

# Minimum characters to be added at front to make string palindrome

we need to check a string it is a palindrome or not, if it is a palindrome then return 0, else we need to add some characters to the front of string such that characters added would be minimum as possible.

For example,

consider string,

$S = "abc"$

$i=0$   
 $j=2$   
 $last\_idx=2$

→ check  
↓  $s[0:2]$   
Not passed  
→  $i=0$   
 $last\_idx--$   
 $j=last\_idx$

Step 2:

$i=0$   $j=1$   $last\_idx=1$

→ Check &  $S = "abc"$   
↓  
Not palin  $s[0:j]$   
→  $i=0$   
 $last\_idx--$   
 $j=last\_idx$

Step 3:

$i=0$   $j=0$   $last\_idx=0$   
→ check &  $S = "abc"$   
↓  
palin  $s[0]=s[0]$   
→  $i=1$   $j=0$

Now the Min length would be

$n-1 - last\_index$

→  $3-1-0 \Rightarrow 2$

Algorithm:

we need to iterate through the string until  $i \leq j$ , we need to stop when it becomes  $i > j$

Now we need to check, if  $str[i] == str[j]$ , initially  $i=0$  &  $j=n-1$ , if it is true, then do  $i++$  &  $j--$

If  $str[i] \neq str[j]$  then, we can assume, we need to add  $str[j]$  to start of the string. so, neglect that index for now. Then do,

$i=0$

$last\_index--$

$j=last\_index$

Initially, last-index is pointing to  $n-1$

we do, return  $n-1 - last\_index$  because, last-index counts num of mismatches which are avoiding the current string to become a palindrome

while ( $i \leq j$ ) {

if ( $str[i] == str[j]$ ) {

$i++$

$j--$

} else {

$i=0$  &  $last\_index--$ ;

$j=last\_index$

} else case is used to count these mismatches

return ( $n-1 - last\_index$ )

Which returns, minimum characters to be added at front to make string palindrome.