R2019 Question B...

- ❖ Explain about MD5 in detail. (April/May'11, April/May'10 & May/June'12)
- Explain Secure Hashing Algorithm (SHA) (April/May'15, Nov/Dec'13, May/June'13 & April/May'10)

- Explain the process of deriving eighty 64-bit words from the 1024-bits for processing of a single block and also discuss single round function in SHA-512 algorithm. Show the results of W16, W17, W18 and W19. (Nov/Dec'14)
- a. MESSAGE DIGEST 5: MD5
- Developed by Ron Rivest at MIT
- Input: a message of arbitrary length
- Output: 128-bit message digest
- 32-bit word units, 512-bit blocks

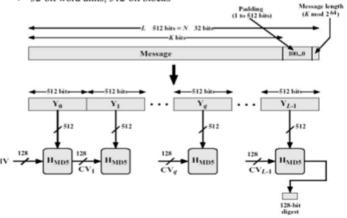


Figure: MD5 Logic

MD5 Logic:

Step 1: Append padding bits

- ✓ The message is Padded so that its bit length

 = 448 mod 512 (i.e., the length) of padded message is 64 bits less than an integer multiple of 512 bits)
- Padding is always added, even if the message is already of the desired length (1 to 512 bits)
- ✓ Padding bits: 1000....0 (a single 1-bit followed by the necessary number of 0-

Step 2: Append length

- ✓ A 64-bit length: contains the length of the original message modulo 264
- ✓ The expanded message is Y₀, Y₁, ..., Y_{L-1}; the total length is L × 512 bits
 ✓ The expanded message can be thought of as a multiple of 16 32-bit words
- ✓ Let M[0 ... N-1] denote the word of the resulting message, where $N = L \times 16$

Step 3: Initialize MD buffer

- ✓ 128-bit buffer (four 32-bit registers A,B,C,D) is used to hold intermediate and final results of the hash function
- ✓ A,B,C,D are initialized to the following values

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A = 67452301

B = EFCDAB89

C = 98BADCFE

D = 10325476

- ✓ Stored in little-endian format (least significant byte of a word in the lowaddress byte position)
 - o E.g. word A: 01 23 45 67 (low address ... high address)
 - o word B: 89 AB CD EF
 - o word C : FE DC BA 98
 - o word D: 76 54 32 10

Step 4: Process message in 512-bit (16-word) blocks

- ✓ Heart of the algorithm called a compression function Consists of 4 rounds
- ✓ The 4 rounds have a similar structure, but each uses a different primitive logical functions, referred to as F, G, H, and I
- Each round takes as input the current 512-bit block (Yq), 128-bit buffer value ABCD and updates the contents of the buffer
- ✓ Each round also uses the table T[1 ... 64], constructed from the sine function;
 - $T[i] = 232 \times abs(sin(i))$
- ✓ The output of 4th round is added to the CVq to produce CVq+1

Step 5: Output

 After all L 512-bit blocks have been processed, the output from the Lth stage is the 128-bit message digest

 $CV_0 = IV$

CVq+1 = SUM32(CVq, RFI[Yq, RFH[Yq, RFG[Yq, RFF[Yq,

CVq]]])

 $MD = CV_L$

Where IV = initial value of the ABCD buffer, defined in step 3

Yq = the qth 512-bit block of the message

L = the number of blocks in the message (including padding and length fields) CVq = chaining variable processed with the qth block of the message

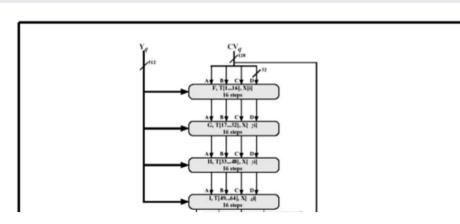
RFx = round function using primitive logical function

MD = final message digest value

SUM32 = addition modulo 232 performed separate

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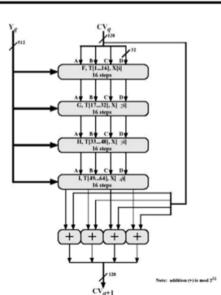


Figure: MD5 processing of a single 512-bit block (MD5 compression function)

MD5 Compression Function:

- ✓ Each round consists of a sequence of 16 steps operating on the buffer ABCD
- ✓ Each step is of the form

$$a \leftarrow b + ((a + g(b, c, d) + X[k] + T[i] <<< s)$$

where

 $a,b,c,d = the \ 4$ words of the buffer, in a specified order that varies across steps g = one of the primitive functions F, G, H, I

<<s = circular left shift (rotation) of the 32-bit arguments by s bits

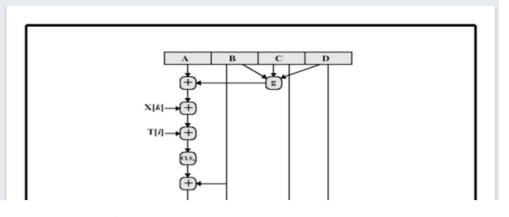
 $X[k] = M[q \times 16 + k]$ = the kth 32-bit word in the qth 512-bit block of the message

T[i] = the ith 32-bit word in table T

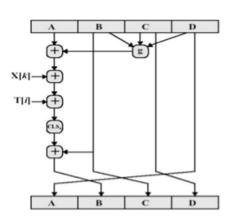
+ = addition modulo 232

MD5 Operation

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- ✓ One of the 4 primitive logical functions is used in each 4 rounds of the algorithm
- Each primitive function takes three 32-bit words as input and produces a 32-bit word output
- ✓ Each function performs a set of bitwise logical operations

Round	Primitive function g	g(b, c, d)		
1	F(b, c, d)	(b ∧ c) ∨ (b' ∧ d)		
2	G(b, c, d)	(b ∧ d) ∨ (c ∧ d')		
3	H(b, c, d)	b ⊕ c ⊕ d		
4	I(b c, d)	c ft (h V d')		

TRUTH TABLE							
b	C	d	F	G	Н	I	
0	0	0	0	0	0	1	
0	0	1	1	0	1	0	
0	1	0	0	1	1	0	
0	1	1	1	0	0	1	
1	0	0	0	0	1	1	
1	0	1	0	1	0	1	
1	1	0	1	1	0	0	

- ✓ The array of 32-bit words X[0..15] holds the value of current 512-bit input block being processed
- Within a round, each of the 16 words of X[i] is used exactly once, during one
 - · The order in which these words is used varies from round to round
 - o In the first round, the words are used in their original order
 - \circ $\,$ For rounds 2 through 4, the following permutations are used

- » $\rho 2(i) = (1 + 5i) \mod 16$ » $\rho 3(i) = (5 + 3i) \mod 16$
- $\rho 4(I) = 7i \mod 16$

b. SECURE HASH ALGORITHM

- Developed by NIST (National Institute of Standards and Technology)
 - Published as a FIPS 180 in 1993
 - A revised version is issued as FIPS 180-1 IN 1995
 - Generally referred to as SHA-1
- SHA is based on the hash function MD4 and its design closely models MD4.
- SHA- 1 produces a hash value of 160 bits.
- Revised version of the standard, FIPS 180-2, that defined three new versions of SHA, with hash value lengths of 256, 384 and 512 bits, known as SHA-256,





