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I pledge my honor that I have abided by the Stevens Honor System. - Joshua Schmidt

Point values are assigned for each question.

1. Find an upper bound for  $f(n)=n^4+10n^2+5$ . Write your answer here: \_\_\_O( $n^4$ )\_\_\_\_ (4 points)

Prove your answer by giving values for the constants c and  $n_0$ . Choose the smallest integral value possible for c. (4 points)

$$f(n)=n^4+10n^2+5<2n^4, n_0=4, c=2$$

2. Find an asymptotically tight bound for  $f(n)=3n^3-2n$ . Write your answer here:  $-\Theta(n^3)$ \_ (4 points)

Prove your answer by giving values for the constants  $C_1$ ,  $C_2$ , and  $C_3$ . Choose the tightest integral values possible for  $C_1$  and  $C_2$ . (6 points)

$$f(n)=3n^3-2n \le 3n^3, n_0=1$$
  
 $f(n)=3n^3-2n \ge 2n^3, n_0=2$   
answer:  $c_1=2, c_2=3, n_0=2$ 

3. Is  $3n-4 \in \Omega(n^2)$ ? Circle your answer: yes / no. (2 points)

If yes, prove your answer by giving values for the constants c and  $n_0$ . Choose the smallest integral value possible for c. If no, derive a contradiction. (4 points) contradiction:

$$3n-4 \ge n^2, n=1 \to -1 \ge 1$$

if  $\Omega(n^2)$ , 3n-4 should be greater than  $n^2$ , but the expression above shows it doesn't work for n=1. The function has a value lower than Omega, which should not be possible.

Generalized: 
$$3n-4 \ge cn^2 \rightarrow 3/n-4/n^2 \ge c$$
,  $n>0$ 

4. Write the following asymptotic efficiency classes in **increasing** order of magnitude.

$$O(n^2)$$
,  $O(2^n)$ ,  $O(1)$ ,  $O(nlgn)$ ,  $O(n)$ ,  $O(n!)$ ,  $O(n^3)$ ,  $O(lgn)$ ,  $O(n^n)$ ,  $O(n^2lgn)$  (2 points each)

$$_{O(1)}_{,O(lg\,n)}_{,O(nlg\,n)}_{,O(nlg\,n)}_{,O(n^2lg\,n)}_{,O(n^2lg\,n)}_{,O(n^3lg\,n)}$$

5. Determine the largest size n of a problem that can be solved in time t, assuming that the algorithm takes f(n) milliseconds. n must be an integer. (2 points each)

a. 
$$f(n) = n$$
,  $t = 1$  second \_\_\_\_1000\_\_\_

```
c. f(n) = n^2, t = 1 hour 1897
```

d. 
$$f(n) = n^3$$
,  $t = 1$  day 442

- e. f(n) = n!, t = 1 minute 8
- 6. Suppose we are comparing two sorting algorithms and that for all inputs of size n the first algorithm runs in 4 n³ seconds, while the second algorithm runs in 64 n lg n seconds. For which integral values of n does the first algorithm beat the second algorithm? \_\_\_\_n=2-6\_\_\_\_\_\_ (4 points) Explain how you got your answer or paste code that solves the problem (2 point): \_\_\_ I tried different values of n, starting at n=1 (undefined), until the runtime for the first algorithm became larger than the second algorithm. This occurred when n=7. Then I concluded that n must be between 2 and 6 in order for the statement to work. I essentially found the intersection between the two functions.
- 7. Give the complexity of the following methods. Choose the most appropriate notation from among  $O, \Theta$ , and  $\Omega$ . (8 points each)

```
int function1(int n) {
    int count = 0;
    for (int i = n / 2; i <= n; i++) {</pre>
         for (int j = 1; j <= n; j *= 2) {
             count++;
    return count;
Answer: \underline{\Theta}(nlgn)
int function2(int n) {
    int count = 0;
    for (int i = 1; i * i * i <= n; i++) {</pre>
         count++;
    return count;
Answer: \Theta(n^{(1/3)})
int function3(int n) {
    int count = 0;
    for (int i = 1; i <= n; i++) {</pre>
         for (int j = 1; j <= n; j++) {</pre>
             for (int k = 1; k <= n; k++) {
                  count++;
         }
    }
```

```
return count;
Answer: \underline{\Theta}(n^3)___
int function4(int n) {
     int count = 0;
     for (int i = 1; i <= n; i++) {</pre>
          for (int j = 1; j <= n; j++) {
                count++;
                break;
           }
     }
     return count;
}
Answer: \underline{\hspace{0.1cm}}\Theta(n)\underline{\hspace{0.1cm}}
int function5(int n) {
     int count = 0;
     for (int i = 1; i <= n; i++) {</pre>
          count++;
     for (int j = 1; j <= n; j++) {</pre>
          count++;
     return count;
}
Answer: \underline{\Theta}(n)
```