

3. OPTIMIZATION OF BASIC BLOCKS

- ❖ Illustrate the optimization of basic blocks with an example. [Nov/Dec 2014]
- ❖ Discuss in detail the process of optimization of basic blocks. Give an example. [May/Jun 2014]
- ❖ Explain in detail the optimization of basic blocks. [Nov/Dec 2011]

OPTIMIZATION OF BASIC BLOCKS

There are two types of basic block optimizations. They are :

1. Structure -Preserving Transformations
2. Algebraic Transformations

1. Structure- Preserving Transformations:

The primary Structure-Preserving Transformation on basic blocks are:

- Common sub-expression elimination
- Dead code elimination
- Renaming of temporary variables
- Interchange of two independent adjacent statements.

Common sub-expression elimination:

Common sub expressions need not be computed over and over again. Instead they can be computed once and kept in store from where it's referenced when encountered again – of course providing the variable values in the expression still remain constant.

Example:

```
a=b+c  
b=a-d  
c=b+c  
d=a-d
```

The 2nd and 4th statements compute the same expression: b+c and a-d

Basic block can be transformed to

```
a=b+c  
d=a-d  
c=d+c
```

Dead code elimination:

It's possible that a large amount of dead (useless) code may exist in the program. This might be especially caused when introducing variables and procedures as part of construction or error -correction of a program – once declared and defined, one forgets to remove them in case they serve no purpose. Eliminating these will definitely optimize the code.

Renaming of temporary variables:

A statement $t := b + c$ where t is a temporary name can be changed to $u := b + c$ where u is another temporary name, and change all uses of t to u .
In this we can transform a basic block to its equivalent block called normal-form block.

Interchange of two independent adjacent statements:

Two statements

$t1 := b + c$

$t2 := x + y$

can be interchanged or reordered in its computation in the basic block when value of $t1$ does not affect the value of $t2$.

2. Algebraic Transformations:

- Algebraic identities represent another important class of optimizations on basic blocks. This includes simplifying expressions or replacing expensive operation by cheaper ones i.e. reduction in strength.
- Another class of related optimizations is constant folding. Here we evaluate constant expressions at compile time and replace the constant expressions by their values. Thus the expression $2 * 3.14$ would be replaced by 6.28 .
- The relational operators $<=$, $>=$, $<$, $>$, $+$ and $=$ sometimes generate unexpected common sub expressions.
- Associative laws may also be applied to expose common sub expressions. For example, if the source code has the assignments

$a := b + c$

$e := c + d + b$

the following intermediate code may be generated:

$a := b + c$

$t := c + d$

$e := t + b$

- Example:
 - $x := x + 0$ can be removed
 - $x := y ** 2$ can be replaced by a cheaper statement $x := y * y$
- The compiler writer should examine the language carefully to determine what rearrangements of computations are permitted, since computer arithmetic does not always obey the algebraic identities of mathematics. Thus, a compiler may evaluate $x * y - x * z$ as $x * (y - z)$ but it may not evaluate $a + (b - c)$ as $(a + b) - c$.