

- \* A problem is ~~an algorithm~~ algorithm that takes as input an instance of the problem and determines whether the answer to that instance is "yes" or "no"

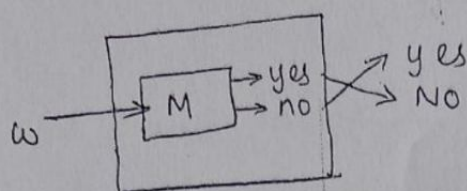
## Closure Properties of recursive and Recursively Enumerable languages.

### Theorem 1

The complement of a recursive language is recursive

### Proof

- \* Let  $L$  be a recursive language and  $M$  a Turing machine that halts on all inputs and accepts  $L$ . ✓
- \* Construct  $M'$  from  $M$  so that if  $M$  enters a final state on input  $w$  then  $M'$  halts without accepting.
- \* If  $M$  halts without accepting,  $M'$  enters a final state.
- \* Since one of these two events occurs,  $M'$  is an algorithm.
- \* Clearly  $L(M')$  is the complement of  $L$  and thus the complement of  $L$  is a recursive language.

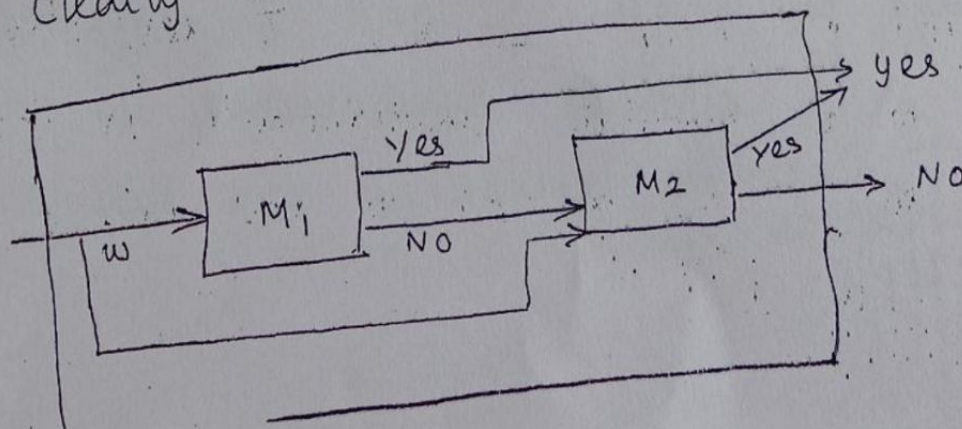


## Theorem 2

- \* The union of two recursive language is recursive.
- \* The union of two recursively enumerable language is recursively enumerable.

### Proof:

- \* Let  $L_1$  and  $L_2$  be recursive languages accepted by algorithms  $M_1$  and  $M_2$ .
- \* We construct  $M$ , which first simulates  $M_1$ .
- \* If  $M_1$  accepts then  $M$  accepts.
- \* If  $M_1$  rejects then  $M$  simulates  $M_2$  and accepts iff  $M_2$  accepts.
- \* Since both  $M_1$  and  $M_2$  are algorithms,  $M$  is guaranteed to halt.
- \* Clearly  $M$  accepts  $L_1 \cup L_2$ .



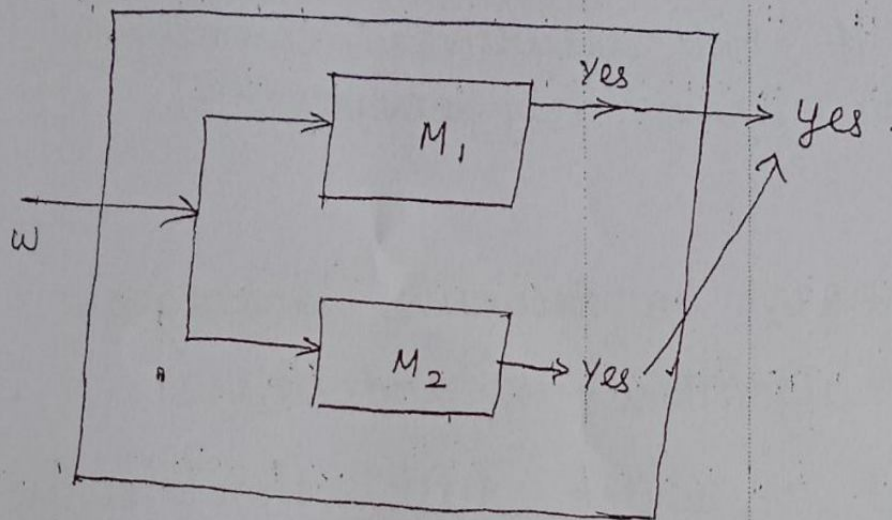


... recursively enumerable ... construction does not work.

\* Since  $M_1$  may not halt.

\* Instead,  $M$  can simultaneously simulate  $M_1$  and  $M_2$  on separate tapes.

\* If either accepts, then  $M$  accepts.



### Theorem 3:

If a language  $L$  and its complement  $\bar{L}$  are both recursively Enumerable, then  $L$  (and hence  $\bar{L}$ ) is recursive.

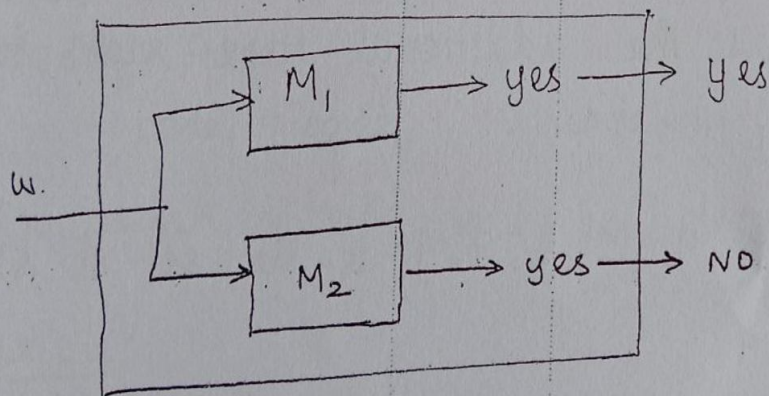
### Proof

- \* Let  $M_1$  and  $M_2$  accept  $L$  and  $\bar{L}$  respectively.
- \* Construct  $M$  to simulate simultaneously  $M_1$  and  $M_2$
- \*  $M$  accepts  $w$  if  $M_1$  accepts  $w$  and rejects  $w$  if  $M_2$  accepts  $w$

Since  $w$  is in either  $L$  or  $\bar{L}$ , we know that exactly one of  $M_1$  or  $M_2$  will accept.

\* Thus  $M$  will always say either 'yes' or 'no' but will never say both.

\* Since  $M$  is an algorithm that accepts  $L$ , it follows that  $L$  is recursive.



### UNIVERSAL TURING MACHINE:

\* Diagonalization is used to show the Problem