

# **15-887: Assignment #2**

Due on Wednesday, October 19, 2016

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## Problem 1

(a)

To run this planner, please use the following command: `python world.py -a prob1.txt`

This planner performs a forward djikstra search in x and y and uses the result as an informed heuristic for a single backward A\* search in x, y, and t.

The 2D djikstra search expands all nodes in the graph - one million nodes for test case one. On my machine, it takes about 8 seconds to complete this phase of the search. The distances calculated by the djikstra search are then used as the heuristic for a backward A\* search which searches in both position and time.

**TODO: Time, cost, number of states expanded**

(b)

To run this planner, please use the following command: `python world.py -b prob1.txt`

One of the advantages of the algorithm I chose for part (a) is its ease of extension to part (b). By weighting the backward A\* search, it is easy to trade optimality for speed, and it comes with the typical weighted A\* performance guarantees.

**TODO: Time, cost, number of states expanded**

## Problem 2

### 2.1

Suppose you have two consistent heuristic functions:  $h_1$  and  $h_2$ . Prove that  $h(s) = \max(h_1(s), h_2(s))$  for all states  $s$  in the graph is also a consistent heuristic function.

A heuristic is consistent if

$$h(n) \leq c(n, n') + h(n')$$

for every node  $n$  and its child node  $n'$ .

Proof:

$$\begin{aligned} h(n) &= \max(h_1(n), h_2(n)) \\ &\leq \max(c(n, n') + h_1(n'), c(n, n') + h_2(n')) \\ &\leq c(n, n') + \max(h_1(n'), h_2(n')) \\ &\leq c(n, n') + h(n') \end{aligned}$$

Suppose you have two consistent heuristic functions:  $h_1$  and  $h_2$ . Prove that  $h(s) = \min(h_1(s), h_2(s))$  for all states  $s$  in the graph is also a consistent heuristic function.

A heuristic is consistent if

$$h(n) \leq c(n, n') + h(n')$$

for every node  $n$  and its child node  $n'$ .

Proof:

$$\begin{aligned} h(n) &= \min(h_1(n), h_2(n)) \\ &\leq \min(c(n, n') + h_1(n'), c(n, n') + h_2(n')) \\ &\leq c(n, n') + \min(h_1(n'), h_2(n')) \\ &\leq c(n, n') + h(n') \end{aligned}$$

## 2.2

d. Monotonically non-increasing sequence

## 2.3

f. None of the above

## 2.4

e. None of the above