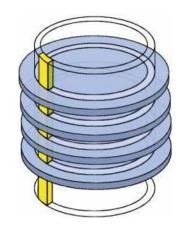




Conception Avancée de Bases de Données







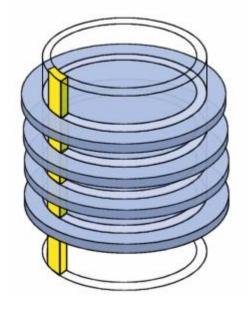






Disque

Mémoire

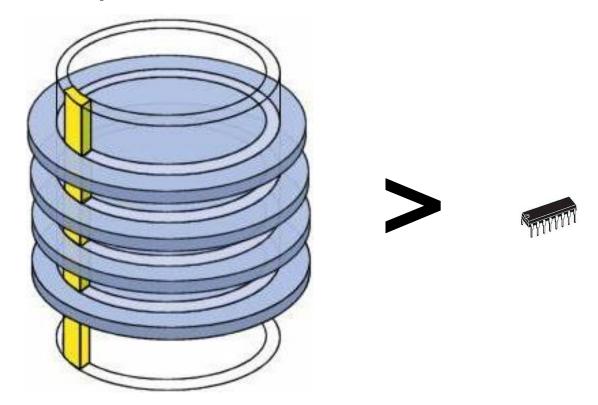




Problème.



• Que se passe-t-il si la taille d'une relation est plus grande que la taille de la mémoire ?



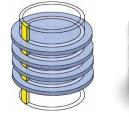
Problème.



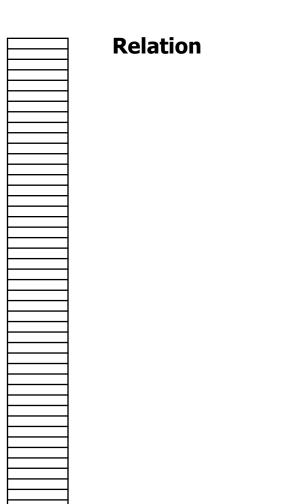
• Que se passe-t-il si la taille d'une relation est plus grande que la taille de la mémoire ?

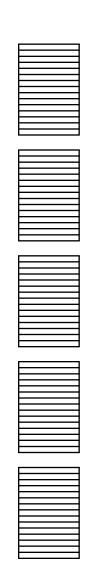
- Il faut découper la relation pour réaliser les opérations par « segments »
- Un SGBD gère lui-même l'espace disque comme, à la place, du système d'exploitation.
 - Les relations sont stockées sur disque sous forme de segments non contigus.

Relations sur disque disques



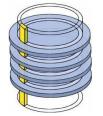




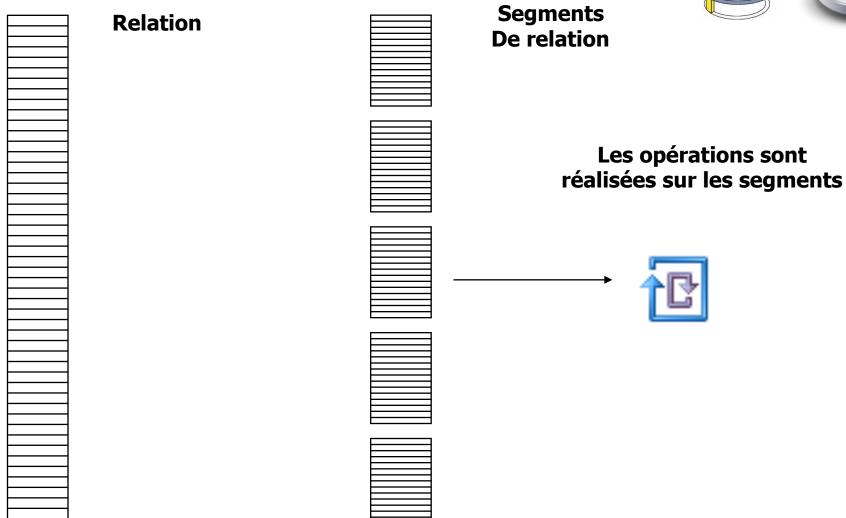


Segments De relation

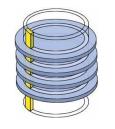
Relations sur disque



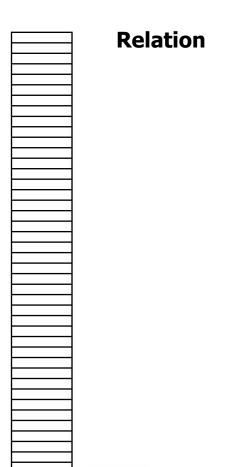




Relations sur disque

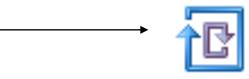








Les opérations sont réalisées sur les segments

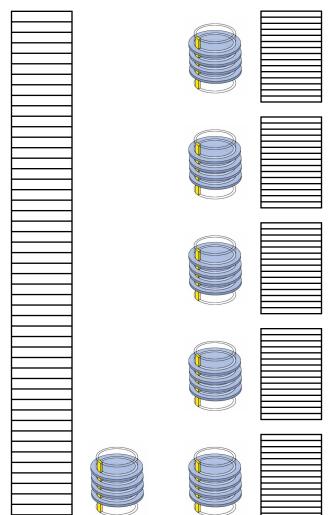


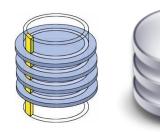


Chaque segment Est chargé en mémoire Au cours de l'opération

Relations sur disque

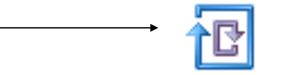
Relation



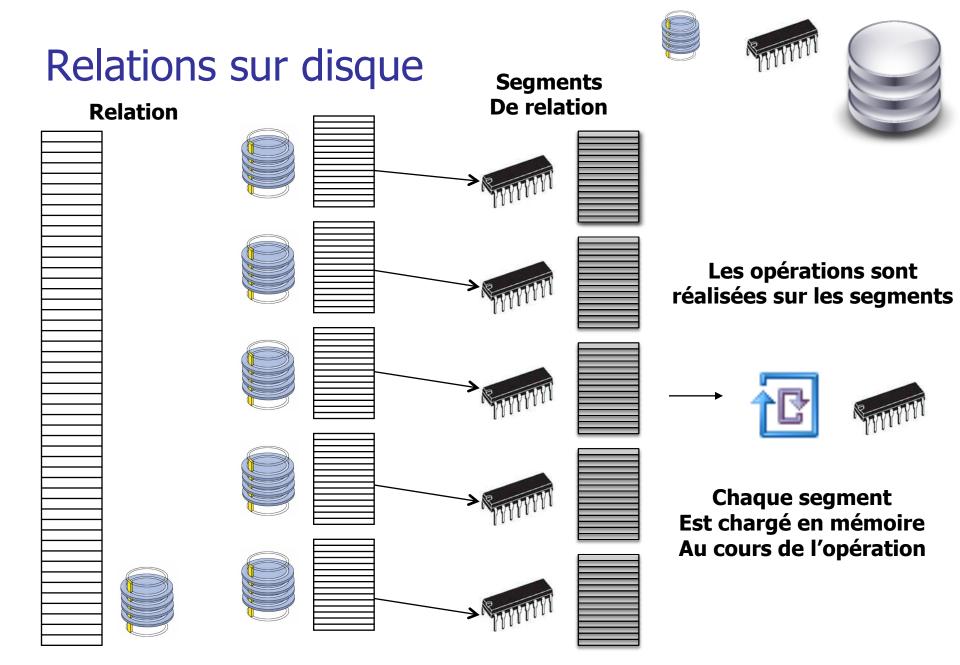


Segments De relation

Les opérations sont réalisées sur les segments



Chaque segment Est chargé en mémoire Au cours de l'opération



Memory join

Nested loop

Merge join

Hash join

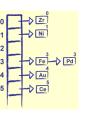














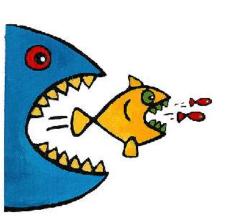
Disk algorithms

- Sort merge
- TBC



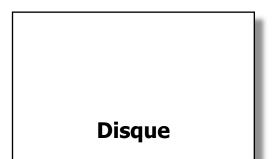


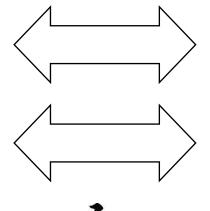




Mixed storage space

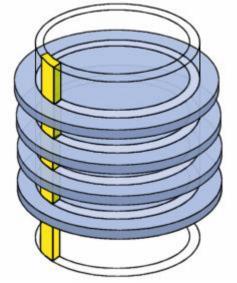






Mémoire







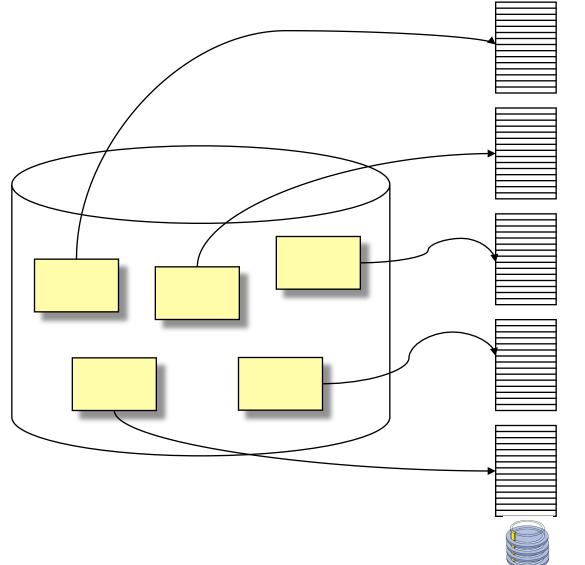
Secondary Memory

Relations et blocks disques

Blocks disques

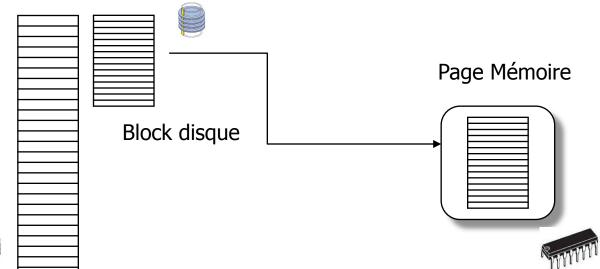






Relations sur disque et en mémoire







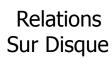


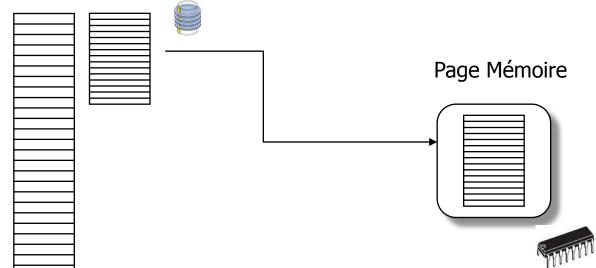


Tailles blocs disques pages mémoires













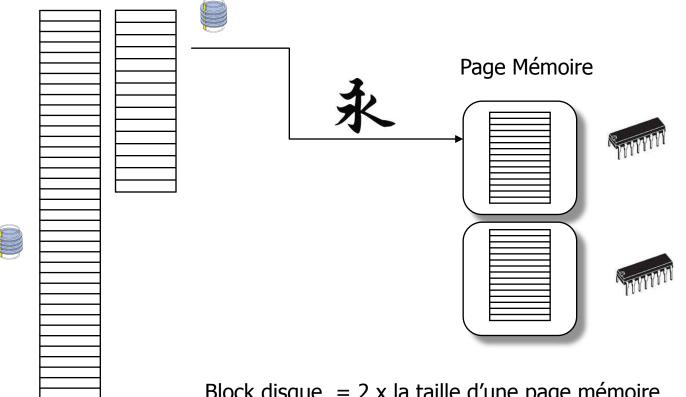
Talle Block disque = Taille page mémoire

Tailles blocs disques pages mémoires





Relations Sur Disque





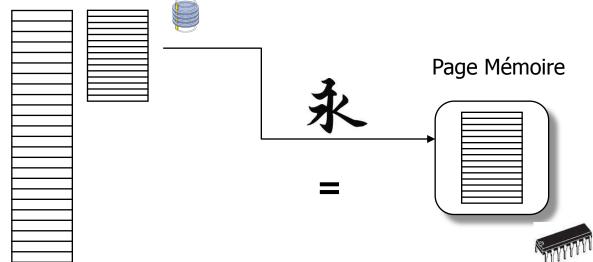
Block disque = $2 \times la$ taille d'une page mémoire

Tailles blocs disques pages mémoires

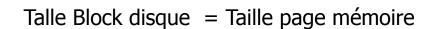












Tailles blocs disque page mémoire







