

item and determines whether
instance is "yes" or "no"

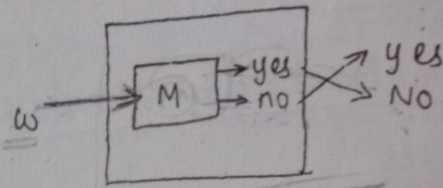
Closure Properties of recursive and recursively Enumerable languages

Theorem 1

The complement of a recursive language is recursive

Proof

- * Let L be a recursive language and M a Turing machine that halts on all inputs and accepts L . \downarrow M enters final state
- * Construct M' from M so that if M enters a final state on input w then M' halts without accepting.
- * If M halts without accepting, M' enters a final state.
- * Since one of these two events occurs, M' is an algorithm.
- * Clearly $L(M')$ is the complement of L and thus the complement of L is a recursive language.

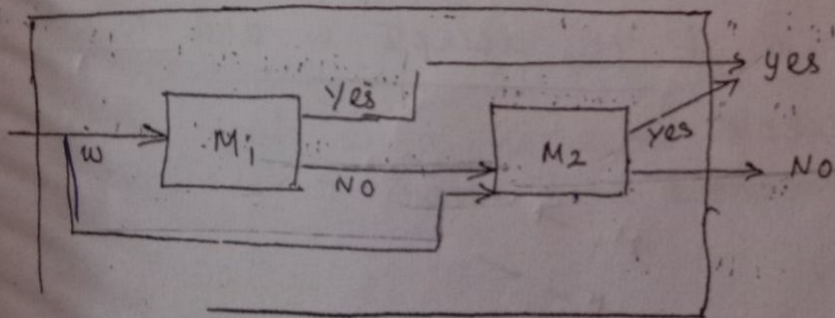


Theorem 2

- * The union of two recursive language is recursive.
- * The union of two recursively enumerable language is recursively enumerable.

Proof:

- * Let L_1 and L_2 be recursive languages accepted by algorithms M_1 and M_2 .
- * We construct M , which first simulates M_1 .
- * If M_1 accepts then M accepts.
- * If M_1 rejects then M simulates M_2 and accepts iff M_2 accepts.
- * Since both M_1 and M_2 are algorithms, M is guaranteed to halt.
- * Clearly M accepts $L_1 \cup L_2$.

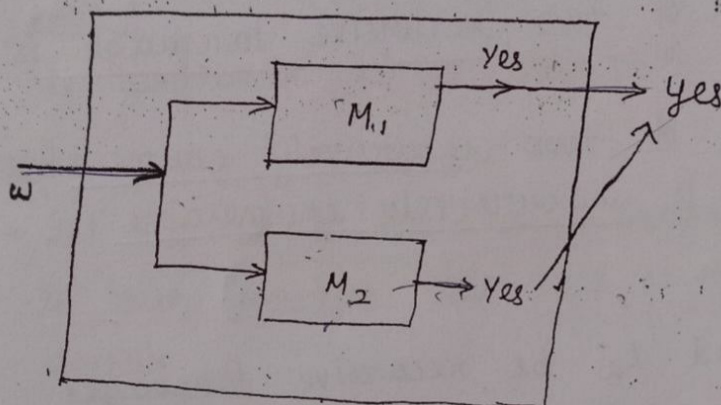


$M_1 \times M_2$

M_1 accept M_2 TM accept

... recursively enumerable ...
construction does not work.

- * Since M_1 may not halt.
- * Instead, M can simultaneously simulate M_1 and M_2 on separate tapes.
- * If either accepts, then M accepts.



Theorem 3:

If a language L and its complement \bar{L} are both recursively Enumerable, then L (and hence \bar{L}) is recursive.

Proof

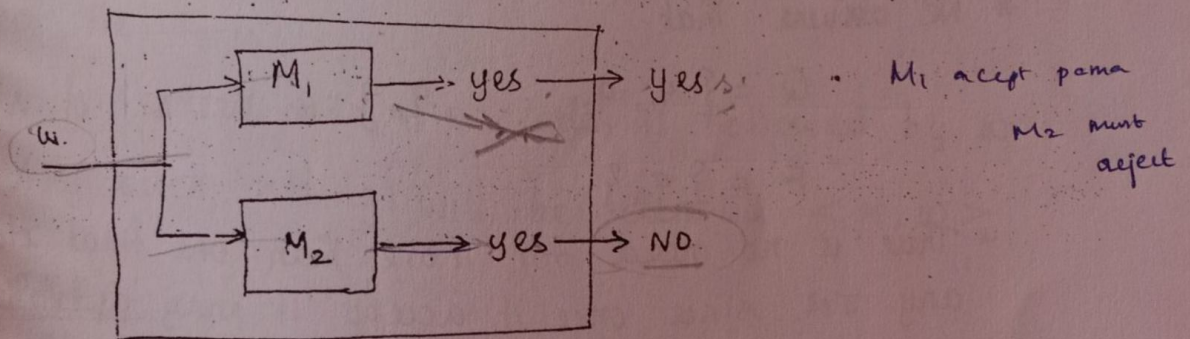
- * Let M_1 and M_2 accept L and \bar{L} respectively.
- * Construct M to simulate simultaneously M_1 and M_2 .
- * M accepts w if M_1 accepts w and rejects w if M_2 accepts w .

$M_1 \rightarrow w$
 $M_2 \rightarrow$

Since w is in either L or \bar{L} , we know that exactly one of M_1 or M_2 will accept.

* Thus M will always say either 'yes' or 'no' but will never say both.

* Since M is an algorithm that accepts L , it follows that L is recursive.



UNIVERSAL TURING MACHINE:

* Diagonalization is used to show the Problem