

### 3. MD5, SECURE HASH ALGORITHM

- ❖ Explain about MD5 in detail. (April/May'11, April/May'10 & May/June'12)
- ❖ Explain Secure Hashing Algorithm (SHA) (April/May'15, Nov/Dec'13, May/June'13 & April/May'10)
- ❖ Explain the process of deriving eighty 64-bit words from the 1024-bits for processing of a single block and also discuss single round function in SHA-512 algorithm. Show the results of  $W_{16}$ ,  $W_{17}$ ,  $W_{18}$  and  $W_{19}$ . (Nov/Dec'14)

#### a. MESSAGE DIGEST 5: MD5

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- ✓ Developed by Ron Rivest at MIT
- ✓ Input: a message of arbitrary length
- ✓ Output: 128-bit message digest
- ✓ 32-bit word units, 512-bit blocks

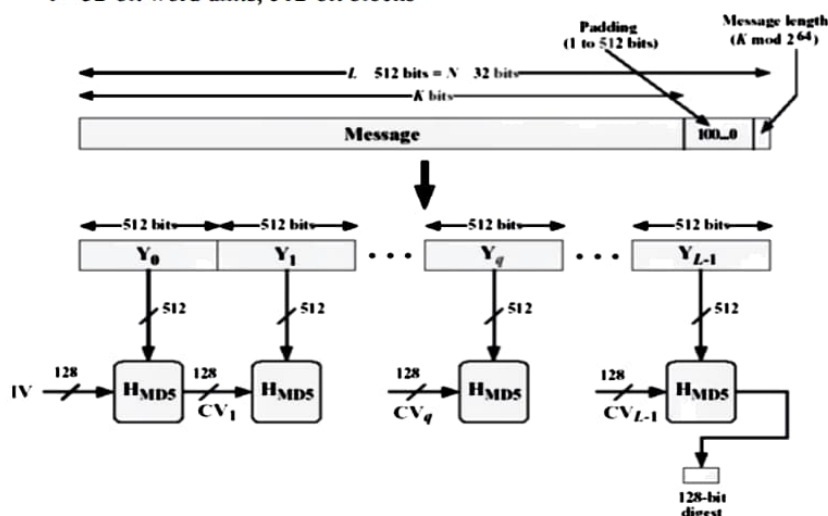


Figure: MD5 Logic

#### MD5 Logic:

##### Step 1: Append padding bits

- ✓ The message is Padded so that its bit length  $\equiv 448 \pmod{512}$  (i.e., the length of padded message is 64 bits less than an integer multiple of 512 bits)
- ✓ Padding is always added, even if the message is already of the desired length (1 to 512 bits)
- ✓ Padding bits: 1000...0 (a single 1-bit followed by the necessary number of 0-bits)

##### Step 2: Append length

- ✓ A 64-bit length: contains the length of the original message modulo 264
- ✓ The expanded message is  $Y_0, Y_1, \dots, Y_{L-1}$ ; the total length is  $L \times 512$  bits
- ✓ The expanded message can be thought of as a multiple of 16 32-bit words
- ✓ Let  $M[0 \dots N-1]$  denote the word of the resulting message, where  $N = L \times 16$

##### Step 3: Initialize MD buffer

- ✓ 128-bit buffer (four 32-bit registers A,B,C,D) is used to hold intermediate and final results of the hash function

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- ✓ A,B,C,D are initialized to the following values

A = 67452301

B = EFCDAB89

C = 98BADCFE

D = 10325476

- ✓ Stored in *little-endian* format (least significant byte of a word in the low-address byte position)
  - o E.g. word A : 01 23 45 67 (low address ... high address)
  - o word B : 89 AB CD EF
  - o word C : FE DC BA 98
  - o word D : 76 54 32 10

#### Step 4: Process message in 512-bit (16-word) blocks

- ✓ Heart of the algorithm called a *compression function* Consists of 4 rounds
- ✓ The 4 rounds have a similar structure, but each uses a different *primitive logical functions*, referred to as F, G, H, and I
- ✓ Each round takes as input the current 512-bit block ( $Y_q$ ), 128-bit buffer value ABCD and updates the contents of the buffer
- ✓ Each round also uses the table  $T[1 \dots 64]$ , constructed from the sine function;
  - $T[i] = 232 \times \text{abs}(\sin(i))$
- ✓ The output of 4th round is added to the  $CV_q$  to produce  $CV_{q+1}$

#### Step 5: Output

- ✓ After all  $L$  512-bit blocks have been processed, the output from the  $L$ th stage is the 128-bit message digest

$$CV_0 = IV$$

$$CV_{q+1} = \text{SUM32}(CV_q, RFI[Y_q, RFH[Y_q, RFG[Y_q, RFF[Y_q, CV_q]]])$$

$$MD = CV_L$$

Where  $IV$  = initial value of the ABCD buffer, defined in step 3

$Y_q$  = the  $q$ th 512-bit block of the message

$L$  = the number of blocks in the message (including padding and length fields)  $CV_q$  = chaining variable processed with the  $q$ th block of the message

$RF_x$  = round function using primitive logical function  $x$

$MD$  = final message digest value

$\text{SUM32}$  = addition modulo 232 performed separately on each word

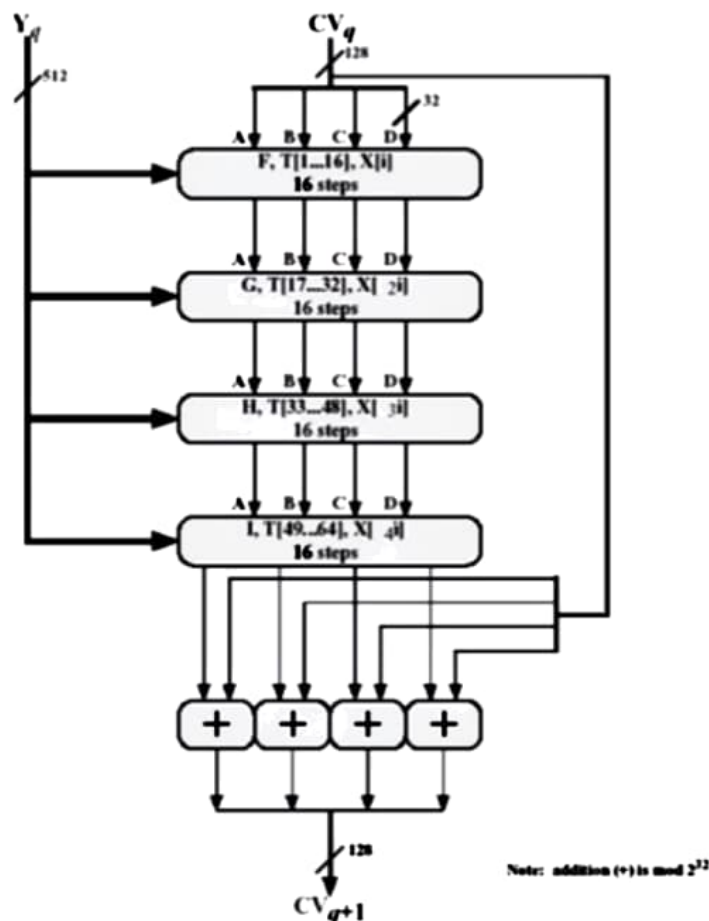


Figure: MD5 processing of a single 512-bit block (MD5 compression function)

### MD5 Compression Function:

- ✓ Each round consists of a sequence of 16 steps operating on the buffer ABCD
- ✓ Each step is of the form

$$a \leftarrow b + ((a + g(b, c, d) + X[k] + T[i] \lll s))$$

where

$a, b, c, d$  = the 4 words of the buffer, in a specified order that varies across steps

$g$  = one of the primitive functions  $F, G, H, I$

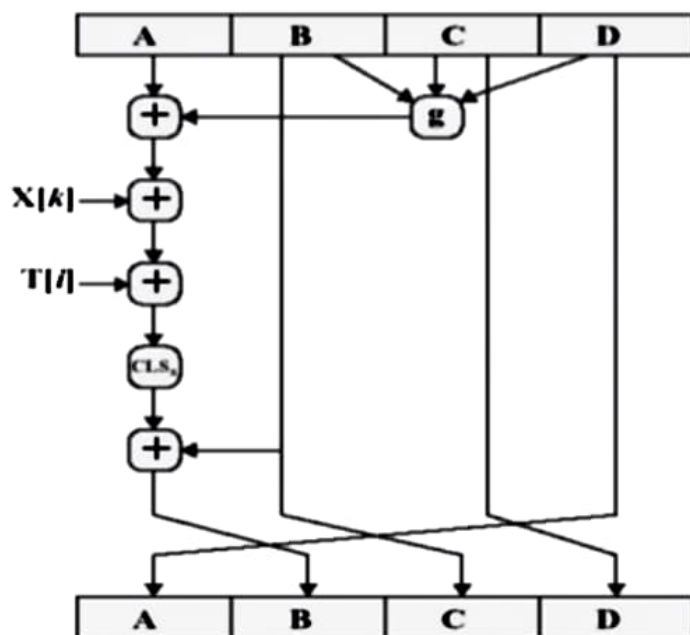
$\lll s$  = circular left shift (rotation) of the 32-bit arguments by  $s$  bits

$X[k] = M[q \times 16 + k]$  = the  $k$ th 32-bit word in the  $q$ th 512-bit block of the message

$T[i]$  = the  $i$ th 32-bit word in table  $T$

$+$  = addition modulo  $2^{32}$

## MD5 Operation



- ✓ One of the 4 primitive logical functions is used in each 4 rounds of the algorithm
- ✓ Each primitive function takes three 32-bit words as input and produces a 32-bit word output
- ✓ Each function performs a set of bitwise logical operations

Round	Primitive function $g$	$g(b, c, d)$
1	$F(b, c, d)$	$(b \wedge c) \vee (b' \wedge d)$
2	$G(b, c, d)$	$(b \wedge d) \vee (c \wedge d')$
3	$H(b, c, d)$	$b \oplus c \oplus d$
4	$I(b, c, d)$	$c \oplus (b \vee d')$

TRUTH TABLE						
b	C	d	F	G	H	I
0	0	0	0	0	0	1
0	0	1	1	0	1	0
0	1	0	0	1	1	0
0	1	1	1	0	0	1
1	0	0	0	0	1	1
1	0	1	0	1	0	1

1	1	0	1	1	0	0
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- ✓ The array of 32-bit words  $X[0..15]$  holds the value of current 512-bit input block being processed
- ✓ Within a round, each of the 16 words of  $X[i]$  is used exactly once, during one step
  - o The order in which these words is used varies from round to round
  - o In the first round, the words are used in their original order
  - o For rounds 2 through 4, the following permutations are used
    - »  $\rho_2(i) = (1 + 5i) \bmod 16$
    - »  $\rho_3(i) = (5 + 3i) \bmod 16$
    - »  $\rho_4(i) = 7i \bmod 16$