ThinLTO Summaries for Incremental JIT Compilation

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LTO in a Nutshell

- → LTO = Link Time Opt. = Whole Program Opt.
 = Cross-Translation-Unit Optimization = Monolithic LTO
- Compilers can only inline code that they see, e.g.:
 C++ definitions from included headers, but not from cpp's
- → Linkers can have the full picture, but traditional LTO is expensive and hard to scale

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Requires IR code in object files

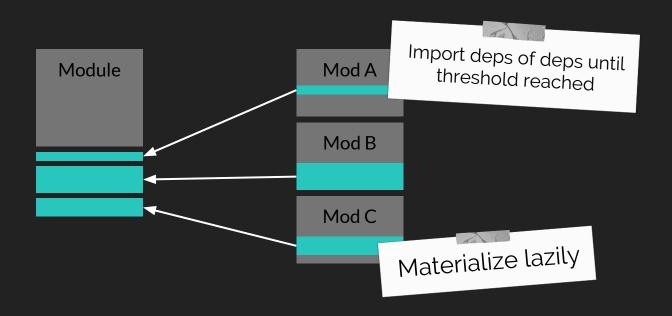
https://llvm.org/docs/BitCodeFormat.html #native-object-file-wrapper-format

ThinLTO in a Nutshell

- Compiler emits summaries with call graph info to object files,
- Linker reads all summaries and builds global combined index for Interprocedural Analysis (Thin-Link) based on GUIDs
- → Then runs ThinLTO Backend per module:
 - Import function definitions from other modules
 - Execute Interprocedural Transformations

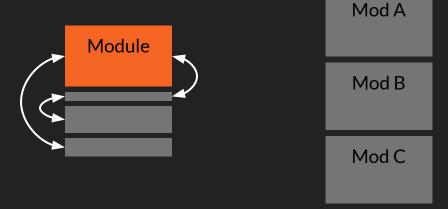
ThinLTO Backends

→ 1. Import functions + dependencies from other modules



ThinLTO Backends

→ 2. Run inlining and IPO



ThinLTO Backends

→ 3. Drop imports



ThinLTO Benefits

- → Small overhead: compact summaries, fast combining
- → Good scalability: actual transformations run in parallel
- → Incremental: simple object file caching + your linker's magic
- → Speedup: close to Monolithic LTO
- → Development: less manual optimization avoids discussions on where to put a definition and can reduce review times dramatically :-)

ThinLTO Command Line

→ Compile:

```
$ clang -flto=thin -02 file1.c file2.c -c
```

→ Link:

```
$ clang -flto=thin -02 file1.o file2.o -o a.out
```

→ All in one:

```
$ clang -flto=thin -02 file1.c file2.c -o a.out
```

More Info

ThinLTO: Scalable and Incremental Link-Time Optimization

Teresa Johnson @ CppCon 2017

https://www.youtube.com/watch?v=p9nH2vZ2mNo

ThinLTO: Scalable and Incremental LTO

Teresa Johnson, LLVM Project Blog 2016

http://blog.llvm.org/2016/06/thinlto-scalable-and-incremental-lto.html

Clang ThinLTO Documentation https://clang.llvm.org/docs/ThinLTO.html

Incremental Compilation in a Nutshell

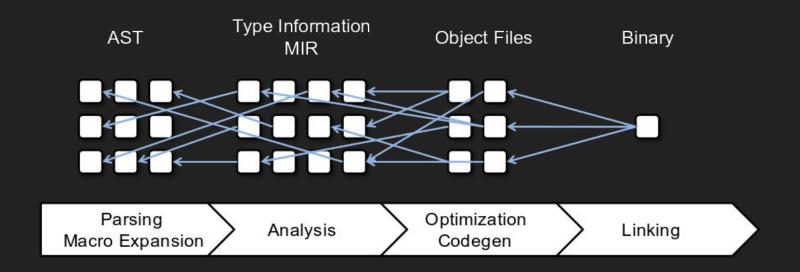
- → Rebuild only what's necessary
- → Terminology:
 - o Incremental Linking ≈ translation unit granularity
 - o Incremental Compilation = higher gran. / smaller entities
- → Complexity in languages with C-like translation units:
 - Incremental Linking << Incremental Compilation

Incremental Compilation in a Nutshell

- → Rebuild only what's necessary
- Terminology:
 - o Incremental Linking ≈ translation unit granularity
 - o Incremental Compilation = higher gran. / smaller entities
- Con Dynamic linking to ↓turnaround times! e.a. libLI_VM
 #Entities↑ ⇒ Overhead↑
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 ⇒ Eats up benefits

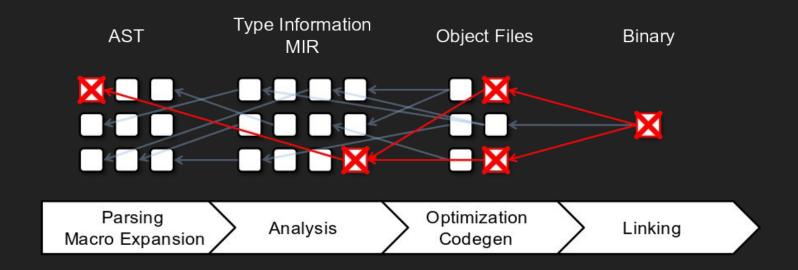
Incremental Compilation in the rustc Compiler

→ AST node granularity



Incremental Compilation in the rustc Compiler

→ AST node granularity gets complex quickly



Incremental Compilation in the rustc Compiler

- → AST node granularity gets complex quickly
- Compile-time savings hard to judge: "Up to 15% or More" *
- → Conceptually promising, kind-of-academic research

More Info

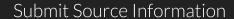
Incremental Compilation

Michael Woerister, Rust Programming Language Blog 2016 https://blog.rust-lang.org/2016/09/08/incremental.html

JIT Compilation

- → JIT Compiler does both, compile and link
 - → Source representation is LLVM IR
 - → Incremental Compilation ≈ Incremental Linking
 No simple separation

"Monolithic" JIT Compilation

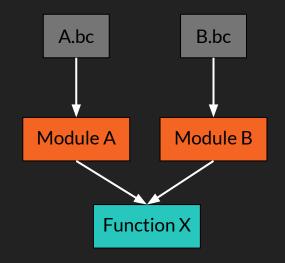


E.g. file names, ninja build description, etc.

Request Callable Function

E.g. mangled name or name + signature

Function Handle / Error

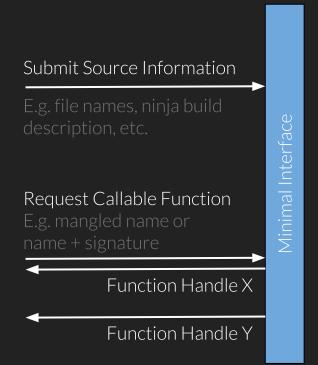


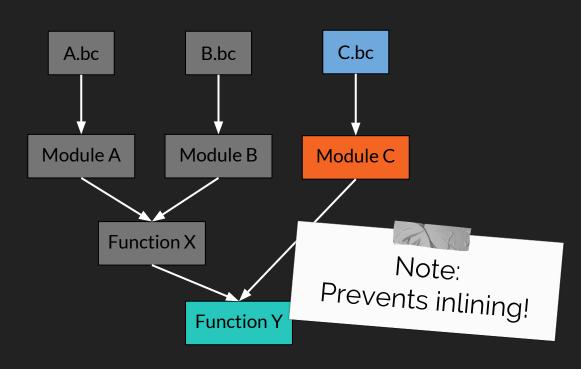
C.bc

Compile relevant sources into native code modules.

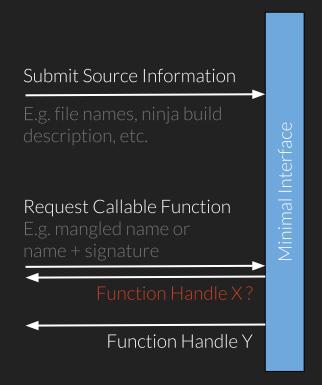
Link native code modules, run static initializers and obtain function pointer.

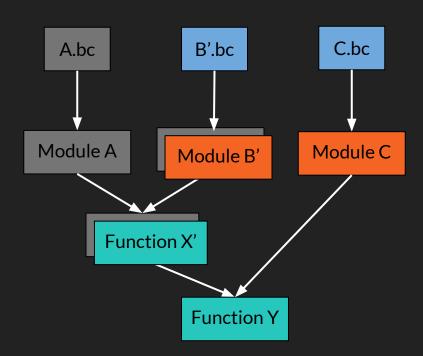
Incremental JIT Compilation





Incremental Just-in-Time Compilation





Incremental JIT Compilation



- Requires quite some housekeeping
 - → Track file changes on disk

Source Module Hashes

→ Determine source dependencies for functions



- → Decide what needs to be recompiled
- → Link against existing code
- → Determine static initializers to rerun
- → Invalidate existing function handles (and rebind?)

Incremental JIT Compilation

→ Requires quite some housekeeping

- → More importantly we also run the code!
 - → Can we get Stateful Incremental JIT Compilation?

Stateful Incremental JIT Compilation



- Requires even more housekeeping
 - → Track file changes on disk

Source Module Hashes

→ Determine source dependencies for functions



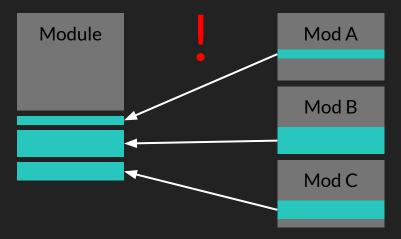
- → Decide what needs to be recompiled
- → Link against existing code



- → Rename duplicate globals and resolve to existing ones
- → Invalidate existing function handles (and rebind?)

Name Promotion in ThinLTO Backends

Names of imported globals must not clash with local names!



Name Promotion in ThinLTO Backends

main.cpp

```
int readGlobal();
void writeGlobal(int x);

static int GlobalVar = 10;

int main() {
  writeGlobal(1);
  return readGlobal() + GlobalVar;
}
```

```
$ clang main.cpp global.cpp -o test
$ ./test
$ echo $?
11
```

Name collision on import

global.cpp

```
static int GlobalVar = 0;
int readGlobal() { return GlobalVar; }
void writeGlobal(int x) { GlobalVar = x; }
```

Name Promotion in ThinLTO Backends

Imported pseudo-code main.cpp

```
int readGlobal() { return GlobalVar.llvm.98887675656456543548795; }
void writeGlobal(int x) { GlobalVar.llvm.98887675656456543548795 = x; }

static int GlobalVar = 10;
extern int GlobalVar.llvm.98887675656456543548795;

int main() {
    writeGlobal(1);
    return readGlobal() + GlobalVar;
}

Postfix with source
    module hash
```

→ Like the regular linker, JIT can use hash to find module with the symbol definition

global.cpp

```
#include <cstdio>
static int GlobalVar = 0;
int readGlobal() { return GlobalVar; }

void writeGlobal(int x) {
  printf("Modified GlobalVar");
  GlobalVar = x;
}
```

→ Modification changes module hash

Adjusted pseudo-code global.cpp

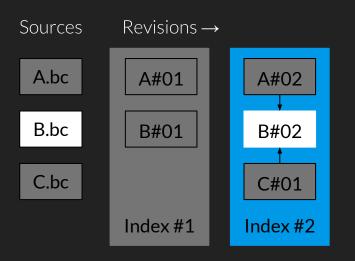
```
#include <cstdio>
extern int GlobalVar.llvm.98887675656456543548795;
int readGlobal() { return GlobalVar.llvm.98887675656456543548795; }

void writeGlobal(int x) {
  printf("Modified GlobalVar");
  GlobalVar.llvm.98887675656456543548795 = x;
}
```

Manually rename own global, so it can be resolved to existing instance in previous revision

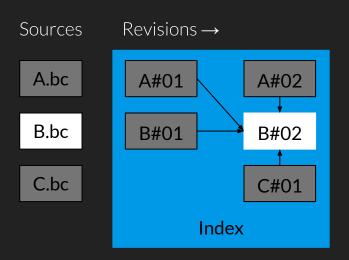
→ ThinLTO tools tailored for ahead-of-time compilation, e.g.

ModuleSummaryIndex uses module ID as key and not Hash



- Cannot (easily) keep one Index over multiple revisions.
- Diff summaries to obtain on-disk changes during Thin-Link
- Promote stateful globals manually in Thin Backends (lots of extra maintenance!)

→ If ModuleSummaryIndex keys were hashes, we could easily import from previous revisions:



- → Other ThinLTO implementations used:
 - → llvm::getLazyBitcodeModule()
 - → llvm::readModuleSummaryIndex()
 - → llvm::computeDeadSymbols()
 - Ilvm::ComputeCrossModuleImportForModule()

Overlap: ThinLTO ↔ (Stateful) Incremental JIT

- → Require global view of the program
- → Computationally expensive & huge scale
- → Run in two stages
 - → Global analysis (sequential)
 - → Per-module recompilation (in parallel)
- → Focus on optimizations in the compiled program
 - → JIT: keep acceptable runtime with high granularity!

Short-Term Goals

- → Hack Clang to emit Hashes in ModuleSummaryIndex.
- → Build infrastructure upon ThinLTO implementations that directly targets stateful incremental recompilation to get rid of hacky workarounds.
- → Extend management infrastructure, e.g. handle cases like function signature changes.
- → Load ninja.build descriptions instead of single files.

- → Git-like revision management for program state
 - → Fast-forward merge side-effect-free changes
 - → Let developer resolve conflicts
 - → Allow saving and restoring state snapshots
 - → Fall back to restart on fail

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- Emit program summary data to executable (or extra file)
 - → Like debug information for program state

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- Emit program summary data to executable (or extra file)
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- → Package JIT as dylib, so it can be used like a sanitizer
- → Implement support in common developer tools (e.g. LLDB) to modify executables based on program summary data

Thanks!

Let's dive into the code!

https://github.com/weliveindetail/ThinLtoJit