

CS5002. Prof Higger

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Problem 1

$$\begin{aligned} \text{i. } 753 &= 376 \times 2 + 1 \\ 376 &= 188 \times 2 + 0 \\ 188 &= 94 \times 2 + 0 \\ 94 &= 47 \times 2 + 0 \\ 47 &= 23 \times 2 + 1 \\ 23 &= 11 \times 2 + 1 \\ 11 &= 5 \times 2 + 1 \\ 5 &= 2 \times 2 + 1 \\ 2 &= 1 \times 2 + 0 \\ 1 &= 0 \times 2 + 1 \end{aligned}$$

$$\Rightarrow 753_{10} = 101111000_2$$

$$\begin{aligned} \text{ii. } (10100101)_2 &= 1 \times 2^7 + 0 \times 2^6 + 1 \times 2^5 + 0 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= 128 + 32 + 4 + 1 \\ &= 165 \end{aligned}$$

$$\begin{aligned} \text{iii. } 45263 &= 2828 \times 16 + 15 \quad (F) \\ 2828 &= 176 \times 16 + 12 \quad (C) \\ 176 &= 11 \times 16 + 0 \\ 11 &= 0 \times 16 + 11 \quad (B) \end{aligned}$$

$$\Rightarrow 45263_{10} = B0CF_{16}$$

$$\begin{aligned} \text{iv. } (BFF)_{16} &= 11 \times 16^2 + 15 \times 16^1 + 15 \times 16^0 \\ &= 2816 + 240 + 15 \\ &= 3071 \end{aligned}$$

Problem 2.

$$i. \quad 179_{10} = 2^7 + 2^5 + 2^4 + 2^1 + 2^0 = 10110011_2$$

$$55 = 2^5 + 2^4 + 2^2 + 2^1 + 2^0 = 00110111_2$$

$$223 = 2^7 + 2^6 + 2^4 + 2^3 + 2^2 + 2^1 + 2^0 = 11011111_2$$

$$12 = 2^3 + 2^2 = 00001100$$

the binary-hex convert table as follows:

0000	0	0100	4	1000	8	1100	C
0001	1	0101	5	1001	9	1101	D
0010	2	0110	6	1010	A	1110	E
0011	3	0111	7	1011	B	1111	F

$$(101100110011011110111100001100)_2 = (B337DF0C)_{16}$$

$$= 3006783244_{10}$$

$$ii. \quad BAC2A78F = 10111010.11000010.10100111.10001111$$

$$= 186.194.167.143 \dots \dots \text{dotted decimal}$$

$$= 3133319055 \dots \dots \text{decimal}$$

Problem 3

$$i. \quad 10000000_2 \Rightarrow 01111111 \Rightarrow 10000000 = -2^7 = -128$$

$$11110011_2 \Rightarrow 00001100 \Rightarrow 00001101 = -(2^3 + 2^2 + 2^0) = -13$$

$$01111111_2 \Rightarrow 2^6 + 2^5 + 2^4 + 2^3 + 2^2 + 2^1 + 2^0 = 2^7 - 1 = 127$$

$$ii. \quad 55 \Rightarrow 2^5 + 2^4 + 2^2 + 2^1 + 2^0 = 00110111$$

$$83 \Rightarrow 2^6 + 2^4 + 2^1 + 2^0 = 01010011$$

$$79 = 2^6 + 2^3 + 2^2 + 2^1 + 2^0 = 01001111 \xRightarrow{\text{flip}} 10110000 \xRightarrow{+1} 10110001$$

$$\therefore -79 = 10110001_2$$

$$88 = 2^6 + 2^4 + 2^3 = 01011000 \Rightarrow 10100111 \Rightarrow 10101000$$

$$\therefore -88 = 10101000_2$$

ii. $-79 + 55 = -24$

$$\begin{array}{r} -79 \\ + 55 \Rightarrow \\ \hline 10110001 \\ + 00110111 \\ \hline 11101000 \end{array}$$

$$11101000 \Rightarrow 00010111 \Rightarrow 00011000 = -(2^4 + 2^3) = -24$$

correct !

It's not an overflow

$$-79 - 88 = -167$$

$$\begin{array}{r} -79 \\ + -88 \Rightarrow \\ \hline 10110001 \\ + 10101000 \\ \hline 101011001 \end{array}$$

It's an overflow

$$101011001 \Rightarrow 010100110 \Rightarrow 010100111 = -(2^7 + 2^5 + 2^2 + 2^1 + 2^0) = -167$$

$$83 + 55 = 138$$

$$\begin{array}{r} 83 \\ + 55 \Rightarrow \\ \hline 01010011 \\ + 00110111 \\ \hline 10001010 \end{array} = 10001010 = 2^7 + 2^3 + 2^1 = 138$$

It's an overflow

Problem 4.

1. When in 8-bit two's complement, the range is $[-128, 127]$.

So we can figure in 10-bit two's complement, $\Rightarrow [-2^9, 2^9 - 1]$. NEZ

the range is $[-2^9, 2^9 - 1] \Rightarrow [-512, 511]$.

2. $\because 2^{11} < 2100 < 2^{12}$ ($2^{11} = 2048, 2^{12} = 4096$)

We need the range $[-2^{12}, 2^{12} - 1]$ to represent 2100

\therefore the minimum number of bits = 13

Problem 5.

$$\begin{array}{r} i. \quad \begin{array}{r} 100111 \\ + \quad 1010 \\ \hline 110001 \end{array} \end{array}$$

$$\begin{array}{r} ii. \quad \begin{array}{r} 110111 \\ + \quad 1011 \\ \hline 1000010 \end{array} \end{array}$$

$$\begin{array}{r} iii. \quad \begin{array}{r} 101000 \\ \times \quad 1101 \\ \hline 101000 \\ 000000 \\ 101000 \\ \hline 101000 \\ \hline 1000001000 \end{array} \end{array}$$

$$\begin{array}{r} iv. \quad \begin{array}{r} 110111 \\ \times \quad 1100 \\ \hline 000000 \\ 000000 \\ 110111 \\ \hline 110111 \\ \hline 101001000 \end{array} \end{array}$$

Problem 6.

$$i. \quad 8 \bmod 4 = 0 \Rightarrow x = 0$$

$$ii. \quad x \bmod 5 = 3 \Rightarrow x = 5n + 3, \quad n \in \mathbb{N}. \quad \text{eg. } x \in \{3, 8, 13, 18, \dots\}$$

$$\begin{aligned} iii. \quad 145 + 174 \bmod 17 &= x \\ \Rightarrow 319 \bmod 17 &= 13 \Rightarrow x = 13 \end{aligned}$$

$$\begin{aligned} iv. \quad 145 \bmod 17 &= 9 \\ 174 \bmod 17 &= 4 \Rightarrow x = 9 + 4 = 13 \end{aligned}$$

$$v. \quad (x+y) \bmod n = x \bmod n + y \bmod n, \quad x, y, n \in \text{integer}$$

It is always true.

$$\begin{aligned} \text{Prove: assume } x \bmod n &= k_1 \Rightarrow x = n \times a + k_1 & k_1, k_2 < n \\ y \bmod n &= k_2 \Rightarrow y = n \times b + k_2 & \therefore a+b = A+B \\ \Rightarrow (x+y) \bmod n &= k_3 = x+y = n \times (A+B) + k_3 & k_1+k_2 = k_3 \\ \Rightarrow n \times a + k_1 + n \times b + k_2 &= n \times (A+B) + k_3 \end{aligned}$$

To add some explanation here, the upper equation shows that $n(a+b) + (k_1+k_2) = n(A+B) + k_3$, from linear algebra we know that only if the coefficients are the same can the two line be coincident.

Problem 7.

So the person is his/her own ancestor.

the total number of ancestor is $1+2+4+8 \dots$

$$= (1+10+100+1000 \dots)_2$$

$$\Rightarrow 2^{n+1} - 1 = (\underbrace{11111 \dots 1}_{n+1})_2$$

So each "1" in the binary format represents

the number added when generation n is added by 1.