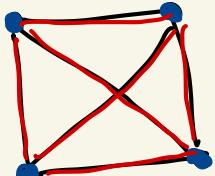



HW 9 due this Friday! (has before the break material)

Section 2.6 Graph Theory

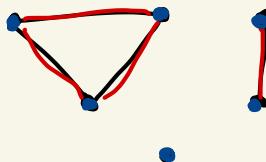
Define: A graph G is a collection of vertices and edges. A vertex v is represented by a dot and an edge e is a line segment between two vertices.

Ex



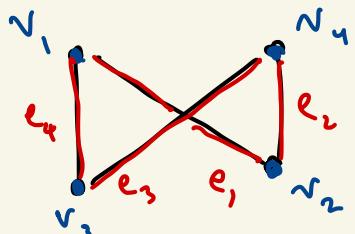
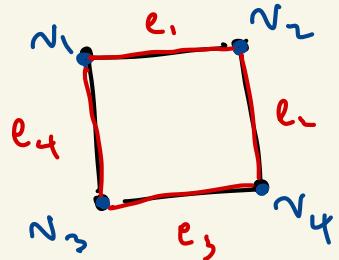
4 vertices ✓
6 edges ✓

Ex



6 vertices
4 edges

There are 2 different ways to draw the same graph.



All that matters
is which
edges are connected
to which vertices!

Edges have no direction!

$$G = (E, V)$$

$$V = \{v_1, \dots, v_n\}$$

E = sets of pairs from V

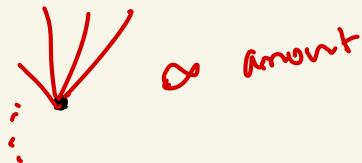
$$V = \{v_1, v_2, v_3, v_4\}$$

$$E = \left\{ \{v_1, v_2\}, \{v_2, v_4\}, \{v_3, v_4\}, \{v_1, v_3\} \right\}$$

Typically : Only a finite amount of edges or vertices



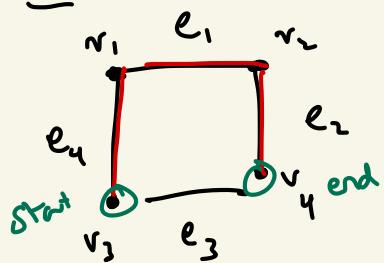
If you want to consider infinite graphs, you might assume only finitely many edges connected to a vertex.



We will only consider finite graphs!

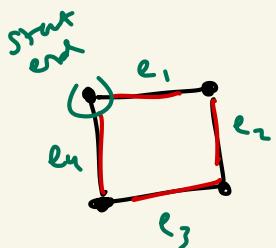
Def: A path within a graph G is a series of edges e_1, \dots, e_n such that e_i and e_{i+1} share a vertex.

Ex



e_4, e_1, e_2 is a path

e_4 and e_1 share vertex 1 in common
Cuz
 e_4 and e_1 share vertex 1 in common
Not a circuit



e_1, e_2, e_3, e_4 is another path

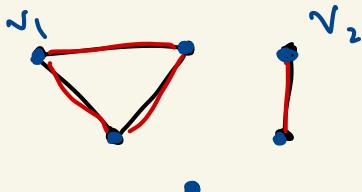
This is a circuit.

Def: A path in G is called a circuit if it starts

and ends at the same vertex.

Def : We say a graph is connected if there exists a path from any vertex to any other vertex.

Ex



This is NOT a connected graph.
There's no path from v_1 to v_2 .



This smaller graph is connected!

Note :



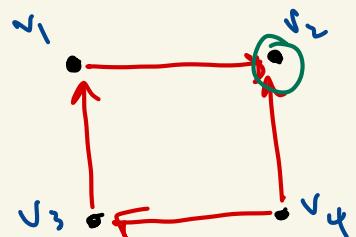
In some books, this would be considered a graph. In others this would be a multigraph.

We won't consider more than 1 edge between two vertices.

Diagraphs Now we add directions to our edges.

Def: A diagraph is a graph but with additional information about where edges start and end.

$$e = (v_i \rightarrow v_j) \text{ Diagraph}$$

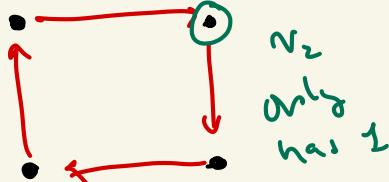


In diagraphs,

via arrows, edges are represented by arrows.

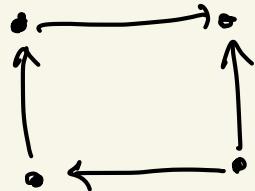
$$e = \{v_i, v_j\} \text{ graph}$$

Both of these
are connected.



Then are two different diagraphs.

v2
only
has 1



is connected

since the underlying graph is connected.

Despite the fact there's no path out of v_2 .

Def Given a digraph G , we can associate 2 vector spaces to

b. Suppose G has vertices $v_1 \dots v_n$
edges $e_1 \dots e_m$

$C_1 = \text{all } \boxed{\text{formal}} \text{ linear combinations of } e_1 \dots e_m$.

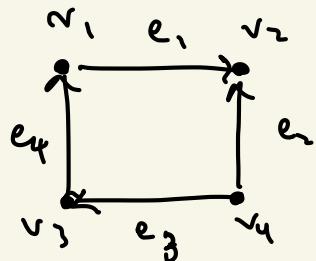
$$e_1 = 1e_1 + 0e_2 + \dots + 0e_m \in C^1$$

$$2e_1 + \frac{1}{2}e_2 + 5e_3 \in C^1.$$

formal means that
 $e_1 \dots e_m$
is a "basis"
in an abstract sense

If somehow $e_1 \dots e_m$ were a basis to a vector space,
 what would that vector space look like? You get C_1 .

$C_0 = \text{all } \underline{\text{formal}} \text{ linear combinations of } v_1 \dots v_n.$

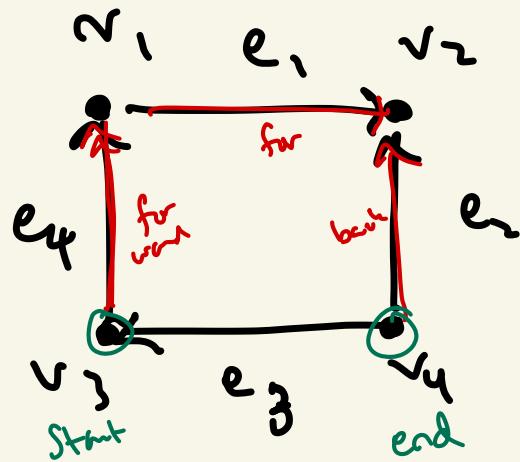


$$C_1 = \text{Span}(e_1, e_2, e_3, e_4)$$

$$C_0 = \text{Span}(v_1, v_2, v_3, v_4)$$

We can just turn vertices and edges in basis vectors
 of a vector space arbitrarily.

We can write paths / circuits as certain linear combinations
 of edges.



e_4, e_1, e_2 is a path.

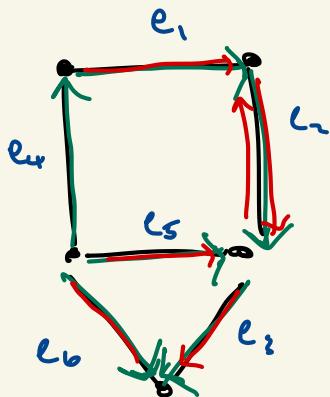
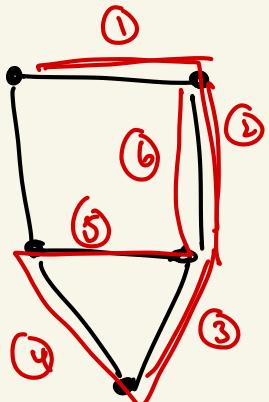
$$e_4 + e_1 - e_2$$

↓

backwards!

- + represents going forward
- represents going backward.

This path $e_4 + e_1 - e_2 \in C_1$ lies in C_1 .



So the path — as a

linear combination is

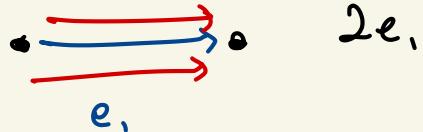
$$e_1 + e_2 + e_3 - e_5 + e_6 - e_2 \\ = e_1 + e_3 + e_5 - e_6$$

$$= e_1 + e_5 + e_3 - e_6$$

$$= -e_6 + e_1 + e_3 + e_5$$

i) Label edges

ii) Add some directions | make this
a digraph



$$-x^*{}^T f \leq 0$$

$$x^* = K^{-1}f$$

$$-x^*{}^T f = -f^T \underbrace{(K^{-1})^T}_{\text{por def}} f \leq 0$$

$$-f^T(K^{-1})^T f \leq 0 \quad \text{for } K \text{ pos def}$$