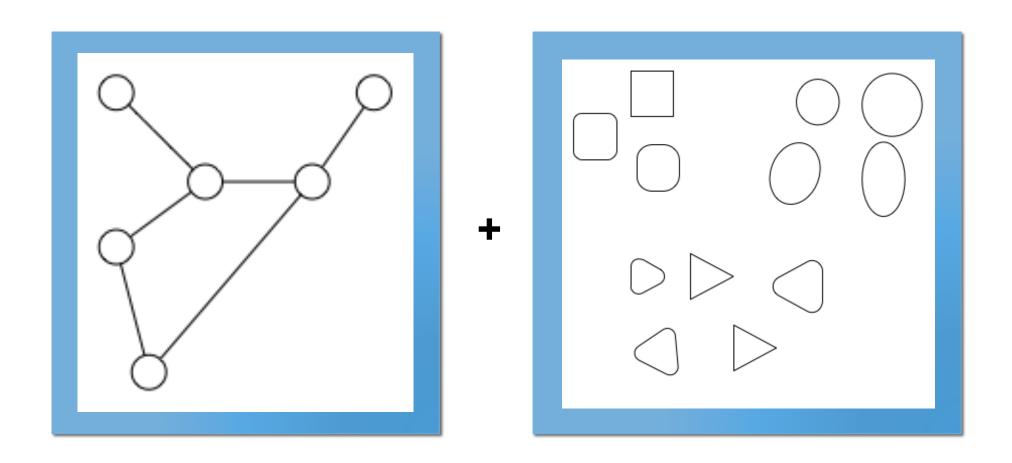
GRAPHS + CLUSTERING

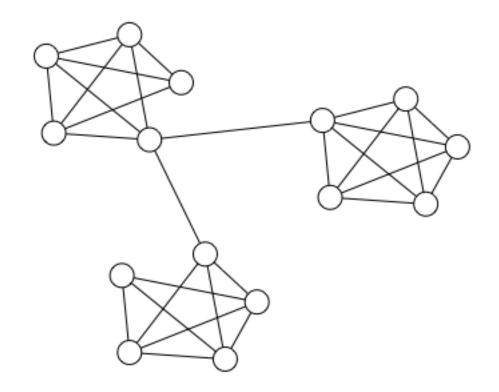
Together At Last!

April 2016

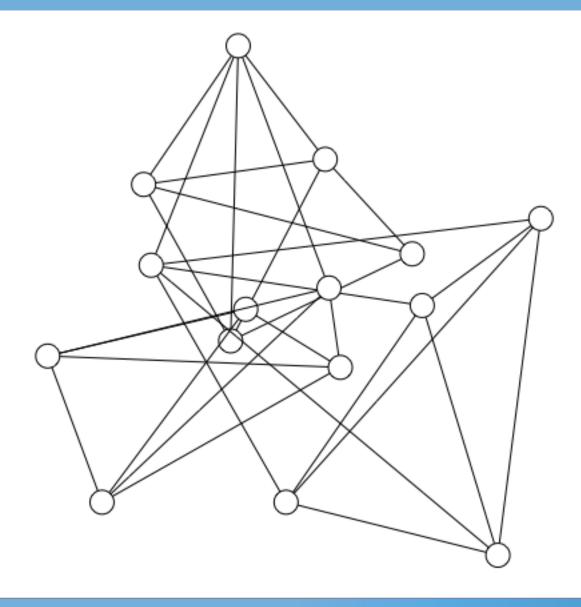
Jen Schellinck



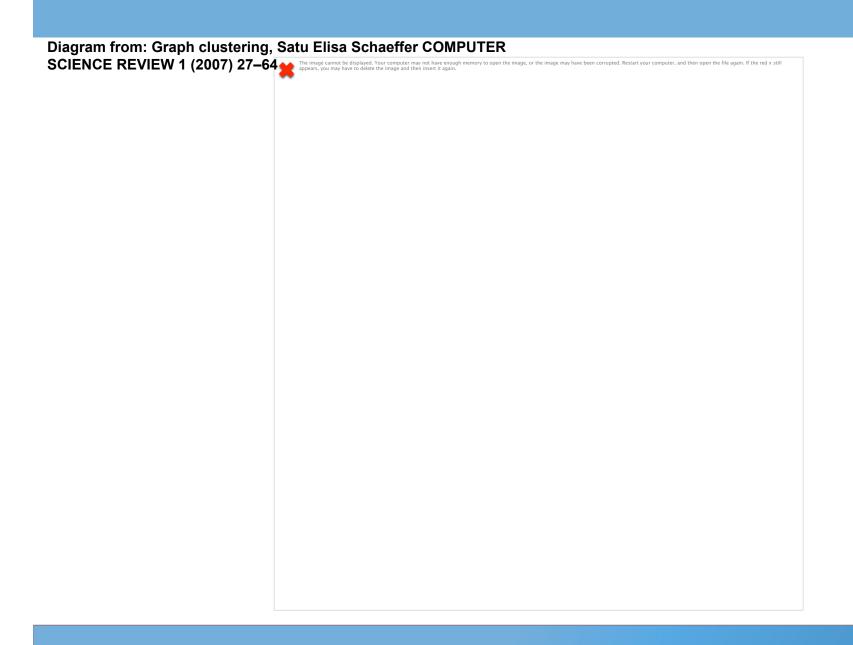
graphs + clustering = a lot to talk about!



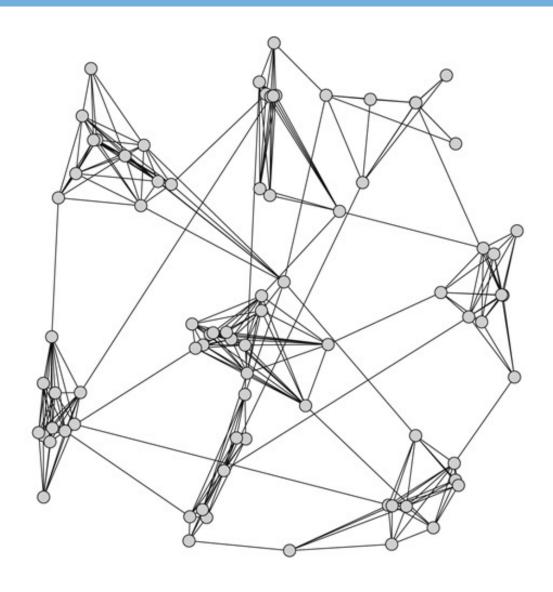
Where are the clusters? Obvious?



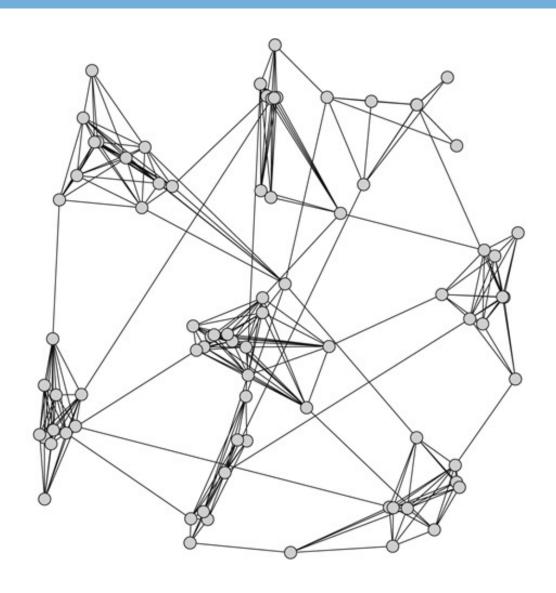
Where are the clusters? Not obvious?



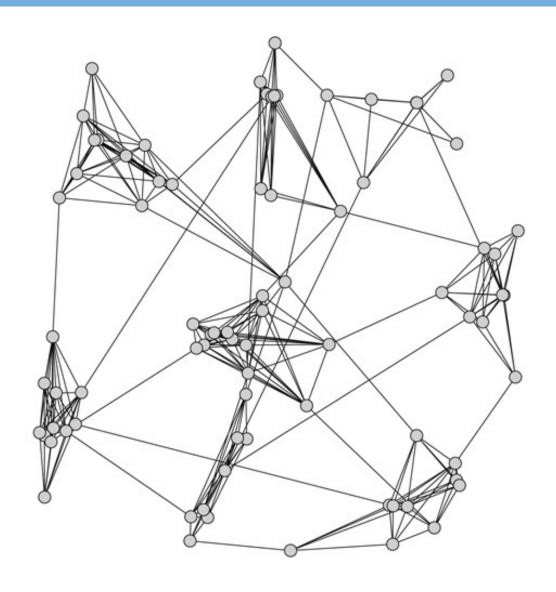
high density separated by low density



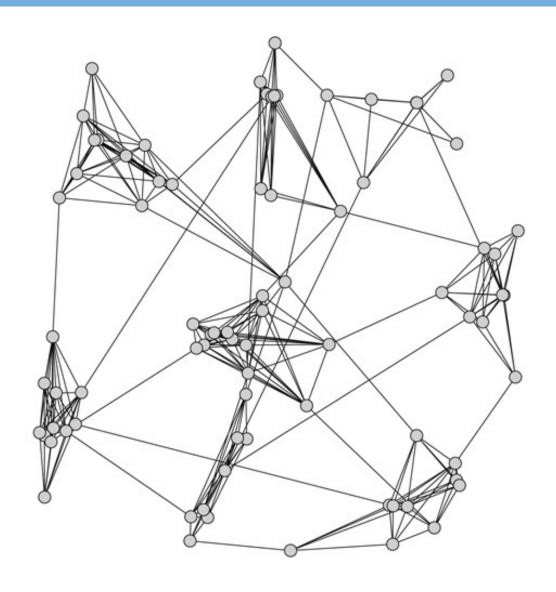
What if this graph represented: a road network?



What if this graph represented: flight paths?



What if this graph represented: neuron connections?



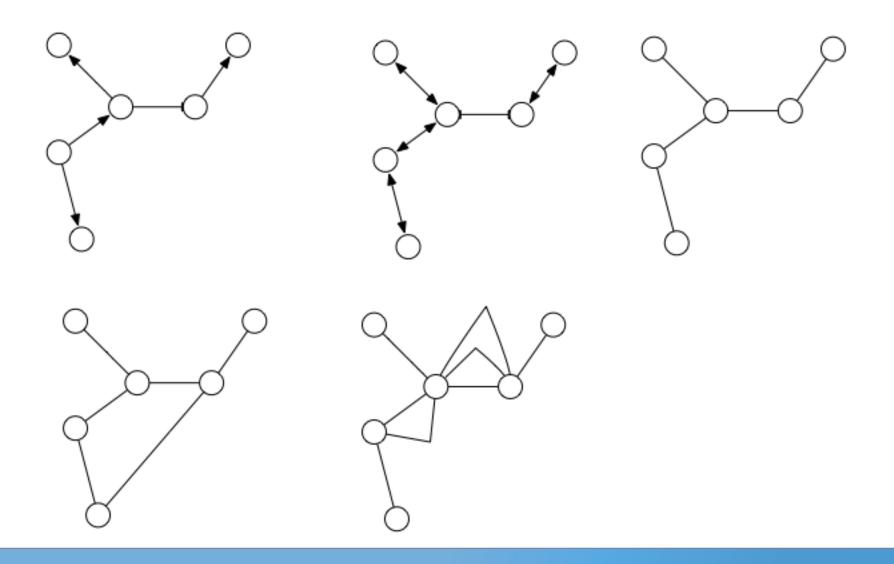
What if this graph represented: people's friends?



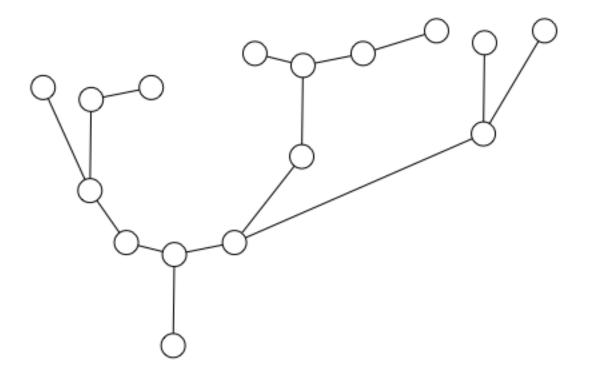
PRESS FOR CLUSTERING

Clustering: The 'Dark Art'

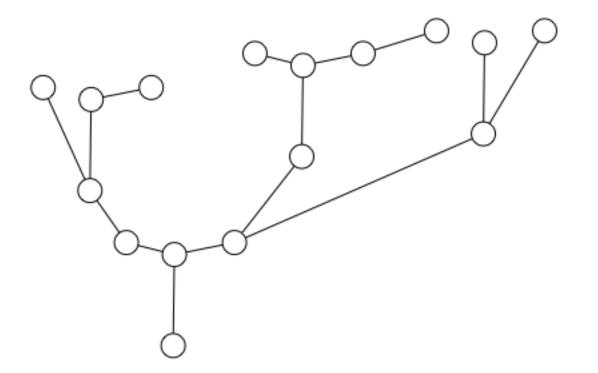
- No agreed upon definition of a cluster
- Different algorithms use different definitions
- Algorithms may not be deterministic (due to NP hard problem)



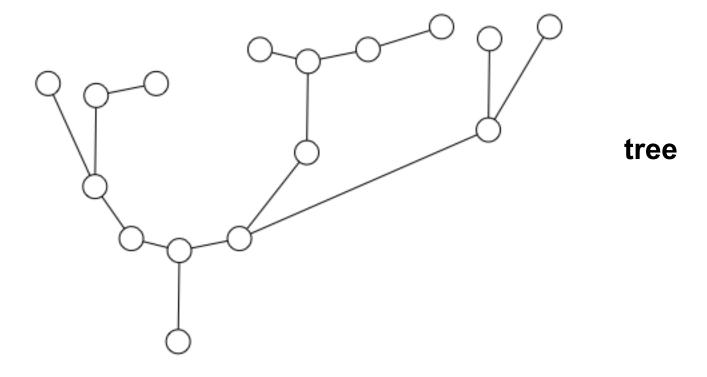
Some basics (not too poetic yet)

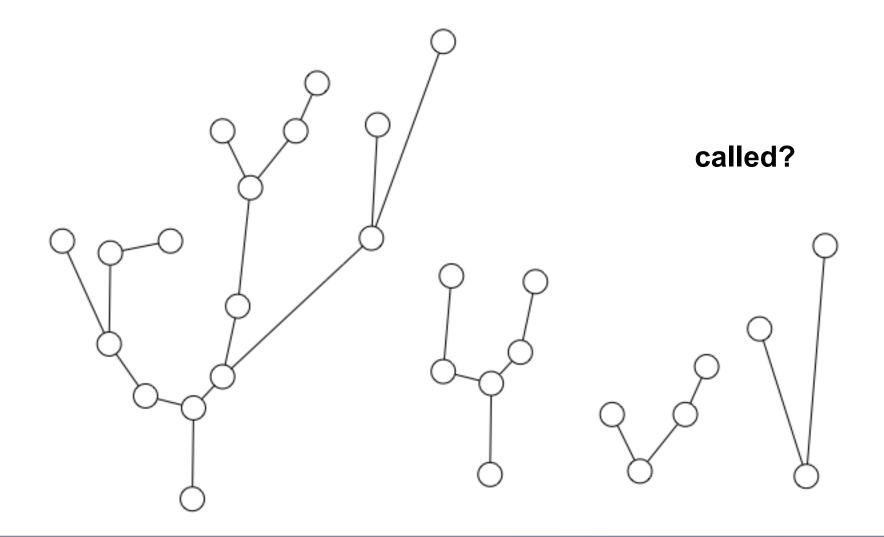


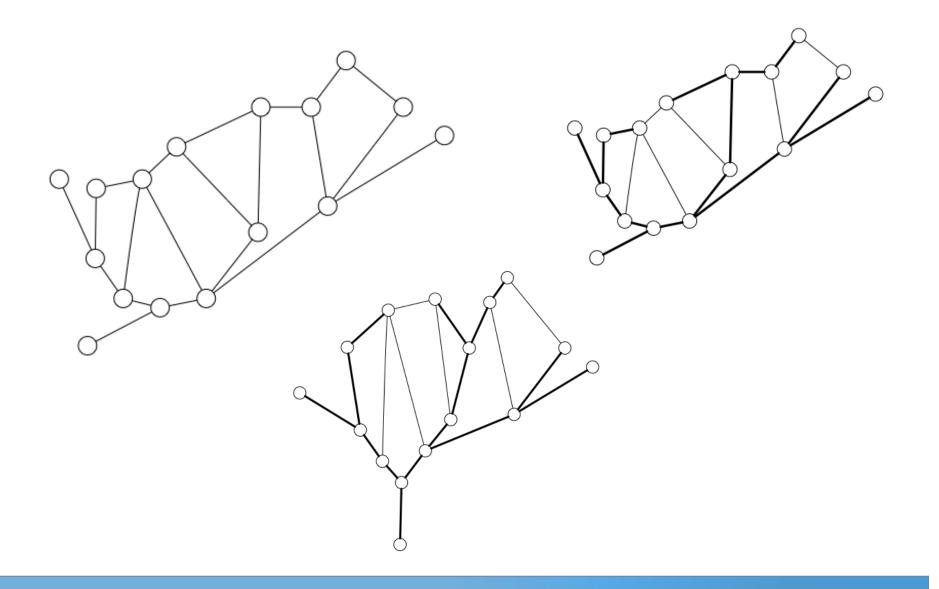
acyclic graph



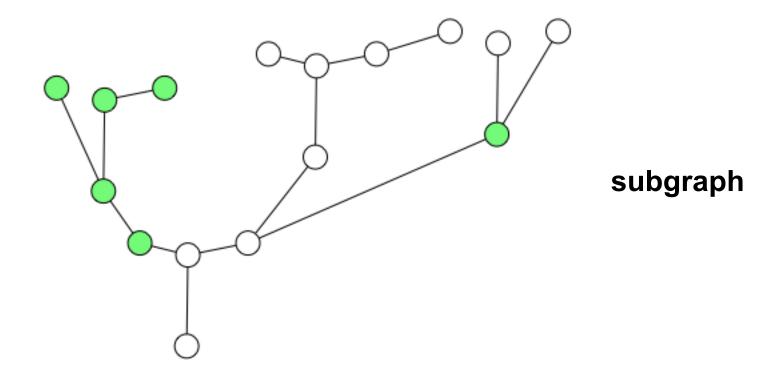
also called?

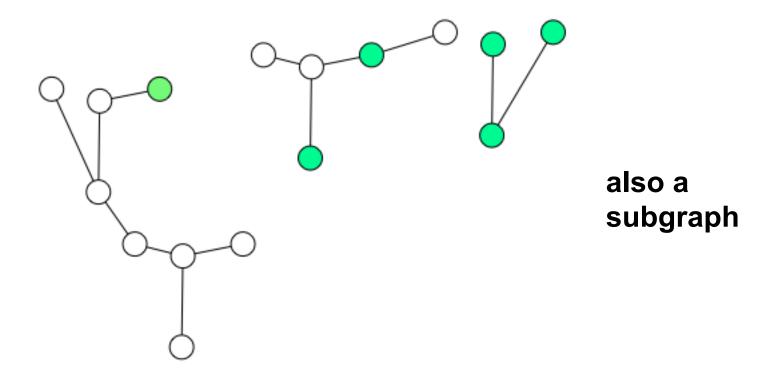


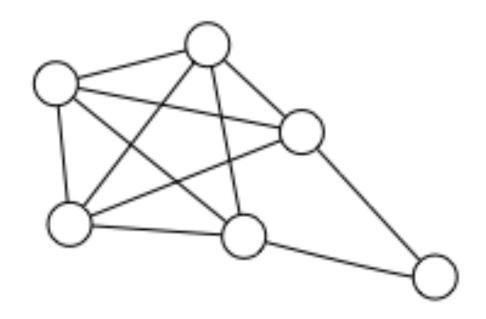




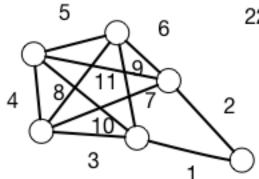
a spanning tree





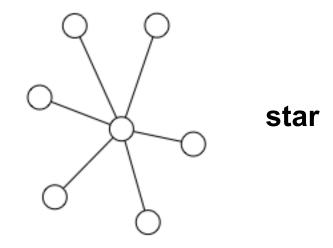


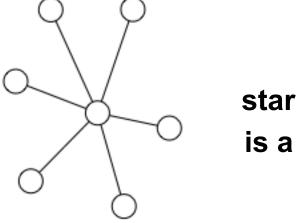
density = number of connections/total number of possible connections



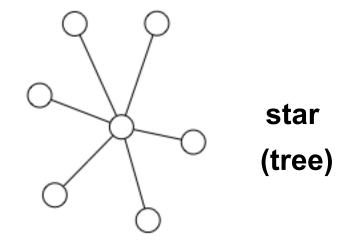
22 connections (both directions)/30 possible connections = 0.73 density

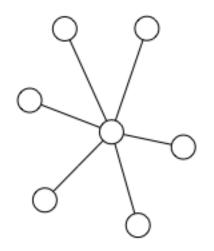
density = number of connections/total number of possible connections



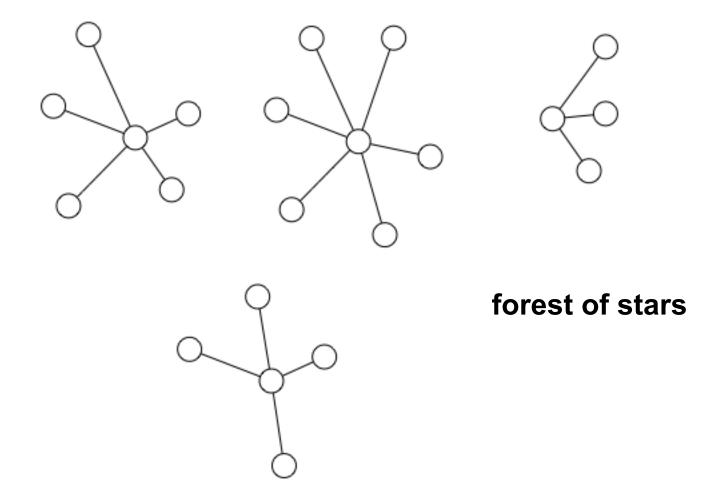


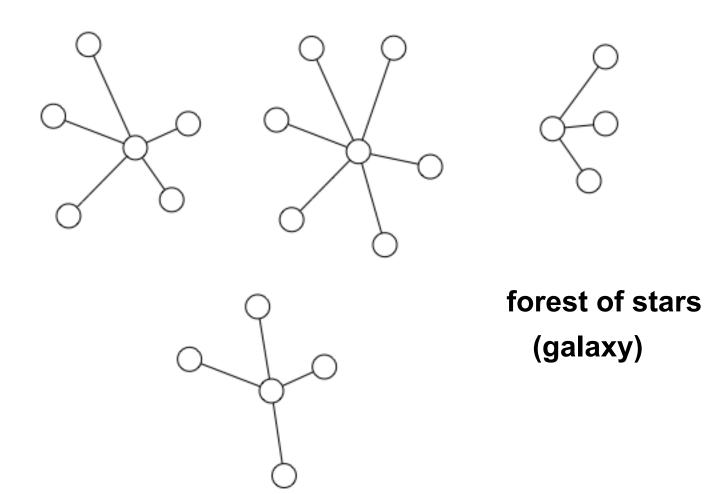
star is a type of...?



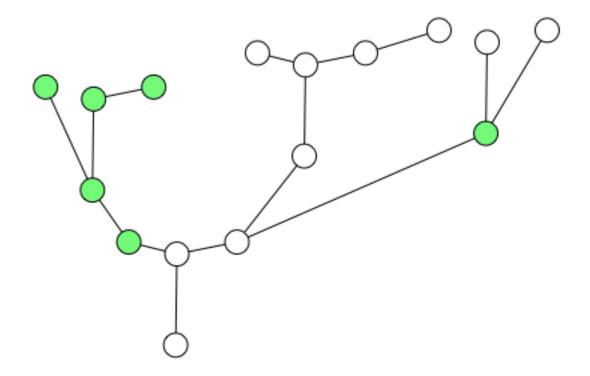


edge-graceful star





"Star arboricity is the minimum number of forests that a graph can be partitioned into such that each tree in each forest is a star" "Star arboricity is the minimum number of forests that a graph can be **partitioned** into such that each tree in each forest is a star"



partitioned: divided into sub groups

$$D = \frac{2|E|}{|V|\left(|V| - 1\right)}$$

Είναι όλα ελληνικά για μένα

$$D = \frac{2|E|}{|V|(|V| - 1)}$$

(Translation: graph density equals the number of graph edges, considered in both directions, divided by the total possible number of edges that the graph could in theory have, which is, combinatorally, the number of vertices times the number of vertices minus 1)

Είναι όλα ελληνικά για μένα

(Translation: It's all greek to me)

a note about mathematical notation

A graph G is a pair of sets G=(V,E). V is the set of vertices and the number of vertices n=|V| is the order of the graph. The set E contains the edges of the graph. In an undirected graph, each edge is an unordered pair $\{v,w\}$. In a directed graph (also called a digraph in much literature), edges are ordered pairs. The vertices v and w are called the endpoints of the edge. The edge count |E|=m is the size of the graph. In a weighted graph, a weight function $\omega:E\to\mathbb{R}$ is defined that assigns a weight on each edge. A graph is planar if it can be drawn in a plane without any of the edges crossing.

$$\delta(G) = \frac{m}{\binom{n}{2}}.$$
 $a_{v,u}^G = \begin{cases} 1, & \text{if } \{v,u\} \in E, \\ 0, & \text{otherwise.} \end{cases}$

The number of edges incident on a given vertex v is the degree of v and is denoted by deg(v). A graph is regular if all of the vertices have the same degree; if $\forall v \in V$ in G = (V, E) we have deg(v) = k, the graph G is k-regular. The diagonal degree matrix of a graph G = (V, E) is

$$D = \begin{pmatrix} \deg(v_1) & 0 & 0 & \dots & 0 & 0 \\ 0 & \deg(v_2) & 0 & \dots & 0 & 0 \\ 0 & 0 & \deg(v_3) & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & \deg(v_{n-1}) & 0 \\ 0 & 0 & 0 & \dots & 0 & \deg(v_n) \end{pmatrix}. (5)$$

A path from v to u in a graph G=(V,E) is a sequence of edges in E starting at vertex $v_0=v$ and ending at vertex $v_{k+1}=u$;

$$\{v, v_1\}, \{v_1, v_2\}, \dots, \{v_{k-1}, v_k\}, \{v_k, u\}.$$
 (8)

The cut size is the number of edges that connect vertices in S to vertices in V\S:

$$c(S, V \setminus S) = |\{\{v, u\} \in E \mid u \in S, v \in V \setminus S\}|.$$
(6)

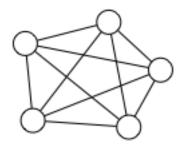
We denote by

$$\deg(S) = \sum_{v \in S} \deg(v) \tag{7}$$

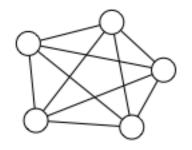
An induced subgraph of a graph G = (V, E) is the graph with the vertex set $S \subseteq V$ with an edge set E(S) that includes all such edges $\{v, u\}$ in E with both of the vertices v and u included in the set S:

$$E(S) = \{ \{v, u\} \mid v \in S, u \in S, \{v, u\} \in E \}.$$
 (9)

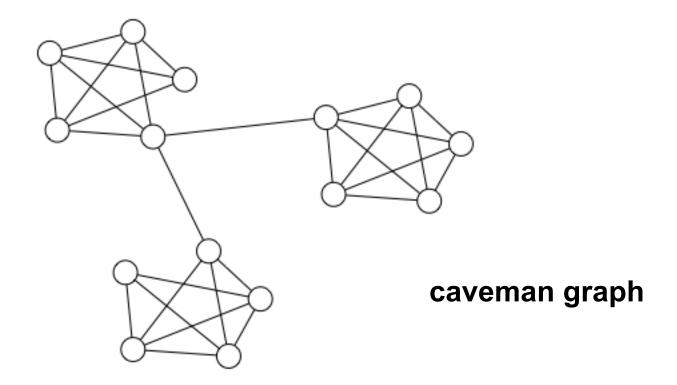
Graph clustering, Satu Elisa Schaeffer COMPUTER SCIENCE REVIEW 1 (2007) 27–64

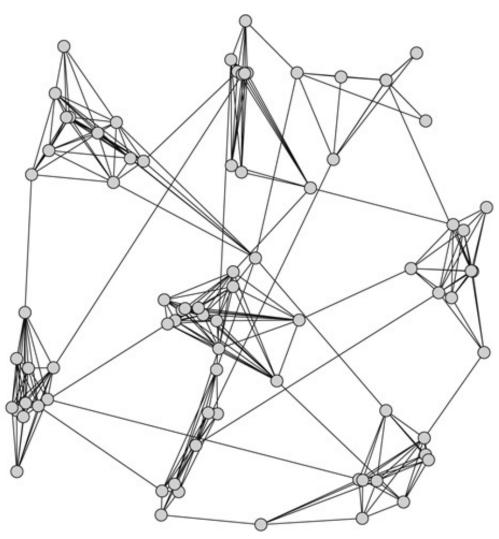


clique

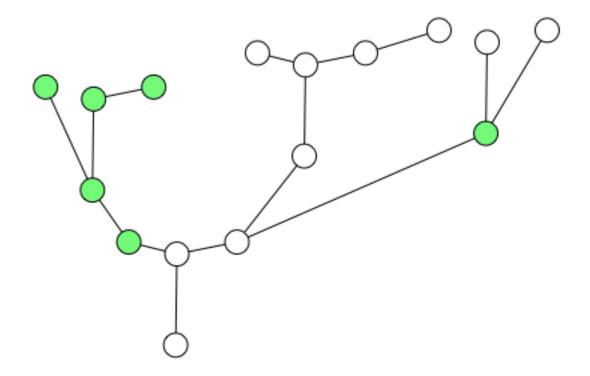


clique (cave)





relaxed caveman graph

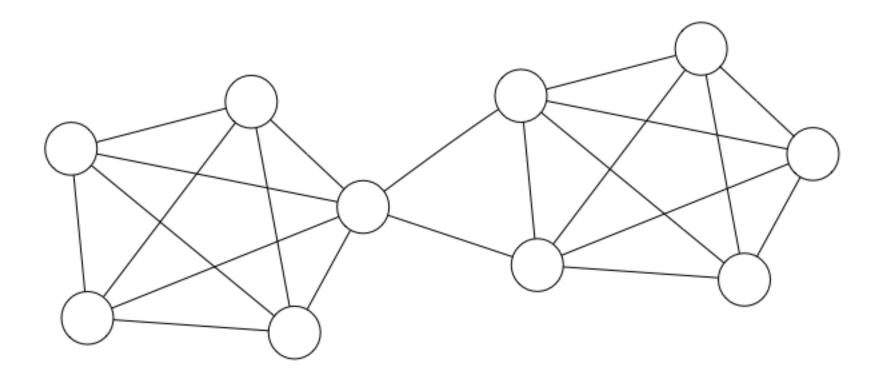


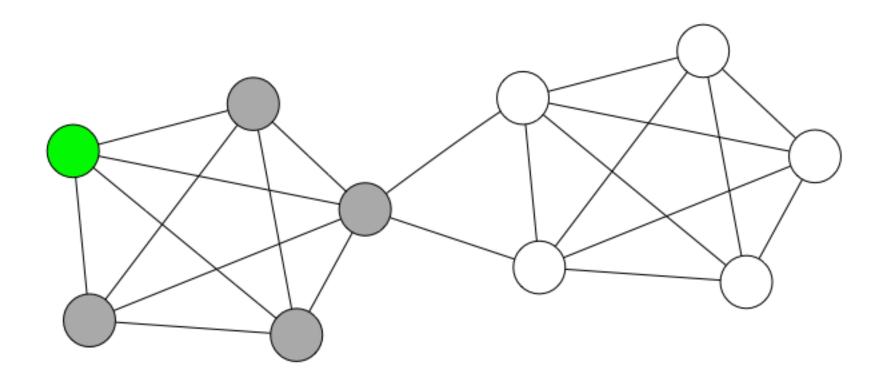
randomly 'clustered' (but not really clustering, per se)

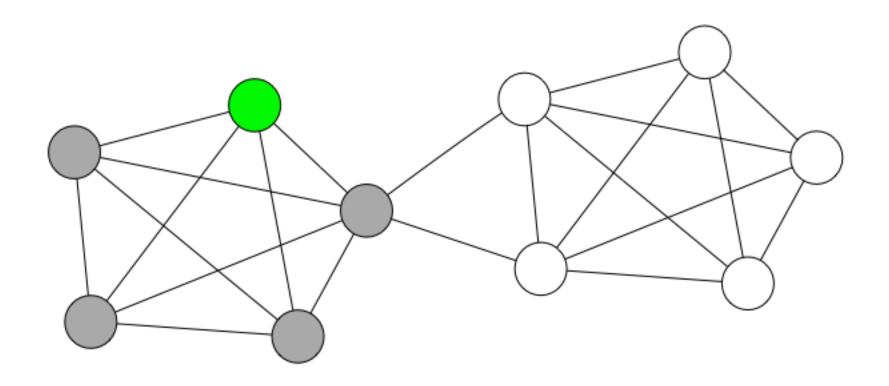
Defining clusters

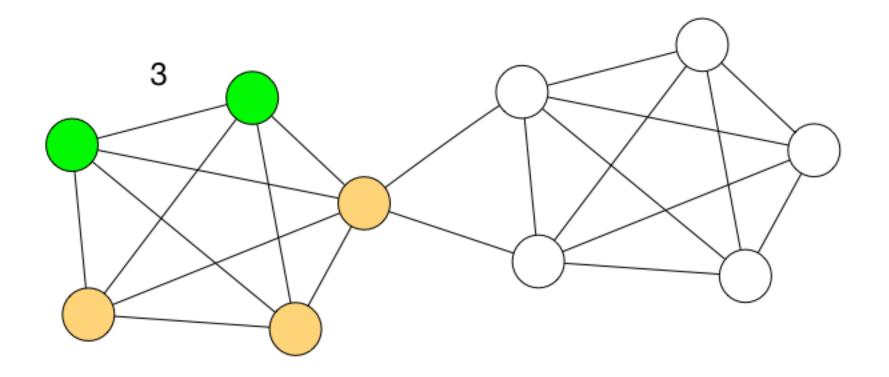
Partition based on:

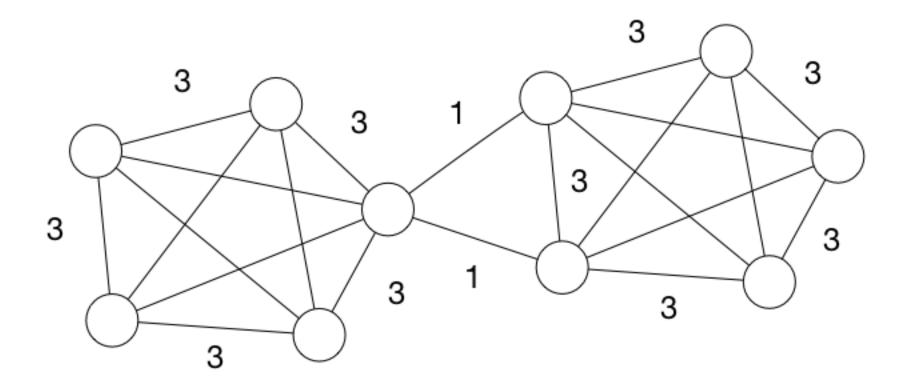
- High-Low density
- Shared neighbours
- Traversal probability (random walk)
- Cuts required to separate
- Other?

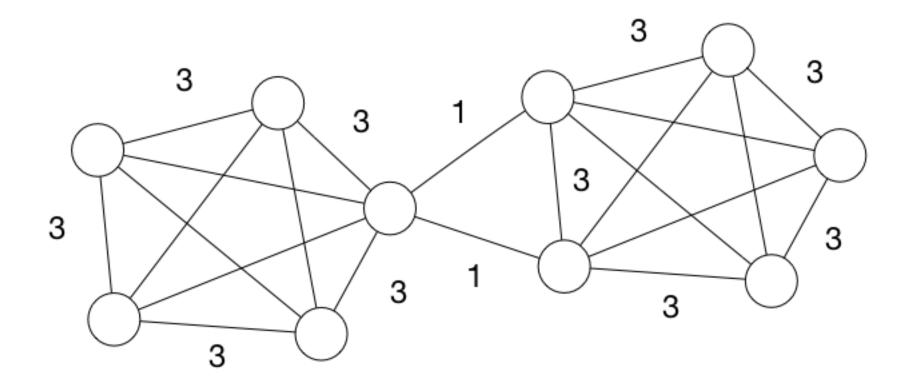


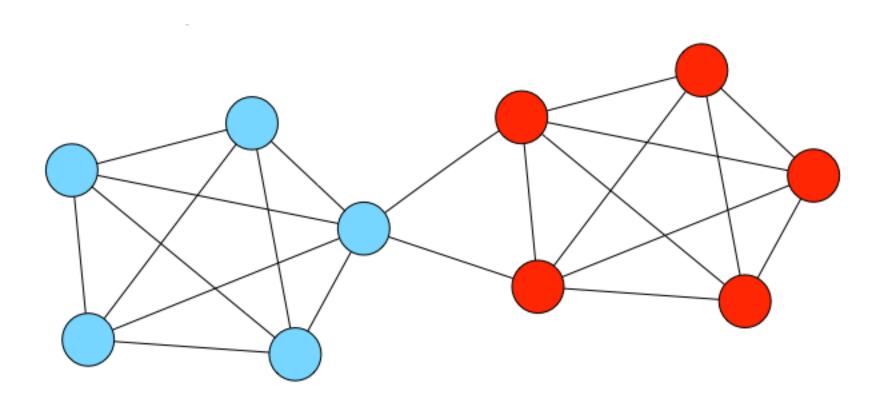




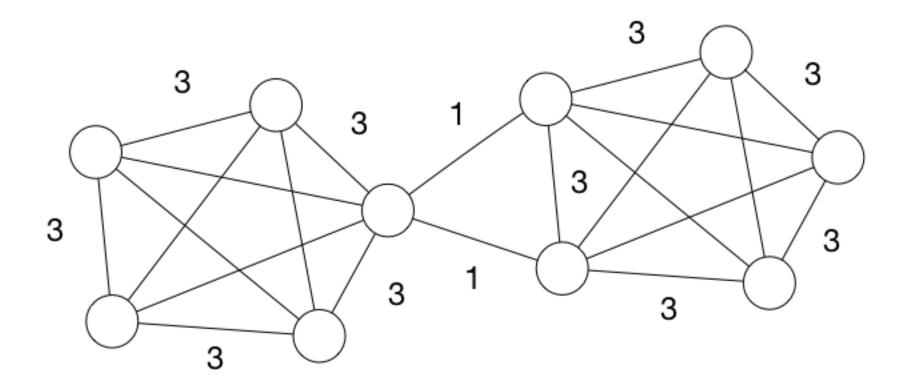








create clusters based on threshold



Some other clustering algorithms

- k-Spanning Tree
- Betweenness Centrality
 Based
- Highly Connected Components
- Maximal Clique Enumeration
- The Markov Cluster Algorithm
- HCS algorithm

QUESTIONS

