def label(tree):
    return tree[0]
def branches(tree):
    return tree[1:]
```

Tree ADT

***Note that the concepts are still the same!***

Drawing out diagrams for trees, terminology associated with trees, etc

# is\_leaf

`is_leaf` is a **function** that returns whether the tree is a leaf (a node with no branches)

```
def is_leaf(self):  
    return not self.branches
```

Tree class

```
def is_leaf(tree):  
    return not branches(tree)
```

Tree ADT

Note that `is_leaf` is **not an attribute** of the Tree class!

## count\_leaves: Revisited

```
def count_leaves(t):  
    if is_leaf(t):  
        return 1  
    else:  
        branch_counts = 0  
        for b in branches(t):  
            branch_counts += count_leaves(b)  
        return branch_counts
```

Let's try to implement this using the Tree class!

## count\_leaves: Revisited

```
def count_leaves_2(t):  
    if t.is_leaf():  
        return 1  
    else:  
        branch_counts = 0  
        for b in t.branches:  
            branch_counts += count_leaves_2(b)  
        return branch_counts
```

As shown: We can apply a majority of the concepts we've already learned to solving problems using the Tree class.

# String Representation of Trees



# String Representations

- Strings are important - they represent language and programs
- Printing strings is exactly how the interpreter communicates back to us
  - We should be able to use strings to represent objects!
- Python gives us two ways/functions to represent objects
  - **str**: returns a representation that a **human** should be able to read
    - used by the print function
  - **repr**: returns a representation that the **interpreter** should be able to read
    - used by the interpreter for printing results
  - these are usually the same, but not always!

# String representation of trees

- When creating tree objects using code, it often is hard to visualize the properties of each node, which make up the overall tree structure
- The String representation of trees makes visualizing trees easier!

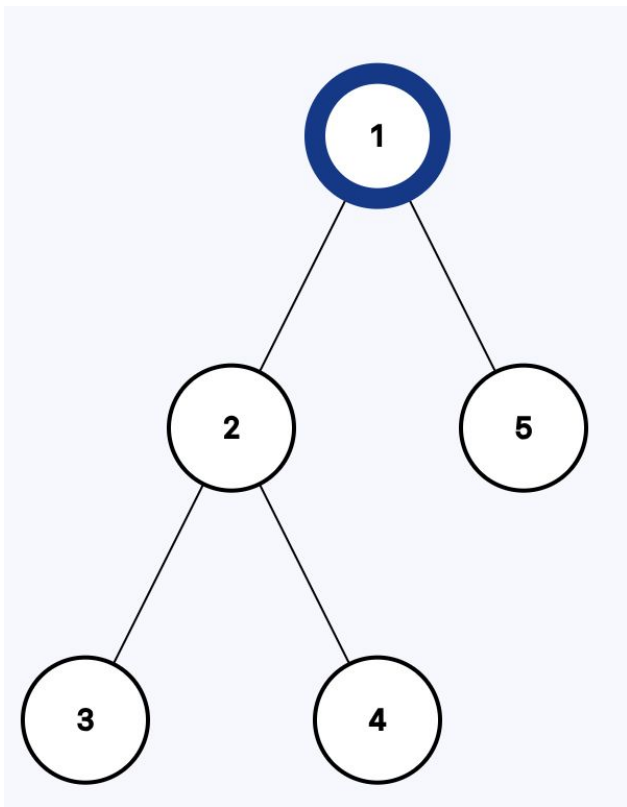
```
def __repr__(self):  
    if self.branches:  
        branch_str = ', ' + repr(self.branches)  
    else:  
        branch_str = ''  
    return 'Tree({0}{1})'.format(self.label, branch_str)
```

`repr` method  
“Computer readable”

```
def __str__(self):  
    def print_tree(t, indent=0):  
        tree_str = ' ' * indent + str(t.label) + "\n"  
        for b in t.branches:  
            tree_str += print_tree(b, indent + 1)  
        return tree_str  
    return print_tree(self).rstrip()
```

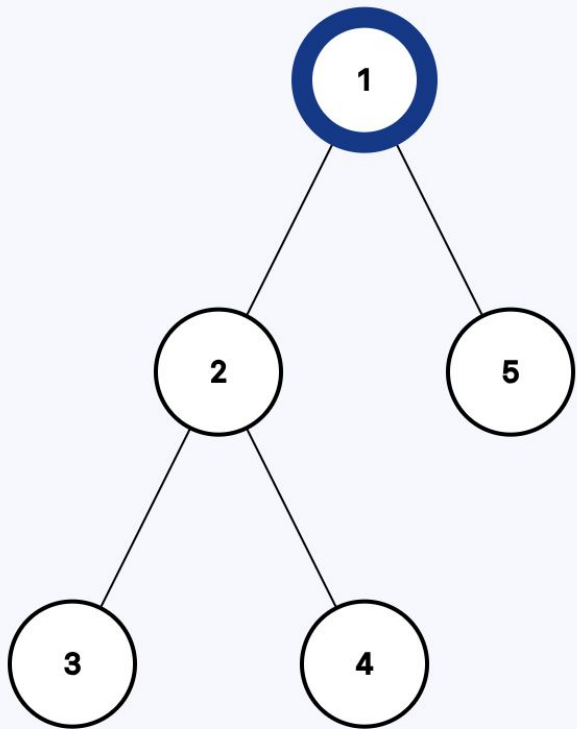
`str` method  
“Human readable”

# Examples



```
>>> t = Tree(1, [Tree(2, [Tree(3), Tree(4)]), Tree(5)])
>>> t
Tree(1, [Tree(2, [Tree(3), Tree(4)]), Tree(5)])
>>> print(t)
1
  2
    3
    4
  5
```

# Examples



```
>>> t = Tree(1, [Tree(2, [Tree(3), Tree(4)]), Tree(5)])
>>> repr(t)
'Tree(1, [Tree(2, [Tree(3), Tree(4)]), Tree(5)])'
>>> str(t)
'1\n_2\n    3\n    4\n  5'
```

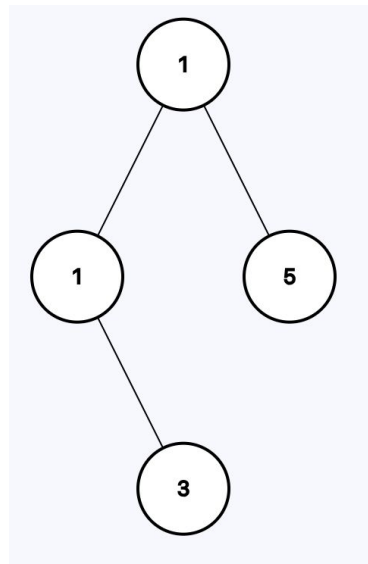
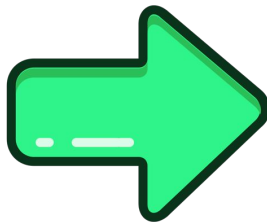
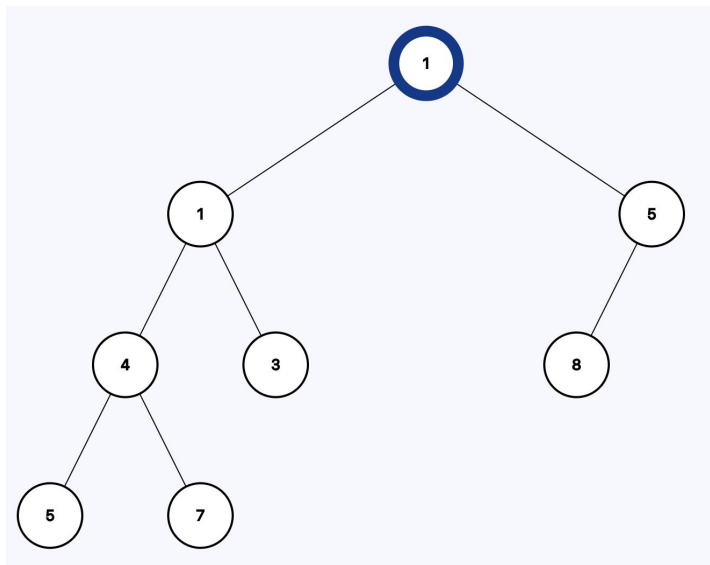
Invoking an object directly in the interpreter calls `repr` on the object, and prints out the output

Printing a string calls `str` on the object, and prints out the output

# Examples

## Example: prune\_evens

- Problem: Given an input tree `t`, **mutate** the tree such that all labels with an even value are removed from the tree.
  - If a node containing an even value is removed, all branches belonging to that node are removed as well



# prune\_evens: Solution

`is_even` is a function we defined in lecture that takes in a value and returns whether it is even

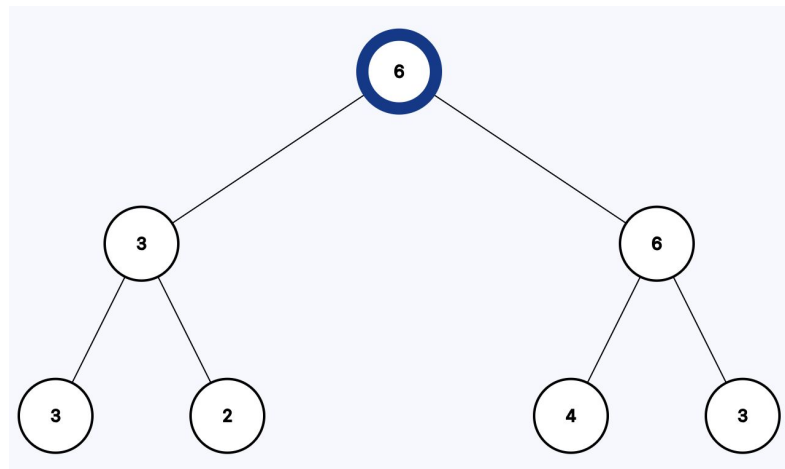
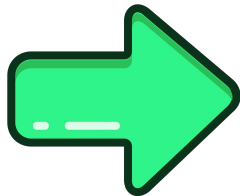
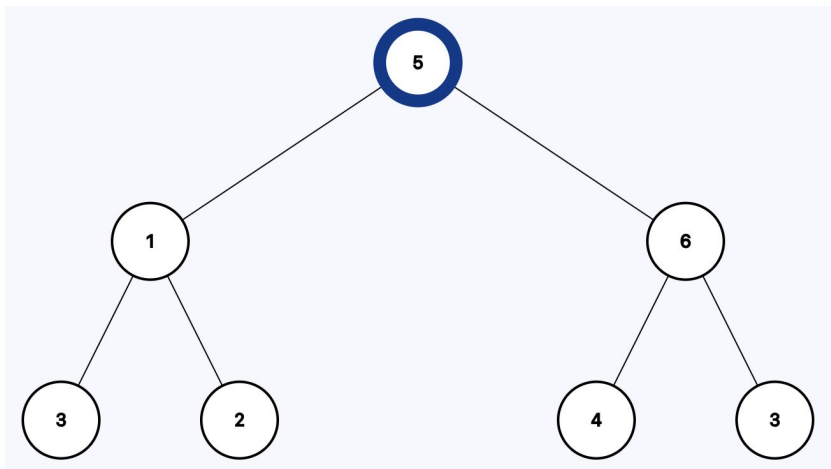
```
def prune_even(t):  
    """Mutate the tree t such that all nodes containing even values are removed."""  
    if is_even(t.label):  
        return  
    new_branches = []  
    for b in t.branches:  
        if not is_even(b.label):  
            prune_even(b)  
            new_branches.append(b)  
    t.branches = new_branches # This is where the mutative step happens!
```

To mutate an object, we need to modify one of its attributes (in this case, modifying the `branches` attribute)

The return in the base case is used to stop the recursion, not to explicitly return a value used by a previous frame

## Example: largest\_of\_group

- Problem: Given an input tree  $t$ , **mutate** each node such that the label of a node is the largest among all of its children plus its original value label.





## largest\_of\_group: Solution #1

```
def largest_of_group(t):  
    largest_labels = [t.label]  
    for b in t.branches:  
        largest_of_group(b)  
        largest_labels.append(b.label)  
    t.label = max(largest_labels)
```

## largest\_of\_group: Solution #2

```
def largest_of_group(t):  
    for b in t.branches:  
        largest_of_group(b)  
    t.label = max([t] + t.branches, key=lambda n : n.label).label
```

Here, we were taking the max over a list of tree nodes. Python doesn't know how to compare tree nodes directly, so we pass in a **key** argument which specifies what attribute of tree nodes we use during comparison. Note that max still evaluates to a tree node, and not the attribute that the **key** function was using!

Break

## Example: keep\_k\_largest

- Problem: Given a tree  $t$  and a value  $k$ , **mutate** the tree such that each node has no more than  $k$  branches. If a node has more than  $k$  branches, keep the  $k$  largest ones.

## keep\_k\_largest: Solution

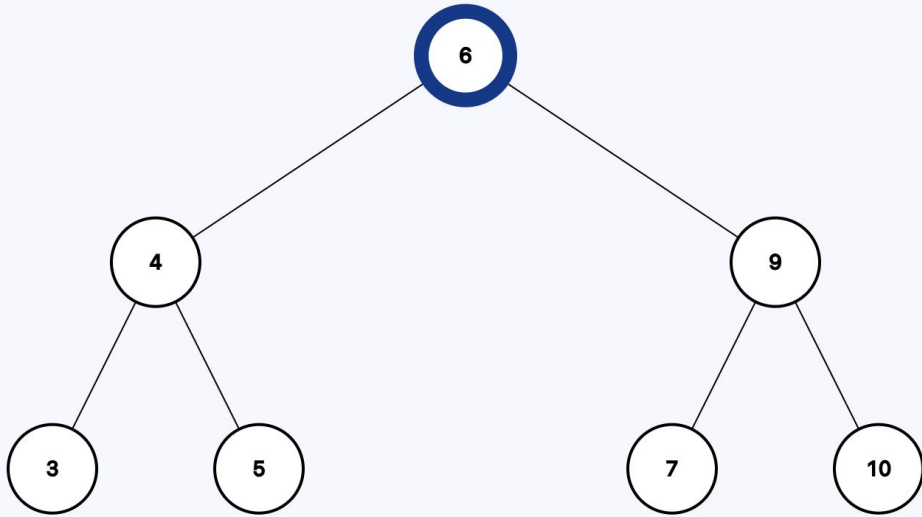
```
def keep_k_largest(t, k):  
    n = len(t.branches)  
    counter = k  
    while counter > n:  
        smallest = min(t.branches, key=lambda n : n.label)  
        t.branches.remove(smallest)  
        counter -= 1  
    for b in t.branches:  
        keep_k_largest(b, k)
```

# Extra: Binary Search Trees

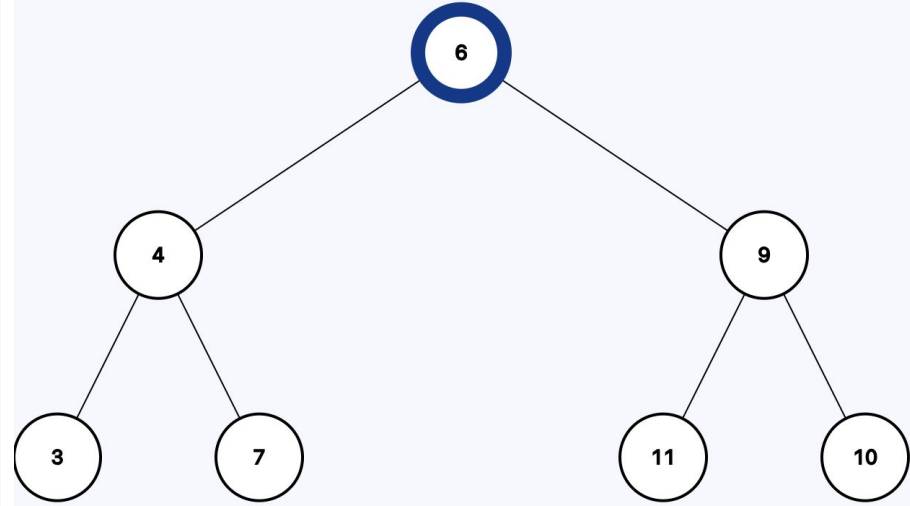
# Binary Search Trees

- Binary Search Trees (BSTs) are a specific type of tree with the following properties:
  - Every node has 0, 1, or 2 branches (binary)
  - Every value to the **left** of a node is **less** than the label of the current node, and every value to the **right** of a node is **greater** than the label of the current node
    - Definitions vary, but oftentimes BSTs do not allow duplicate values. We'll assume this implementation
- We'll also assume that with our Tree class, if a node has exactly one branch, it is the left branch
- These properties of Binary Search Trees make them very efficient to search for values

# BSTs vs Non-BSTs



BST



Not a BST





## Example: search\_tree

- Problem: Given a BST `t` and a value `x`, return whether or not `x` exists as a label in `t`.

## search\_tree: Solution #1

```
def search_tree(t, x):  
    if t.label == x:  
        return True  
    for b in t.branches:  
        found = search_tree(b, x)  
        if found:  
            return True  
    return False
```

With this approach, are we taking advantage of the properties of BSTs?

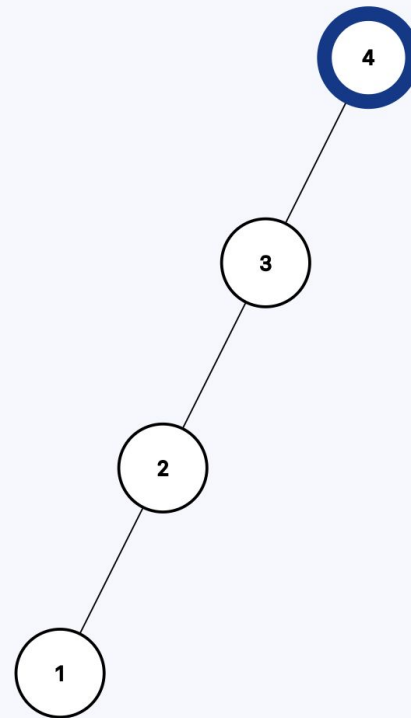
## search\_tree: Solution #2

```
def search_tree(t, x):  
    if t.label == x:  
        return True  
    n = len(t.branches)  
    if n == 0: # Leaf node  
        return False  
    elif n == 1: # Only has left branch  
        if x < t.label:  
            return search_tree(t.branches[0])  
        else:  
            return False  
    else: # Has both left and right branch  
        if x < t.label:  
            return search_tree(t.branches[0])  
        else: # x > t.label  
            return search_tree(t.branches[1])
```

Assuming we have a “balanced” BST, searching for a value is logarithmic/ $O(\log n)$  runtime with this approach.

# Is searching BSTs always logarithmic?

- No. Consider the case of a “stringy” BST
- This is a valid BST!
- However, it doesn’t allow us to get the behavior of halving all values at each step
- In fact, this structure is closer to a Linked List (our topic for tomorrow’s lecture!)



# Summary

- The Tree class is the OOP version of the Tree ADT
  - ADTs are a delayed introduction to OOP
- Same concepts from Tree ADT can be applied to Tree OOP!
- Trees created from the Tree class are mutable
  - No longer have to create new trees to solve problems
- Extra: Binary Search Trees are an efficient data structure for lookup of data