DIGITAL SYSTEM DESIGN APPLICATIONS

(CRN: 11275)

THE REPORT OF EXPERIMENT – 7



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1. Circuit That Detects Four Consecutive 1 or 0

Finite State Machine (FSM) Encoding Methodes:

• **Binary Encoding:** In state machines, there could be a variety of states. In HDL coding these states are represented by an encoding number. Binary encoding method requires assigning these numbers one by one to minimize the length of the state vector, which is good for FPGA and PLA designs.

State	Value					
Idle	000					
r1	001					
r2	010					
r3	011					
r4	100					
c	101					
p1	110					
p2	111					

Binary Encoded State Values Table¹

• **Gray Encoding:** In gray encoding, only one-bit changes when moving between various states. As a result of this behavior, gray encoding consumes less power than binary coding. To explain furthermore, between each state jump, there are transistors on and off to represent 1s and 0s. This change in transistors causes the parasitic capacitors of the transistors to charge or discharge which requires power. With gray encoding, all the state jumps done by changing only one-bit that means the unchanged bits have not consumed any energy. Also, it can reduce glitches in the implementation.

State	Value					
Idle	000					
r1	001					
r2	011					
r3	010					
r4	110					
c	111					
p1	101					
p2	100					

Gray Encoded State Values Table¹

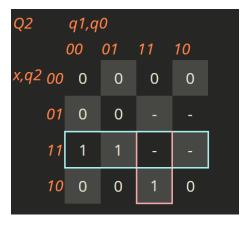
• One-Hot Encoding: One-Hot Encoding simplifies the circuits above via using one bit for each state. This simplifies the circuit and reduces propagation delays which results in making the circuit compatible with higher frequency clocks. However, the trade-off is that one-hot encoding increases the number of flip-flops used to store the state of the system.

State	Value				
Idle	0000001				
ri	0000010				
r2	00000100				
r3	00001000				
r4	00010000				
С	00100000				
p1	01000000				
p2	10000000				

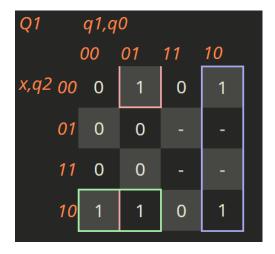
One-Hot Encoded State Values Table¹

state	binary code
Α	000
В	001
C	010
D	011
E	100
F	101

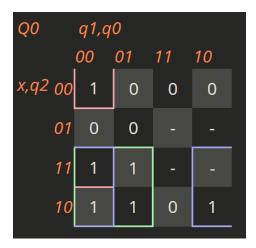
Binary Encoding



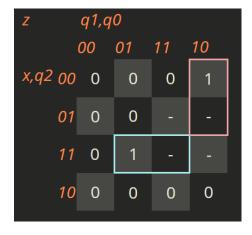
Karnaugh Map of Q2



Karnaugh Map of Q1



Karnaugh Map of Q0



Karnaugh Map of Output z

REDUCED EQUATIONS

$$z(x,q2,q1,q0) = x'q1q0' + xq2q0$$

$$Q2(x,q2,q1,q0) = xq1q0 + xq2$$

$$Q1(x,q2,q1,q0) = q2'q1'q0 + q1q0' + xq2'q1'$$

$$Q0(x,q2,q1,q0) = q2'q1'q0' + xq1' + xq0'$$

```
`timescale 1ns / 1ps
module FSM1 (
    input clk,
    input rst,
    input x,
    output z
    );
    reg q0, q1, q2;
    assign z = (-x & q1 & -q0) | (x & q2 & q0);
    always @(posedge clk or posedge rst) begin
        if(rst) begin
             q0 <= 1'b0;
             q1 <= 1'b0;
             q2 <= 1'b0;
        end else begin
             q2 \leftarrow (x & q1 & q0) | (x & q2);
             q1 \leftarrow (-q2 & -q1 & q0) | (q1 & -q0) | (x & -q2 & -q1);
             q0 \leftarrow (-q2 & -q1 & -q0) | (x & -q1) | (x & -q0);
    end
endmodule
```

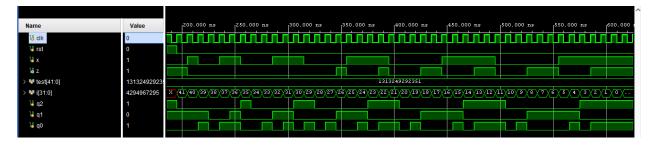
FSM Verilog Code

!!! This is an addition page to report!!!

After making FSM2, I fixed my code for FSM1 but after correcting my code there are not any faulty one or zero occurred in the simulation. Because of that I don't know what to do so, I put the fixed code with simulation in this addition page. (I use the same testbench)

```
module FSM1 (
    input clk,
    input rst,
    input x,
    output z
    );
    reg q0, q1, q2;
    wire Q0, Q1, Q2;
    assign z = (\sim x \& q1 \& \sim q0) | (x \& q2 \& q0);
    assign Q2 = (x & q1 & q0) | (x & q2);
    assign Q1 = (\sim q2 \& \sim q1 \& q0) | (q1 \& \sim q0) | (x \& \sim q2 \& \sim q1);
    assign Q0 = (\simq2 & \simq1 & \simq0) | (x & \simq1) | (x & \simq0);
    always @(posedge clk or posedge rst) begin
         if(rst) begin
             q0 <= 1'b0;
             q1 <= 1'b0;
             q2 <= 1'b0;
         end else begin
             q2 <= Q2;
             q1 <= Q1;
             q0 <= Q0;
         end
    end
endmodule
```

FSM1 Verilog Code (works correctly)



FSM1 Behavioral Simulation (no faulty)

From now on the report continues as wanted order.

```
`timescale 1ns / 1ps
module FSM1_tb();
          clk = 0;
    reg
          rst = 0;
    reg
              = 0;
          X
    reg
          z;
    FSM1 dut(
        .clk(clk),
        .rst(rst),
        .x(x),
        .z(z));
    always #5 clk = ~clk;
    reg [41:0] test =
42'b01 0011 0001 1100 0011 1100 0001 1111 0000 0011 1111;
    integer i;
    initial begin
        repeat(80) @(posedge clk);
        rst = 1;
        repeat(20) @(posedge clk);
        rst = 0;
        for(i=41; i>=0; i=i-1) begin
            x = test[i];
            @(posedge clk);
        end
    end
endmodule
```

FSM Testbench Code



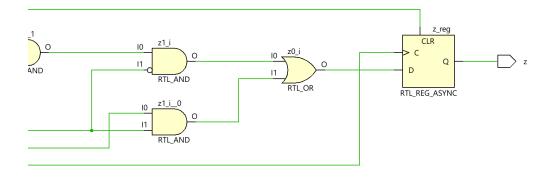
FSM Behavioral Simulation

This design works as a Mealy Machine which resulted in the output to go high within three clock cycles rather than four. This is because of the way Mealy Machine gives output. Whenever the state registered the output changes without waiting the clock which means the device works as a "Three consecutive 1s or 0s". This is rather a design choice.

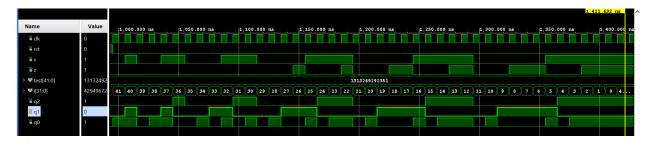
```
`timescale 1ns / 1ps
module FSM1(
    input clk,
    input rst,
    input x,
    output reg z
    reg q0, q1, q2;
    always @(posedge clk or posedge rst) begin
         if(rst) begin
              q0 <= 1'b0;
              q1 <= 1'b0;
              q2 <= 1'b0;
              z <= 1'b0;
         end else begin
              q2 \le (x \& q1 \& q0) | (x \& q2);
              q1 \leftarrow (\sim q2 \& \sim q1 \& q0) | (q1 \& \sim q0) | (x \& \sim q2 \& \sim q1);
              q0 \leftarrow (-q2 \& -q1 \& -q0) | (x \& -q1) | (x \& -q0);
              z \leftarrow (-x & q1 & -q0) | (x & q2 & q0);
         end
     end
endmodule
```

Moore FSM Verilog Code

To convert a Mealy Machine into Moore Machine, the output must be synchronized with the clock. With the code above, the output z can be synchronized with the clock using "always" block. In the RTL Schematic, the DFF connected to the output z can be seen.



RTL Schematic of The Output Z (Moore)

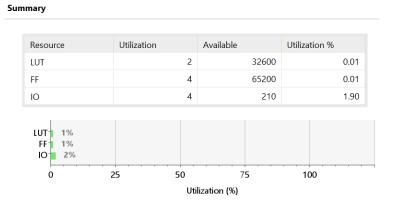


FSM Moore Behavioral Simulation

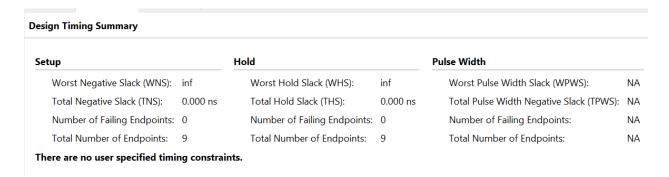
By making the design Moore, the output value synchronized with the clock which resulted an increasement in clock cycle of detecting the 1s and 0s. This is caused by the one clock cycle time requirement on the output due to the DFF. Instead of giving the output immediately after the 3rd clock cycle, while waiting for the clock for the output it receives the 4th state value correctly. As it can be seen on the simulation figure above, z only gives "1" when the input is either zero or one for four clock cycles instead of the three in the Mealy Machine. This time the design works as intended.



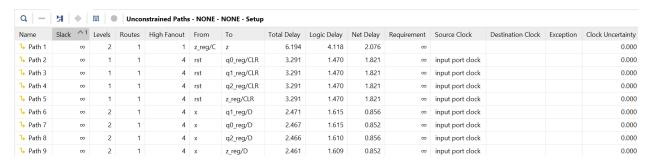
FSM Moore Post-Implementation Timing Simulation



FSM Moore Utilization Summary

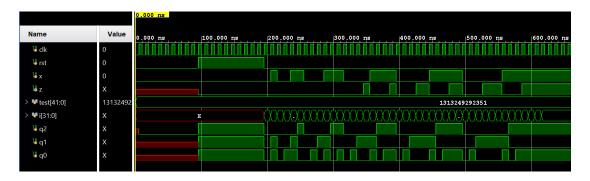


FSM Moore Design Timing Summary



FSM Moore Path Delays - Setup

STUCKING IN UNDESIREBLE STATES:



Behavioral Simulation for "111" state



Behavioral Simulation for "110" state

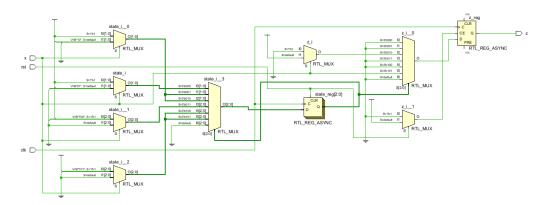
As it can be seen in the simulations initiating from the states "111" and "110" do not force device to be stuck in arbitrary states. However, the first z output value of the "110" state initiation is wrong.

BEHAVIORAL DESIGN:



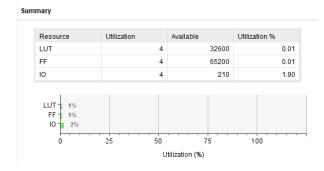
Behavioral FSM Behavioral Simulation

The simulations are the same for behavioral and Moore design. However, Moore design requires significantly more pre-computations. Behavioral design is better in terms of coding simplicity.

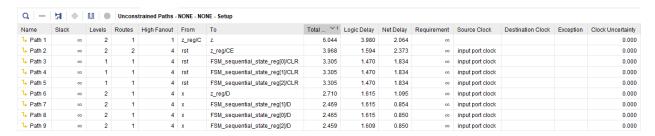


Behavioral FSM RTL Schematic

In terms of RTL placement, in behavioral design, instead of using numerous gates and paths, all the cases are represented with MUXs and registers. However, this does not mean that the technology schematic would decrease similarly.



Behavioral FSM Utilization Summary



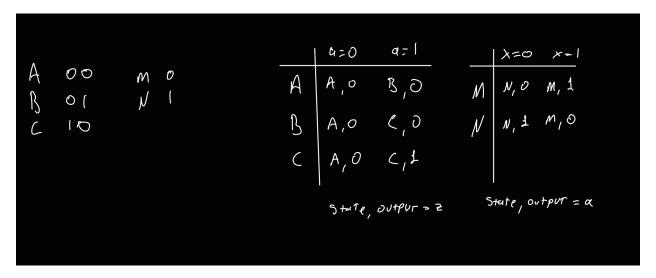
Behavioral FSM Path Delays - Setup

As it can be seen in the Utilization Summaries, the LUT usage is increase to 4 from 2 in the behavioral design. However, the longest path delay is decreased considerably. Beside the increase in "rst to z reg" path, all the other paths decreased a little.



Behavioral FSM Post-Implementation Timing Simulation

2. Design with Divided State Diagrams

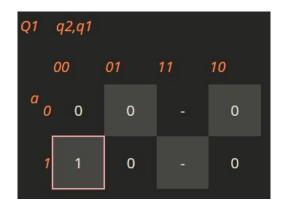


а	q2	q1	Q2	Q1	Z	Х	q0	Q0	а
0	0	0	0	0	0	0	0	1	0
0	0	1	0	0	0	0	1	1	1
0 0	1	0	0	0	0	1	0	0	1
0	1	1	Х	X	х	1	1	0	0
1	0	0	0	1	0				
1	0	1	1	0	0				
1	1	0	1	0	1				
1	1	1	Х	X	Х				

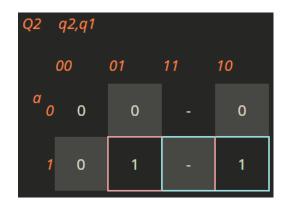
$$a = x \wedge q0$$
$$Q0 = ! x$$



$$z(a,q2,q1) = aq2$$



Q1(a,q2,q1) = aq2'q1'



$$Q2(a,q2,q1) = aq1 + aq2$$

REDUCED EQUATIONS

$$a = x ^ q0$$
 $Q0 = !x$
 $Q1 = a & (!q1) & (!q2)$
 $Q2 = (q1 & a) | (a & q2)$
 $Z = a & q2$

```
`timescale 1ns / 1ps
module FSM2 (
   input clk,
    input rst,
    input x,
    output reg z
    );
    reg q0, q1, q2;
    wire a, Z;
    wire Q0, Q1, Q2;
    assign a = x ^ q0;
    assign Q0 = !x;
    assign Q1 = a & (!q1) & (!q2);
    assign Q2 = (q1 \& a) | (a \& q2);
    assign Z = a & q2;
    always @(posedge clk or posedge rst)
begin
        if (rst) begin
            q2 <= 1'b0;
            q1 <= 1'b0;
            q0 <= 1'b0;
            z <= 1'b0;
        end else begin
            q2 <= Q2;
            q1 <= Q1;
            q0 <= Q0;
            z <= Z;
        end
    end
endmodule
```

FSM2 Verilog Code

```
`timescale 1ns / 1ps
module FSM2 tb();
    reg clk = 0;
        rst = 0;
    reg
        x = 0;
    reg
    wire z;
    FSM2 dUT(
       .clk(clk),
        .rst(rst),
        .x(x),
       .z(z)
    );
    always #5 clk = ~clk;
    reg [41:0] test =
42'b0100 1100 0111 0000 1111 0000 0111 1100 0000 1111 11;
    integer i;
    initial
    begin
        repeat(80) @(posedge clk);
        rst = 1;
        repeat(20) @(posedge clk);
        rst = 0;
        for(i=41; i>=0; i=i-1)
        begin
            x = test[i];
            @(posedge clk);
        end
    end
endmodule
```

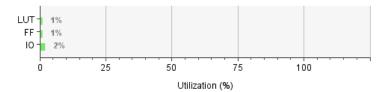
FSM2 Testbench Code



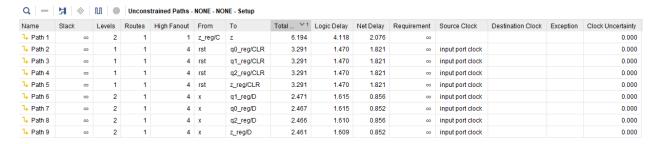
FSM2 Behavioral Simulation

Summary

Resource	Utilization	Available	Utilization %
LUT	2	32600	0.01
FF	4	65200	0.01
IO	4	210	1.90



Utilization Summary



FSM2 Path Delays

	Hold		Pulse Width	
inf	Worst Hold Slack (WHS):	inf	Worst Pulse Width Slack (WPWS):	NA
0.000 ns	Total Hold Slack (THS):	0.000 ns	Total Pulse Width Negative Slack (TPWS):	NA
0	Number of Failing Endpoints:	0	Number of Failing Endpoints:	NA
9	Total Number of Endpoints:	9	Total Number of Endpoints:	NA
	0.000 ns 0	inf Worst Hold Slack (WHS): 0.000 ns Total Hold Slack (THS): 0 Number of Failing Endpoints:	inf Worst Hold Slack (WHS): inf 0.000 ns Total Hold Slack (THS): 0.000 ns 0 Number of Failing Endpoints: 0	inf Worst Hold Slack (WHS): inf Worst Pulse Width Slack (WPWS): 0.000 ns Total Hold Slack (THS): 0.000 ns Total Pulse Width Negative Slack (TPWS): Number of Failing Endpoints: 0 Number of Failing Endpoints:

Timing Summary



FSM2 Post-Implementation Timing Simulation

FSM2 BEHAVIORAL:

```
`timescale Ins / lps
module FSM2 behav(
   input clk,
   input rst,
   input x,
   output reg z
   );
                          reg [1:0] state1;
reg state2;
reg a;
                     reg state2;
reg a;
always @ (posedge clk or posedge rst) begin
    if(rst) begin
    state1 <= A;
    z <= 1'b0;
end else begin
        state1 <= B;
        z <= 1'b0,
    end else begin
        state1 <= B;
        z <= 1'b0,
    end else begin
        state1 <= A;
        z <= 1'b0;
    end else begin
        state1 <= A;
        z <= 1'b0;
    end
    end
    B: begin
        if(a) begin
        state1 <= C;
        z <= 1'b0;
    end else begin
        state1 <= C;
        z <= 1'b0;
    end else begin
        state1 <= A;
        z <= 1'b0;
    end
    end
    c: begin
    if(a) begin
    state1 <= A;
    z <= 1'b1;
    end else begin
    state1 <= C;
    z <= 1'b1;
    end else begin
    state1 <= A;
    z <= 1'b0;
    end
    end
    default: begin
    state1 <= A;
    z <= 1'b0;
    end
    end
    end
end
end
end
end
end
end
always @ (posedge clk or posedge rst) begin
    if(rst) begin</pre>
```

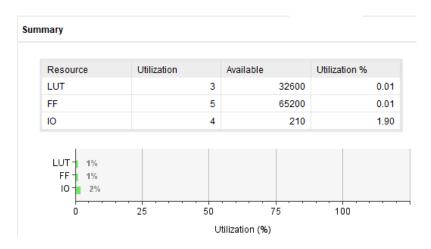
FSM2 Behavioral Verilog Code

```
`timescale 1ns / 1ps
module FSM2 behav tb();
          clk = 0;
    reg
          rst = 0;
    reg
             = 0;
          Х
    reg
    wire
         z;
    FSM2 behav DUT(
        .clk(clk),
        .rst(rst),
        .x(x),
        .z(z)
    );
    always #5 clk = ~clk;
    reg [41:0] test =
42'b0100 1100 0111 0000 1111 0000 0111 1100 0000 1111 11;
    integer i;
    initial
    begin
        repeat(80) @(posedge clk);
        rst = 1;
        repeat(20) @(posedge clk);
        rst = 0;
        for(i=41; i>=0; i=i-1)
        begin
            x = test[i];
            @(posedge clk);
        end
    end
endmodule
```

FSM2 Behavioral Testbench Code



FSM2 Behavioral - Behavioral Simulation



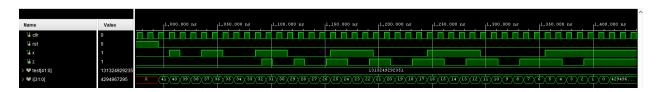
FSM2 Behavioral Utilization Summary



FSM2 Behavioral Timing Summary

Name	Slack	Levels	Routes	High Fanout	From	То	Total V 1	Logic Delay	Net Delay	Requirement	Source Clock	Destination Clock	Exception	Clock Uncertainty
3 Path 1	00	2	1	1	z_reg/C	Z	6.194	4.118	2.076	00				0.000
3 Path 2	00	1	1	5	rst	FSM_sequentialte1_reg[0]/CLR	3.291	1.470	1.821	∞	input port clock			0.000
3 Path 3	00	1	1	5	rst	FSM_sequentialte1_reg[1]/CLR	3.291	1.470	1.821	00	input port clock			0.000
3 Path 4	00	1	1	5	rst	a_reg/CLR	3.291	1.470	1.821	∞	input port clock			0.000
3 Path 5	00	1	1	5	rst	state2_reg/CLR	3.291	1.470	1.821	00	input port clock			0.000
3 Path 6	00	1	1	5	rst	z_reg/CLR	3.291	1.470	1.821	∞	input port clock			0.000
3 Path 7	00	2	1	2	x	state2_reg/D	2.844	1.644	1.200	00	input port clock			0.000
3 Path 8	00	2	1	2	х	a_reg/D	2.815	1.615	1.200	∞	input port clock			0.000
3 Path 9	00	2	1	3	a_reg/C	z_reg/D	1.486	0.608	0.878	00				0.000
3 Path 10	00	2	1	3	a_reg/C	FSM_sequential_state1_reg[1]/D	1.458	0.580	0.878	00				0.000

FSM2 behavioral Path Delays - Setup



FSM2 Behavioral Post-Implementation Timing Simulation

To compare two different models of FSM2, it can be seen on the utilization summaries that both the LUT and FF usage is increased by 1 for behavioral model. These additions are caused by the registers that contain states. However, besides the addition of the two new paths, the existing path delays have not changed at all.

To conclude, in FSM2 the behavioral model's coding and design is simpler but it takes more space than the first design. However, unlike FSM1, making the design behavioral does not affect the path delays of FSM2.

References

1- https://www.allaboutcircuits.com/technical-articles/encoding-the-states-of-a-finite-state-machine-vhdl/