

Session 12. Advanced Decision Models

* Game Theory

- **Game theory** attempts to mathematically capture behavior in *strategic situations*, in which an individual's success in making choices depends on the choices of others.

Ex] Soccer penalty kicks

P[Scoring a goal]		Goalie	
		b_1 : Dive left	b_2 : Dive right
Right-footed Kicker	a_1 : Kick left	0.637	0.944
	a_2 : Kick right	0.893	0.437

- The *optimal* strategy for the kicker?
- The *optimal* strategy for the goalie?
- The *value* of the game?

* Multi-Objective Decision Making

- Many real-world decision problems have **multiple objectives**. For example, when choosing a medical treatment plan, we want to maximize the **efficacy** of the treatment, but also minimize the **side effects**.
- These objectives typically **conflict**, e.g., we can often increase the **efficacy** of the treatment, but at the cost of more severe **side effects**.
- Multi-objective methods
 - Analytical hierarchical process (**AHP**)
 - Multi-objective **linear programming** (**MOLP**)
 - Multi-attribute **utility** theory



A. Game Theory*

* Game Theory

- Two or more decision-makers have *conflicting* interests.
- Game theory attempts to mathematically capture behavior in *strategic situations*, in which an individual's success in making choices depends on the choices of others.
- Game theory has been widely recognized as an important tool in many fields. Eight *game theorists* have won *Nobel prizes* in economics.

* Basic Assumptions

1. *Both* players are *rational*.
2. *Both* players choose their strategies solely to *promote their own welfare* (no compassion for the opponent).

* Payoff Table

- It shows the *gains* (or payoffs) for *Player A* that would result from each combination of strategies for the two players.

Payoff Table		Player B			
		b_1	b_2	...	b_n
Player A	a_1	r_{11}	r_{12}	...	r_{1n}
	a_2	r_{21}	r_{22}	...	r_{2n}

	a_m	r_{m1}	r_{m2}	...	r_{mn}

* Classifications

- *Two person* game
 - Zero-sum game
 - Constant-sum game
 - Non-constant-sum game
- *n-person* game



* Solving Two-Person Zero-Sum Games

- **Step 1.** Games with **dominated** strategies:

Eliminate any of Player A's **dominated** strategies. Looking at the *reduced* payoff table, eliminate any of Player B's **dominated** strategies. Now eliminate any of Player A's dominated strategies. Continue in this fashion until no more dominated strategies can be found.

- **Step 2.** Games with **pure** strategies:

Find the **maximin** strategy for Player A that maximizes the minimum payoff. Find the **minimax** strategy for Player B that minimizes the maximum loss. If the maximin value equals to the minimax value, then we find the **stable** solution (*saddle point*) for both players. Players A and B should exclusively use their maximin and minimax strategies, respectively.

- **Step 3.** Games with **mixed** strategies:

Introduce variables that represent the probabilities of selecting each strategy for the players. Find the **maximin** mixed strategy for Player A that maximizes the *expected* minimum payoff. For Player B, find the **minimax** mixed strategy that minimizes the *expected* maximum loss.

* Solution Methods for **Mixed** Strategies



- Graphical solution (or **calculus**):

If one of the players has only **two** strategies, we can find the optimal mixed strategies from a 2-dimensional plane.

- **Linear programming (LP)** formulation

Use Microsoft Excel to solve the LP model.

Ex 1] Games with Dominated Strategies

The head football coach of LSU is attempting to come up with a strategy to deal with Alabama. **LSU** is on offense, and **Alabama** is on defense. The LSU coach has **5** preferred plays, but is not sure which to select. He knows, however, that Alabama usually employs one of **3** defensive strategies. Over the years, he has diligently recorded the average yardage gained by his team for each combination of strategies used:

Yardage gained		Alabama		
		b_1	b_2	b_3
LSU	a_1	0	-1	5
	a_2	7	5	10
	a_3	15	-4	-5
	a_4	5	0	10
	a_5	-5	-10	10

Which of the five plays should the LSU coach select?

- Step 1:
- Step 2:
- Step 3: |
- Step 4:
- The optimal strategies are



Ex 2] Games with Pure Strategies: Players A and B simultaneously call out one of the numbers, one or two.

		Player B		MAX
		b_1 : One	b_2 : Two	min
Player A	a_1 : One	+\$3	+\$2	
	a_2 : Two	+\$1	-\$6	
MIN	max			

- The *optimal* strategies =

Ex 3] Games with Mixed Strategies

Suppose that **Player A** chooses a_1 with probability x and chooses a_2 with probability $1-x$, and **Player B** chooses b_1 with probability y and chooses b_2 with probability $1-y$.

Player A's Payoff v_{ij}		Player B $b_1 \quad b_2$		P[a_i]	Joint Probabilities		Player B $b_1 \quad b_2$		P[a_i]
Player A	a_1	2	4	x	Player A	a_1	xy	$x(1-y)$	x
	a_2	3	1	$1-x$		a_2	$(1-x)y$	$(1-x)(1-y)$	$1-x$
P[b_j]		y	$1-y$	1.0	P[b_j]		y	$1-y$	1.0

- The **expected value** v of the game to Player A is

$$v = \sum_{i=1}^m \sum_{j=1}^n x_i r_{ij} y_j =$$



- Player A's** **maximin** strategy, $(x^*, 1-x^*)$

Find the first-order derivative with respect to y and set it equal to 0. Solve the equation for x .

$$\frac{d}{dy} v = \frac{d}{dy} (-4xy + 3x + 2y + 1) =$$

Thus, the optimal strategy is $x^* =$ and $1-x^* =$

- Player B's** **minimax** strategy, $(y^*, 1-y^*)$

Find the first-order derivative with respect to x and set it equal to 0. Solve the equation for y .

$$\frac{d}{dx} v = \frac{d}{dx} (-4xy + 3x + 2y + 1) =$$

Thus, the optimal strategy is $y^* =$ and $1-y^* =$

- Value** of the game, v^*

$$v^* = -4 x^* y^* + 3 x^* + 2 y^* + 1 =$$

Ex 4] Pure strategies: Consider the following payoff table, which represents player A's gain. Is this a fair game?

Payoff Table for A		Player B				Maxi min
		b_1	b_2	b_3	b_4	
Player A	a_1	3	2	4	2	
	a_2	6	-4	-8	-3	
	a_3	4	2	3	2	
	a_4	-5	-3	7	-4	
Minimax						

- Dominated strategies?
- Optimal strategies?
- Value of the game?

Ex 5] Mixed Strategies: Suppose that **Player A** chooses a_1 with probability x and chooses a_2 with probability $1-x$, and **Player B** chooses b_1 with probability y and chooses b_2 with probability $1-y$.

Player A's Payoff r_{ij}'		Player B		P[a_i]	Joint Probabilities		Player B		P[a_i]
		b_1	b_2				b_1	b_2	
Player A	a_1	9	13	x	Player A	a_1	xy	$x(1-y)$	x
	a_2	11	7	$1-x$		a_2	$(1-x)y$	$(1-x)(1-y)$	$1-x$
P[b_j]		y	$1-y$	1.0	P[b_j]		y	$1-y$	1.0

- Player A's *maximin* strategy, $(x^*, 1-x^*)$
The optimal strategy is $x^* =$ and $1-x^* =$
- Player B's *minimax* strategy, $(y^*, 1-y^*)$
The optimal strategy is $y^* =$ and $1-y^* =$
- Value of the game, v^*

$$v^* = 6x^* + 4y^* - 8x^*y^* + 7 =$$

Invariance under the change of *location* and *scale*: $r_{ij}' = 2r_{ij} + 5$

B. LP Formulation for Mixed Strategies*

Any game with **mixed** strategies can be solved by transforming the problem to a **linear programming** problem.

Payoff Table		Player B				P[a_i]	w_i
		b_1	b_2	...	b_n		
Player A	a_1	r_{11}	r_{12}	...	r_{1n}	x_1	w_1
	a_2	r_{21}	r_{22}	...	r_{2n}	x_2	w_2

	a_m	r_{m1}	r_{m2}	...	r_{mn}	x_m	w_m
P[b_j]		y_1	y_2	...	y_n	1.0	$\max\{w_j\}$
v_j		v_1	v_2	...	v_n	$\min\{v_i\}$	v^*

- (a) **Player A:** Let v be the Player A's *minimum* expected **gain**. Since Player A is trying to maximize the minimum expected gain v , the LP model is

$$\text{Max } z = v$$

$$\sum_{i=1}^m x_i r_{ij} \geq v \quad \text{for } j = 1, 2, \dots, n \text{ and}$$

$$\sum_{i=1}^m x_i = 1 \quad \text{where } x_i \geq 0 \text{ and } v \text{ is unrestricted in sign.}$$

- (b) **Player B:** Let w be the Player B's *maximum* expected **loss**. Since Player B is trying to minimize the maximum expected profit w , the LP model is

$$\text{Min } z = w$$

$$\sum_{j=1}^n r_{ij} y_j \leq w \quad \text{for } i = 1, 2, \dots, m \text{ and}$$

$$\sum_{j=1}^n y_j = 1 \quad \text{where } y_j \geq 0 \text{ and } w \text{ is unrestricted in sign.}$$



(c) **Combined LP formulation**

$$\min_{1 \leq j \leq n} \sum_{i=1}^m x_i r_{ij} = v \leq w = \max_{1 \leq i \leq m} \sum_{j=1}^n r_{ij} y_j$$

Player A's *maximin* value \leq Player B's *minimax* value

- It can be shown from the **dual theorem** that the *optimal* objective function values v and w are equal.
- Thus, player A's *floor* (maximin value) equals player B's *ceiling* (minimax value).
- This result is often known as the *minimax theorem*.

- **Objective function:**

$$\text{Min } z = w - v$$

- **Constraints**

$$\sum_{i=1}^m x_i r_{ij} \geq v \quad \text{for } j = 1, 2, \dots, n$$

$$\sum_{j=1}^n r_{ij} y_j \leq w \quad \text{for } i = 1, 2, \dots, m$$

$$\sum_{i=1}^m x_i = 1$$

$$\sum_{j=1}^n y_j = 1$$

where $0 \leq x_i < 1$ and $0 \leq y_j \leq 1$ and

v and w are *unrestricted* in sign.



Ex 1] Penalty Kicks in Soccer

Suppose that the **kicker** can only shoot either **right** or **left**. Similarly, the **goalie** must jump either **right** or **left**. (For simplicity, assume that kicking the ball in the **center** or staying in the **center** is not a viable option for either the **kicker** or the **goalie**.) A **kicker** has a higher probability of scoring on his natural side (i.e. on the left side for a right-footed kicker) and the **goalie** has a higher probability of saving the goal if he guesses the correct side.

Right-footed kicker's payoff		Goalie b_1 : Left b_2 Right		Kicker's strategy	Goalie's loss
Kicker	a_1 : Left	63.7%	94.4%	x_1	$\$w_1$
	a_2 : Right	89.3%	43.7%	x_2	$\$w_2$
Goalie's strategy		y_1	y_2	Sum = 1	$w = \max\{w_1, w_2\}$
Kicker's gain		$\$v_1$	$\$v_2$	$v = \min\{v_1, v_2\}$	Min $z = w - v$

▪ LP model

- Objective function

$$\text{Minimize } z = w - v$$



- Constraints

$$0.637 x_1 + 0.893 x_2 = v_1 \qquad 0.637 y_1 + 0.944 y_2 = w_1$$

$$0.944 x_1 + 0.437 x_2 = v_2 \qquad 0.893 y_1 + 0.437 y_2 = w_2$$

$$v_1 \geq v$$

$$w_1 \leq w$$

$$v_2 \geq v$$

$$w_2 \leq w$$

$$x_1 + x_2 = 1$$

$$y_1 + y_2 = 1$$

$$0 \leq x_1 \text{ and } x_2 \leq 1$$

$$0 \leq y_1 \text{ and } y_2 \leq 1$$

v and w are *unrestricted* in sign.

- Decision variables

$$(x_1, x_2, y_1, y_2) \text{ and } (v, w).$$

▪ Various strategies

Right-footed kicker's payoff		Goalie	
		<i>Randomly</i>	<i>Strategically</i>
Kicker	<i>Randomly</i>	Case 1	Case 2
	<i>Strategically</i>	Case 3	Case 4

- Case #1

- Mixed strategy: $\mathbf{x} = (0.5, 0.5)$ and $\mathbf{y} = (0.5, 0.5)$.
- The kicker's average success rate is **72.8%**.



- Case #2

- Mixed strategy: $\mathbf{x} = (0.5, 0.5)$ and $\mathbf{y} = (0.664, 0.336)$.
- The kicker's average success rate is **69.1%**.

- Case #3

- Mixed strategy: $\mathbf{x} = (0.598, 0.402)$ and $\mathbf{y} = (0.5, 0.5)$.
- The kicker's average success rate is **79.1%**.

- Case #4


- Mixed strategy: $\mathbf{x} = (0.598, 0.402)$ and $\mathbf{y} = (0.664, 0.336)$.
- The kicker's average success rate is **74.0%**.

▪ Kicker's success rate for each case

Right-footed kicker's payoff		Goalie	
		<i>Randomly</i>	<i>Strategically</i>
Kicker	<i>Randomly</i>	72.8%	69.1%
	<i>Strategically</i>	79.1%	74.0%

Their **mixed** strategies help them maximize their **expected equilibriums**, thereby providing them with the best response for each other's actions.

Ex 2] Search Game

Dr. Chun writes down one of the numbers (**1**, **2**, **3**) and you must repeatedly **guess** this number until you get it, losing \$1 for each **wrong** guess. After each guess, Dr. Chun will say  whether your guess is too **high**, too **low**, or correct.

- **Game theory**
 - Dr. Chun's strategy a_i : How to pick his **number i** .
 - Your strategy b_j : How to sequence your **guesses**
- **Case #1.** Dr. Chun **randomly** picks the number.
 - **Pure** strategy: you always choose the strategy, (2, 1 or 3).
 - The value of the game is **\$0.67**.

		Your strategy					$P[a_i]$
		1,2,3	1,3,2	2,1or3	3,1,2	3,2,1	
His strategy	1	0	0	1	1	2	1/3
	2	1	2	0	2	1	1/3
	3	2	1	1	0	0	1/3
$E[\text{Payoff}]$		\$1	\$1	\$0.67	\$1	\$1	\$0.67

- **Case #2.** Dr. Chun **strategically** picks the number.
 - **Mixed** strategies
 - The value of the game is **\$0.80**.

		Your strategy					$P[a_i]$
		1,2,3	1,3,2	2,1or3	3,1,2	3,2,1	
His strategy	1	0	0	1	1	2	0.4
	2	1	2	0	2	1	0.2
	3	2	1	1	0	0	0.4
$P[b_j]$		0	0.2	0.6	0.2	0	\$0.80

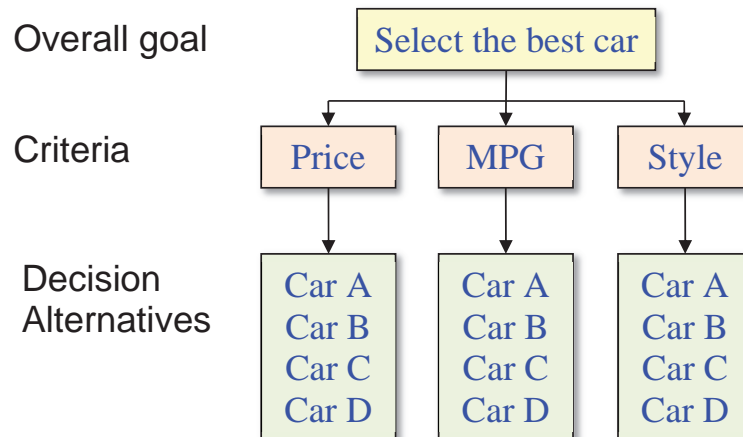
C. Analytic Hierarchy Process*

* Analytic Hierarchy Process

- The **Analytic Hierarchy Process (AHP)** is a structured technique for organizing and analyzing complex decisions.
- Based on mathematics and psychology, it was developed by **Thomas L. Saaty** in the 1970s and has been extensively studied and refined since then.
- Users of the AHP first decompose their decision problem into a **hierarchy** of more easily comprehended sub-problems, each of which can be analyzed independently.

* Developing the Hierarchy

1. Overall goal
2. Criteria
3. Decision alternatives



* Applications

Choice, ranking, prioritization, resource allocation, benchmarking, quality management, conflict resolution, etc.

* Procedures

- **Step 1. Criteria:** The decision maker specifies judgments about the relative importance of each **criterion** in terms of its contribution to the achievement of the **overall goal**.
- **Step 2. Alternatives:** The decision maker indicates a preference or priority for each **decision alternatives** in terms of how it contributes to each **criterion**.
- **Step 3. Overall preference:** The output of AHP is a prioritized ranking indicating the **overall preference** for each **decision alternative**.



* Measurement Scale

Criteria →	Score	← Alternatives
Absolutely more important	9	Extremely preferred
Very strongly more important	7	Very strongly preferred
Strongly more important	5	Strongly preferred
Weakly more important	3	Moderately preferred
Of equal importance	1	Equally preferred

* Criticisms

- While the general consensus is that it is both technically valid and practically useful, the **AHP** does have its critics. Most of the criticisms involve a phenomenon called **rank reversal**.
- Decision-making involves ranking alternatives in terms of criteria or attributes of those alternatives. It is an **axiom** of some decision theories that when new alternatives are added to a decision problem, the ranking of the old alternatives must not change — that "**rank reversal**" must not occur.

Ex] Soulmate: Kevin is ready to select his mate for life and has determined that **beauty**, **intelligence**, and **personality** are the **key factors** in selecting a mate.



- **Step 1. Criteria:** Obtain the weight for each criterion.

<i>Criteria</i>	Beauty	Intelligence	Personality
Beauty	1		
Intelligence	4	1	
Personality	8	4	1
Total			

<i>Criteria</i>	Beauty	Intelligence	Personality	Total	Weight
Beauty	0.077	0.048	0.091	0.216	
Intelligence				0.680	
Personality	0.615	0.762	0.727	2.104	
Total	1.000	1.000	1.000	3.000	1.000

- **Step 2. Alternatives:** Three girlfriends (Mary, Melanie, and Molly) are begging to be the Kevin's mate.

(a) **Beauty**

	Mary	Melanie	Molly
Mary		5	3
Melanie			
Molly		2	
Total			

	Mary	Melanie	Molly	Total	<i>Score</i>
Mary	0.652	0.625	0.667	1.944	
Melanie				0.366	
Molly	0.217	0.250	0.222	0.689	
Total	1.000	1.000	1.000	3.000	1.000

(b) *Intelligence*

	Mary	Melanie	Molly
Mary			
Melanie	6		2
Molly	4		
Total			

	Mary	Melanie	Molly	Total	Score
Mary	0.091	0.100	0.077	0.268	
Melanie					
Molly	0.364	0.300	0.308	0.972	
Total	1.000	1.000	1.000	3.000	1.000

(c) *Personality*

	Mary	Melanie	Molly
Mary	1	4	
Melanie		1	
Molly	4	9	1
Total			

	Mary	Melanie	Molly	Total	Score
Mary	0.190	0.286	0.184	0.660	
Melanie					
Molly	0.762	0.643	0.725	2.130	
Total	1.000	1.000	1.000	3.000	1.000

▪ **Step 3. Overall preference:**

	Beauty	Intelligence	Personality	Weighted Score
Mary	0.648	0.089	0.220	
Melanie				
Molly	0.230	0.324	0.713	
Weight	0.072	0.227	0.701	1.000

- Thus, Kevin should choose **Molly** with **0.590**.

D. Multi-Objective Linear Programming**

* **Original MOLP problem** in **decision space** (x_1, x_2)

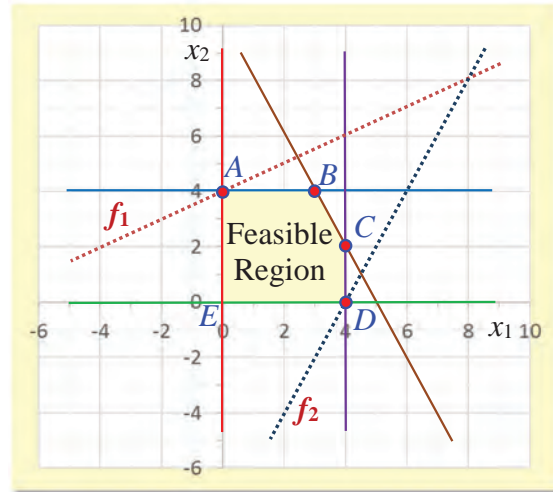
▪ **Objective functions**

$$\text{Max } f_1(\mathbf{x}) = -x_1 + 2x_2$$

$$\text{Max } f_2(\mathbf{x}) = 2x_1 - x_2$$

▪ **Constraints**

- (1) $x_1 \leq 4$
- (2) $x_2 \leq 4$
- (3) $2x_1 + x_2 \leq 10$
- (4) $x_1 \geq 0$
- (5) $x_2 \geq 0$



▪ **Efficient set** = {A, B, C, D}

- Point A = (0, 4)	$f_1(A) = 8$	$f_2(A) = -4$
- Point B = (3, 4)	$f_1(B) = 5$	$f_2(B) = 2$
- Point C = (4, 2)	$f_1(C) = 0$	$f_2(C) = 6$
- Point D = (4, 0)	$f_1(D) = -4$	$f_2(D) = 8$

* **Reformulated MOLP problem** in **criterion space** (f_1, f_2)

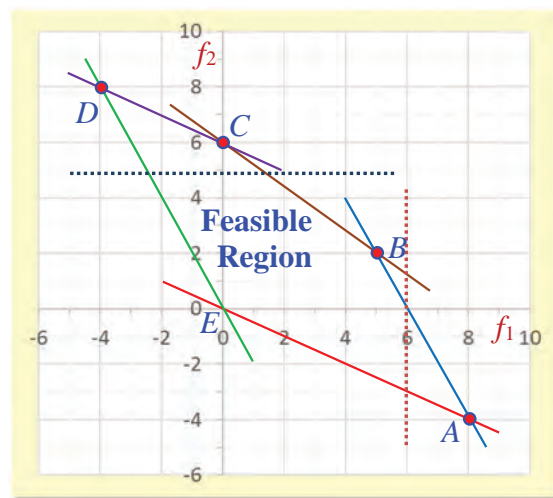
▪ **Objective functions**

$$\text{Max } f_1$$

$$\text{Max } f_2$$

▪ **Constraints**

- (1) $f_1 + 2f_2 \leq 12$
- (2) $2f_1 + f_2 \leq 12$
- (3) $4f_1 + 5f_2 \leq 30$
- (4) $f_1 + 2f_2 \geq 0$
- (5) $2f_1 + f_2 \geq 0$



▪ **Efficient set** = {A, B, C, D}

* Solving MOLP

- **Method 1: Weighted sum** with the weight w .

- Two objective functions in the original LP are

$$f_1(\mathbf{x}) = -x_1 + 2x_2 \text{ and } f_2(\mathbf{x}) = 2x_1 - x_2.$$

Thus, the **weighted sum** is $f_c(\mathbf{x}) = w f_1(\mathbf{x}) + (1 - w) f_2(\mathbf{x})$

$Max f_c(\mathbf{x})$	$f_1(A)=8$ $f_2(A)=-4$	$f_1(B)=5$ $f_2(B)=2$	$f_1(C)=0$ $f_2(C)=6$	$f_1(D)=-4$ $f_2(D)=8$
$w=0.3$	-0.4	2.9	4.2	4.4
$w=0.4$	0.8	3.2	3.6	3.2
$w=0.5$	2.0	3.5	3.0	2.0
$w=0.8$	5.6	4.4	1.2	-1.6

- **Method 2: Goal programming**



- Each objective is viewed as a "goal".
- Deviation variables, a_i and b_i , are the amounts a targeted goal i is **underachieved** or **overachieved**, respectively.
- The goals themselves are added to the **constraint set** with a_i and b_i , acting as the **slack** and **surplus** variables.
- Goal constraints:
 - Goal 1: $f_1(\mathbf{x}) = -x_1 + 2x_2 > s_1 \rightarrow -x_1 + 2x_2 + a_1 - b_1 = s_1$
 - Goal 2: $f_2(\mathbf{x}) = 2x_1 - x_2 > s_2 \rightarrow 2x_1 - x_2 + a_2 - b_2 = s_2$
- Assume that $s_1=6$ and $s_2=5$, and the new **objective function** is $Min z = \$2 a_1 + \$1 a_2$.

$Min z = 2a_1 + 1a_2$	$f_1(A)=8$ $f_2(A)=-4$	$f_1(B)=5$ $f_2(B)=2$	$f_1(C)=0$ $f_2(C)=6$	$f_1(D)=-4$ $f_2(D)=8$
(a_1, b_1) for Goal 1	(0, 2)		(6, 0)	(10, 0)
(a_2, b_2) for Goal 2	(9, 0)		(0, 1)	(0, 3)
Objective function z	\$9		\$12	\$20