# **Homework 5**

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#### **Problem 1**

Prove:  $X \to Y$ ,  $Z \to W \Rightarrow XZ \to YW$ 

$$X \to Y \Rightarrow X \: Z \to Y \: Z$$

$$Z \to W \Rightarrow YZ \to YW$$

 $XZ \rightarrow YZ, YZ \rightarrow YW \Rightarrow XZ \rightarrow YW \blacksquare$ 

#### **Problem 2**

Prove:  $W \rightarrow Y$ ,  $X \rightarrow Z \Rightarrow W X \rightarrow Y$ 

$$W \to Y \Rightarrow W \: X \to Y \: X$$

$$YX\supset Y\Rightarrow YX\rightarrow Y$$

 $WX \to YX, YX \to Y \Rightarrow WX \to Y \blacksquare$ 

#### **Problem 3**

Prove:  $X \to Y$ ,  $X \to W$ ,  $W Y \to Z \Rightarrow X \to Z$ 

 $X \to Y W, W Y \to Z \Rightarrow X \to Z \blacksquare$ 

### **Problem 4**

$$F = \{A \rightarrow C, \, A\, C \rightarrow D, \, E \rightarrow A\, D, \, E \rightarrow H\}$$

$$G = \{A \rightarrow CD, E \rightarrow AH\}$$

Let's try to decompose G to get the FDs in F.

For G:

 $A \to CD \Rightarrow A \to C$ ,  $A \to D$  (assuming we have proved decomposition from Armstrong's axioms)

$$A \to D \Rightarrow A C \to D C, D C \supset D \Rightarrow D C \to D \Rightarrow A C \to D$$

$$E \to A H \Rightarrow E \to A, E \to H$$

$$E \to A$$
,  $A \to CD \Rightarrow E \to CD$ 

$$E \rightarrow CD \Rightarrow E \rightarrow C, E \rightarrow D$$

$$E \to A, E \to D \Rightarrow E \to AD$$

This gives us all four of the FDs in F.

# **Problem 5**

$$R = \{A, B, C, D, E, F, G, H, I, J\}$$

$$F = \{AB \rightarrow C, BD \rightarrow EF, AD \rightarrow GH, A \rightarrow I, H \rightarrow J\}$$

The key must be ABD, as these are the three attributes not determined by a functional dependency. Furthermore, these three attributes are necessar to determine every attribute in R.

Now we decompose. First we find a canonical cover.

 $F = \{AB \rightarrow C, BD \rightarrow EF, AD \rightarrow GH, A \rightarrow I, H \rightarrow J\}$  appears to already be a canonical cover, since all left-hand sides are unique and there are no extraneous attributes.

Now we create a relation for each FD in the cover.

$$R_1 = \{A, B, C\}, R_2 = \{B, D, E, F\}, R_3 = \{A, D, G, H\}, R_4 = \{A, I\}, R_5 = \{H, J\}$$

There are no schemas that are subsets of one another, so we keep them all.

However, there is no schema with the key of R.

We add  $R_6 = \{A, B, D\}$ , and we're done.