

# A slightly better lower bound on the number of edges in (online) list critical graphs

October 15, 2015

## 1 Introduction

Let  $\mathcal{T}_k$  be the Gallai trees with maximum degree at most  $k - 1$ , excepting  $K_k$ . For a graph  $G$ , let  $W^k(G)$  be the set of vertices of  $G$  that are contained in some  $K_{k-1}$  in  $G$ .

**Definition 1.** A graph  $G$  is *AT-reducible* to  $H$  if  $H$  is a nonempty induced subgraph of  $G$  which is  $f_H$ -AT where  $f_H(v) := \delta(G) + d_H(v) - d_G(v)$  for all  $v \in V(H)$ . If  $G$  is not AT-reducible to any nonempty induced subgraph, then it is *AT-irreducible*.

## 2 Reducible Configurations

This first lemma tells us how a single high vertex can interact with the low vertex subgraph. This is the version Hal and I used, it (and more) follows from the classification in “mostlow”.

**Lemma 2.1.** *Let  $k \geq 5$  and let  $G$  be a graph with  $x \in V(G)$  such that:*

1.  $K_k \not\subseteq G$ ; and
2.  $G - x$  has  $t$  components  $H_1, H_2, \dots, H_t$ , and all are in  $\mathcal{T}_k$ ; and
3.  $d_G(v) \leq k - 1$  for all  $v \in V(G - x)$ ; and
4.  $|N(x) \cap W^k(H_i)| \geq 1$  for  $i \in [t]$ ; and
5.  $d_G(x) \geq t + 2$ .

*Then  $G$  is  $f$ -AT where  $f(x) = d_G(x) - 1$  and  $f(v) = d_G(v)$  for all  $v \in V(G - x)$ .*

To deal with more than one high vertex we need the following auxiliary bipartite graph. For a graph  $G$ ,  $\{X, Y\}$  a partition of  $V(G)$  and  $k \geq 4$ , let  $\mathcal{B}_k(X, Y)$  be the bipartite graph with one part  $Y$  and the other part the components of  $G[X]$ . Put an edge between  $y \in Y$  and a component  $T$  of  $G[X]$  iff  $N(y) \cap W^k(T) \neq \emptyset$ . The next lemma tells us that we have a reducible configuration if this bipartite graph has minimum degree at least three.

**Lemma 2.2.** *Let  $k \geq 7$  and let  $G$  be a graph with  $Y \subseteq V(G)$  such that:*

	$k$ -Critical $G$				$k$ -ListCritical $G$		
$k$	Gallai [1] $d(G) \geq$	Kriv [5] $d(G) \geq$	KS [4] $d(G) \geq$	KY [3] $d(G) \geq$	KS [4] $d(G) \geq$	KR [2] $d(G) \geq$	Here $d(G) \geq$
4	3.0769	3.1429	—	3.3333	—	—	
5	4.0909	4.1429	—	4.5000	—	4.0984	<b>4.1</b>
6	5.0909	5.1304	5.0976	5.6000	—	5.1053	<b>5.1082</b>
7	6.0870	6.1176	6.0990	6.6667	—	6.1149	<b>6.1204</b>
8	7.0820	7.1064	7.0980	7.7143	—	7.1128	<b>7.1181</b>
9	8.0769	8.0968	8.0959	8.7500	8.0838	8.1094	<b>8.1143</b>
10	9.0722	9.0886	9.0932	9.7778	9.0793	9.1055	<b>9.1100</b>
15	14.0541	14.0618	14.0785	14.8571	14.0610	14.0864	<b>14.0892</b>
20	19.0428	19.0474	19.0666	19.8947	19.0490	19.0719	<b>19.0738</b>

Table 1: History of lower bounds on the average degree  $d(G)$  of  $k$ -critical and  $k$ -list-critical graphs  $G$ .

1.  $K_k \not\subseteq G$ ; and
2. the components of  $G - Y$  are in  $\mathcal{T}_k$ ; and
3.  $d_G(v) \leq k - 1$  for all  $v \in V(G - Y)$ ; and
4. with  $\mathcal{B} := \mathcal{B}_k(V(G - Y), Y)$  we have  $\delta(\mathcal{B}) \geq 3$ .

Then  $G$  has an induced subgraph  $G'$  that is  $f$ -AT where  $f(y) = d_{G'}(y) - 1$  for  $y \in Y$  and  $f(v) = d_{G'}(v)$  for all  $v \in V(G' - Y)$ .

We also have the following version with asymmetric degree condition on  $\mathcal{B}$ . The point here is that this works for  $k \geq 5$ . As we'll see in the next section, the consequence is that we trade a bit in our size bound for the proof to go through with  $k \in \{5, 6\}$ .

**Lemma 2.3.** *Let  $k \geq 5$  and let  $G$  be a graph with  $Y \subseteq V(G)$  such that:*

1.  $K_k \not\subseteq G$ ; and
2. the components of  $G - Y$  are in  $\mathcal{T}_k$ ; and
3.  $d_G(v) \leq k - 1$  for all  $v \in V(G - Y)$ ; and
4. with  $\mathcal{B} := \mathcal{B}_k(V(G - Y), Y)$  we have  $d_{\mathcal{B}}(y) \geq 4$  for all  $y \in Y$  and  $d_{\mathcal{B}}(T) \geq 2$  for all components  $T$  of  $G - Y$ .

Then  $G$  has an induced subgraph  $G'$  that is  $f$ -AT where  $f(y) = d_{G'}(y) - 1$  for  $y \in Y$  and  $f(v) = d_{G'}(v)$  for all  $v \in V(G' - Y)$ .

### 3 Improved bounds on $\sigma$

Let  $T \in \mathcal{T}_k$ . Then each block of  $T$  is regular. Say  $\text{type}(B) = b$  if  $B$  is  $(b-1)$ -regular. Let  $x \in V(T)$  and let  $B_1, \dots, B_\ell$  be the blocks of  $T$  containing  $x$  where  $B_i$  is of type  $b_i$ . Then we say that  $\text{type}_T(x) = (b_1, \dots, b_\ell)$ . For an endblock  $B$  of  $T$ , let  $x_B$  be the cutvertex of  $T$  contained in  $B$  and put  $T_B := T - (V(B) \setminus \{x\})$ . For  $b \geq 1$ , put  $t(b) := 2 - \frac{2}{b}$ . For  $T \in \mathcal{T}_k$  and  $x \in V(T)$  put

$$\sigma_T(x) := k - 2 + \frac{2}{k-1} - d_T(x).$$

For  $T \in \mathcal{T}_k$  and  $x \in V(T)$  with  $\text{type}_T(x) = (b_1, \dots, b_\ell)$ , put

$$\sigma'_T(x) := \sigma_T(x) - 2 + \sum_{i \in [\ell]} t(b_i).$$

Furthermore, put

$$\sigma(T) := \sum_{x \in V(T)} \sigma_T(x),$$

and

$$\sigma'(T) := \sum_{x \in V(T)} \sigma'_T(x).$$

**Lemma 3.1.** *Let  $T \in \mathcal{T}_k$  where  $k \geq 4$ . Then,*

- (a) *If  $B$  is a block of  $T$ , then  $\sigma(B) = 2$  if  $B = K_{k-1}$  and  $\sigma(B) \geq k - 2 + \frac{2}{k-1}$  otherwise,*
- (b) *If  $B$  is an endblock of  $T$ , then  $\sigma(T) = \sigma(T_B) + \sigma(B) - (k - 2 + \frac{2}{k-1})$ .*

*Proof.* Immediate from the definitions (see Kostochka and Stiebitz [4]). □

**Lemma 3.2.** *If  $T \in \mathcal{T}_k$  and  $k \geq 4$ , then  $\sigma(T) \geq \sigma'(T) + 2$ .*

*Proof.* Suppose the lemma is false and let  $T$  be a counterexample with the minimum number of blocks. First, suppose  $T$  has one block. Then,  $T$  is complete or an odd cycle. If  $T$  is complete, then  $T = K_b$  with  $b \in [k-1]$  and hence  $\sigma'(T) = \sigma(T) + (t(b) - 2)b = \sigma(T) - 2$ . If instead  $T$  is an odd cycle, then  $\sigma'(T) = \sigma(T) + (t(3) - 2)|T| = \sigma(T) - \frac{2}{3}|T| \leq \sigma(T) - 2$ . Hence  $T$  must have at least two blocks.

Let  $B$  be an endblock of  $T$ . Say  $\text{type}(B) = b$  and  $\text{type}_T(x_B) = (b_1, \dots, b_\ell)$  where  $b_\ell = b$ . Then  $\text{type}_{T_B}(x_B) = (b_1, \dots, b_{\ell-1})$ . Therefore, we have

$$\sigma'_{T_B}(x_B) = k - 2 + \frac{2}{k-1} - d_{T_B}(x_B) - 2 + \sum_{i \in [\ell-1]} t(b_i),$$

and

$$\sigma'_B(x_B) = k - 2 + \frac{2}{k-1} - d_B(x_B) + t(b_\ell) - 2,$$

and

$$\sigma'_T(x_B) = k - 2 + \frac{2}{k-1} - d_T(x_B) - 2 + \sum_{i \in [\ell]} t(b_i).$$

Since  $d_T(x_B) = d_{T_B}(x_B) + d_B(x_B)$ , we have

$$\sigma'(T) = \sigma'(T_B) + \sigma'(B) + 2 - \left(k - 2 + \frac{2}{k-1}\right).$$

By our minimality condition on  $T$ , we have

$$\sigma(T_B) \geq \sigma'(T_B) + 2,$$

and

$$\sigma(B) \geq \sigma'(B) + 2.$$

Putting this all together with Lemma 3.1 gives the contradiction

$$\sigma'(T) \leq \sigma(T_B) + \sigma'(B) - \left(k - 2 + \frac{2}{k-1}\right) = \sigma(T) - 2.$$

□

Let  $T \in \mathcal{T}_k$  and  $x \in V(T)$  with  $\text{type}_T(x) = (b_1, \dots, b_\ell)$ . We always have  $\sum_{i \in [\ell]} b_i \leq k + \ell - 1$ . We say that  $x$  is *full* if  $\sum_{i \in [\ell]} b_i = k + \ell - 1$ . When positive integers  $k, b_1, \dots, b_\ell$  are such that  $\sum_{i \in [\ell]} b_i < k + \ell - 1$ , put

$$\Gamma_{k, (b_1, \dots, b_\ell)} := 1 - \frac{3 - 2\ell - \frac{2}{k-1} + \sum_{i \in [\ell]} \frac{2}{b_i}}{k + \ell - 1 - \sum_{i \in [\ell]} b_i},$$

when  $\sum_{i \in [\ell]} b_i = k + \ell - 1$ , put

$$\Gamma_{k, (b_1, \dots, b_\ell)} := \ell.$$

**Lemma 3.3.** *If  $T \in \mathcal{T}_k$ , then  $\sigma'_T(x) \geq \Gamma_{k, \text{type}_T(x)}(k - 1 - d_T(x))$  for all  $x \in V(T)$ .*

*Proof.* We have

$$\sigma'_T(x) = \sigma_T(x) - 2 + \sum_{i \in [\ell]} t(b_i).$$

So,  $\sigma'_T(x) \geq c(k - 1 - d_T(x))$  for some  $c$  and  $x \in V(T)$  with  $\text{type}(x) = (b_1, \dots, b_\ell)$  if and only if

$$(1 - c) \left( k - 1 - \sum_{i \in [\ell]} (b_i - 1) \right) + \sum_{i \in [\ell]} \left( 2 - \frac{2}{b_i} \right) \geq 3 - \frac{2}{k-1}.$$

A quick computation shows that  $\Gamma_{k, (b_1, \dots, b_\ell)}$  is the largest such  $c$  that works when  $x$  is not full. When  $x$  is full, the first term is zero, so  $c$  is irrelevant. We need

$$\sum_{i \in [\ell]} \left( 2 - \frac{2}{b_i} \right) \geq 3 - \frac{2}{k-1}.$$

Since  $x$  is full, we have  $\sum_{i \in [\ell]} b_i = k + \ell - 1$ . In particular, since  $K_k \not\subseteq T$ , we must have  $\ell \geq 2$  and hence  $b_i \geq 2$  for all  $i \in [\ell]$ . So, if  $\ell \geq 3$ , then we have

$$\sum_{i \in [\ell]} \left( 2 - \frac{2}{b_i} \right) \geq \ell \geq 3 - \frac{2}{k-1}.$$

Hence we may assume  $\ell = 2$ . So,  $b_1 + b_2 = k + 1$ . Simplifying the inequality we need gives

$$\frac{k+1}{b_1(k+1-b_1)} = \frac{b_1+b_2}{b_1b_2} \leq \frac{1}{2} + \frac{1}{k-1}.$$

The worst case is when  $b_1 = 2$ , in which case we have

$$\frac{k+1}{b_1(k+1-b_1)} = \frac{1}{2} + \frac{1}{k-1}.$$

□

So, now our task is to prove lower bounds on  $\Gamma_{k,(b_1,\dots,b_\ell)}$ .

**Lemma 3.4.** *Let  $k \geq 6$  and let  $b_1, \dots, b_\ell$  be positive integers. We have*

$$\Gamma_{k,(b_1,\dots,b_\ell)} \geq \begin{cases} 0 & \text{if } \ell = 1 \text{ and } b_1 = k-1, \\ \frac{1}{2} - \frac{1}{(k-1)(k-2)} & \text{if } \ell = 1 \text{ and } b_1 = k-2, \\ 1 - \frac{3k-5}{(k-1)^2} & \text{if } \ell = 1 \text{ and } b_1 = 1, \\ \frac{2}{3} - \frac{4}{3(k-1)(k-3)} & \text{if } \ell = 1 \text{ and } 2 \leq b_1 \leq k-3, \\ 1 - \frac{1}{k-1} & \text{if } \ell = 2, \\ \ell - 2 + \frac{2}{k-1} & \text{if } \ell \geq 3. \end{cases}$$

*Proof.* Just compute. The  $\ell \geq 3$  case can be improved, but not sure we need it. □

The bound in Lemma 3.3 gives  $\sigma'_T(x) \geq 0$  when  $x$  is a full vertex. We can often do better.

**Lemma 3.5.** *Let  $T \in \mathcal{T}_k$  for  $k \geq 5$ . If  $x \in V(T)$  is a full vertex of type  $(b_1, \dots, b_\ell)$ , then*

$$\sigma'_T(x) \geq \begin{cases} 0 & \text{if } \ell = 2, b_1 = 2 \text{ and } b_2 = k-1, \\ \frac{1}{3} - \frac{2}{(k-1)(k-2)} & \text{if } \ell = 2 \text{ and } b_1, b_2 \leq k-2, \\ \frac{1}{3} + \frac{2}{k-1} & \text{if } \ell \geq 3, \end{cases}$$

*Proof.* We have  $\sigma'_T(x) = 2\ell - 3 + \frac{2}{k-1} - \sum_{i \in [\ell]} \frac{2}{b_i}$ . Now, just compute. □

In the proof of the bound in [2], we handled  $K_{k-1}$ 's separately and for the other components we used a bound of  $\frac{1}{2} - \frac{1}{(k-1)(k-2)}$  as a lower bound on  $\Gamma_{k,(b_1,\dots,b_\ell)}$  (which can be seen from Lemma 3.4). As the bound in Lemma 3.4 shows, if we want to improve on this we need to handle the  $\ell = 1$  and  $b_1 = k-2$  case where we get  $\frac{1}{2} - \frac{1}{(k-1)(k-2)}$  exactly. When  $T = K_{k-2}$  is a component, for  $x \in V(T)$  we have

$$\begin{aligned} \sigma'_T(x) &= k-2 + \frac{2}{k-1} - (k-3) + \left(-2 + 2 - \frac{2}{k-2}\right) \\ &= 1 + \frac{2}{k-1} - \frac{2}{k-2} \\ &= \left(\frac{1}{2} - \frac{1}{(k-1)(k-2)}\right) (k-1 - d_T(x)). \end{aligned}$$

So, to improve on  $\frac{1}{2} - \frac{1}{(k-1)(k-2)}$ , we need to do get some extra weight from somewhere. We know  $\sigma(T) \geq \sigma'(T) + 2$  by Lemma 3.2. It turns out that in the proof below, that +2 is wasted on all components except those containing a  $K_{k-1}$ . Since components in general could be arbitrary size, distributing this extra +2 to the vertices of the component will not give us a useful improvement. But in the case that  $T = K_{k-2}$ , this gets us an extra  $\frac{2}{k-2}$  weight for each vertex in  $T$ , which allows us to improve  $\frac{1}{2} - \frac{1}{(k-1)(k-2)}$  to  $\frac{1}{2} + \frac{1}{k-2} - \frac{1}{(k-1)(k-2)} = \frac{1}{2} + \frac{1}{k-1}$  in this case. It remains to find extra weight for  $K_{k-2}$  blocks that are not components. Suppose  $B = K_{k-2}$  is a block in a component  $T$  and let  $x_B$  be the cutvertex of  $T$  contained in  $B$ . We are going to get our extra weight from  $x_B$ . First, if  $x_B$  is not full, then  $\text{type}(x_B) = (2, k-2)$  and Lemma 3.4 gives  $\Gamma_{k, (b_1, b_2)} \geq 1 - \frac{1}{k-1}$ . So,  $x_B$  contributes  $(1 - \frac{1}{k-1})(k-1 - d_T(x_B)) = (1 - \frac{1}{k-1})$  which is an extra  $1 - \frac{1}{k-1} - \left(\frac{1}{2} - \frac{1}{(k-1)(k-2)}\right) = \frac{1}{2} - \frac{k-3}{(k-1)(k-2)}$ . Distributing this evenly amongst the vertices of the  $K_{k-2}$  gets each vertex an extra  $\frac{1}{2(k-2)} - \frac{k-3}{(k-1)(k-2)(k-2)}$ . Suppose instead that  $x_B$  is full. Then Lemma 3.5 shows that there is unused weight  $\frac{1}{3} - \frac{2}{(k-1)(k-2)}$ . Distributing this weight to the at most  $k-3$  non-full vertices in  $T$  gets each vertex an extra  $\frac{1}{3(k-3)} - \frac{2}{(k-1)(k-2)(k-3)}$ . This latter gain is the smaller of the two, adding this to  $\frac{1}{2} - \frac{1}{(k-1)(k-2)}$  gives  $\frac{1}{2} + \frac{k-5}{3(k-2)(k-3)}$ . All of the bounds in Lemma 3.4 are at most this for  $k \geq 5$ , except  $\frac{1}{2} - \frac{1}{(k-1)(k-2)}$  (which we have already fixed) and  $1 - \frac{3k-5}{(k-1)^2}$  which occurs for components with a single vertex. We can easily fix these using the same trick of using the extra +2 like we did for  $K_{k-2}$  components. In the notation of the next section, this proves the following.

**Lemma 3.6.**  $\sigma_k(G) \geq 2|\mathcal{D}| + \left(\frac{1}{2} + \frac{k-5}{3(k-2)(k-3)}\right) \sum_{v \in L'} (k-1 - d_{\mathcal{L}}(v)).$

Below we put this in to get an improved bound in Theorem 4.4.

## 4 The lower bound

We need the following definitions:

$$\begin{aligned} \mathcal{L}_k(G) &:= G[x \in V(G) \mid d_G(x) < k], \\ \mathcal{H}_k(G) &:= G[x \in V(G) \mid d_G(x) \geq k], \\ \sigma_k(G) &:= \left(k - 2 + \frac{2}{k-1}\right) |\mathcal{L}_k(G)| - 2 \|\mathcal{L}_k(G)\|, \\ \tau_{k,c}(G) &:= 2 \|\mathcal{H}_k(G)\| + \left(k - c - \frac{2}{k-1}\right) \sum_{y \in V(\mathcal{H}_k(G))} (d_G(y) - k), \\ g_k(n, c) &:= \left(k - 1 + \frac{k-3}{(k-c)(k-1) + k-3}\right) n. \end{aligned}$$

### 4.1 We only really care about low degree vertices

As proved in [4], a computation gives the following.

**Lemma 4.1.** *Let  $G$  be a graph with  $\delta := \delta(G) \geq 3$  and  $0 \leq c \leq \delta + 1 - \frac{2}{\delta}$ . If  $\sigma_{\delta+1}(G) + \tau_{\delta+1,c}(G) \geq c |\mathcal{H}_{\delta+1}(G)|$ , then  $2 \|G\| \geq g_{\delta+1}(|G|, c)$ .*

With a lower value of  $c$ , we can make it so we only have to care about vertices of degree  $\delta$  and  $\delta + 1$  as follows.

**Lemma 4.2.** *Let  $G$  be a graph with  $\delta := \delta(G) \geq 3$  and  $0 \leq c \leq \frac{\delta+1}{2} + \frac{1}{\delta}$ . Put  $H' := \{v \in V(G) : d_G(v) = \delta + 1\}$ . If  $\sigma_{\delta+1}(G) + \sum_{y \in H'} d_{\mathcal{H}_{\delta+1}}(y) \geq c |H'|$ , then  $2 \|G\| \geq g_{\delta+1}(|G|, c)$ .*

*Proof.* Put  $\mathcal{H} := \mathcal{H}_{\delta+1}$  and  $k := \delta + 1$ . For  $y \in V(\mathcal{H})$ , put

$$\tau_{k,c}(y) := d_{\mathcal{H}}(y) + \left(k - c + \frac{2}{k-1}\right) (d_G(y) - k).$$

We have

$$\begin{aligned} \tau_{k,c}(G) &= \sum_{y \in V(\mathcal{H})} \tau_{k,c}(y) \\ &\geq \sum_{y \in H'} d_{\mathcal{H}}(y) + \sum_{y \in V(\mathcal{H}) \setminus H'} \left(d_{\mathcal{H}}(y) + k - c + \frac{2}{k-1}\right) \\ &\geq \sum_{y \in H'} d_{\mathcal{H}}(y) + \left(k - c + \frac{2}{k-1}\right) |\mathcal{H} - H'| \\ &\geq \sum_{y \in H'} d_{\mathcal{H}}(y) + c |\mathcal{H} - H'|, \end{aligned}$$

where the last inequality follows since  $c \leq \frac{k}{2} + \frac{1}{k-1}$ . Now applying Lemma 4.1 proves the lemma.  $\square$

## 4.2 Finishing the proof

We need the following degeneracy lemma.

**Lemma 4.3.** *Let  $G$  be a graph and  $f : V(G) \rightarrow \mathbb{N}$ . If  $\|G\| > \sum_{v \in V(G)} f(v)$ , then  $G$  has an induced subgraph  $H$  such that  $d_H(v) > f(v)$  for each  $v \in V(H)$ .*

*Proof.* Suppose not and choose a counterexample  $G$  minimizing  $|G|$ . Then  $|G| \geq 3$  and we have  $x \in V(G)$  with  $d_G(x) \leq f(x)$ . But now  $\|G - x\| > \sum_{v \in V(G-x)} f(v)$ , contradicting minimality of  $|G|$ .  $\square$

**Theorem 4.4.** *If  $G$  is an AT-irreducible graph with  $\delta(G) \geq 4$  and  $\omega(G) \leq \delta(G)$ , then  $2 \|G\| \geq g_{\delta(G)+1}(|G|, c)$  where  $c := (\delta(G) - 2) \left(\frac{1}{2} + \frac{\delta(G)-4}{3(\delta(G)-1)(\delta(G)-2)}\right)$  when  $\delta(G) \geq 6$  and  $c := (\delta(G) - 3) \left(\frac{1}{2} + \frac{\delta(G)-4}{3(\delta(G)-1)(\delta(G)-2)}\right)$  when  $\delta(G) \in \{4, 5\}$ .*

*Proof.* Put  $k := \delta(G) + 1$ ,  $\mathcal{L} := \mathcal{L}_k(G)$  and  $\mathcal{H} := \mathcal{H}_k(G)$ . Put  $W := W^k(\mathcal{L})$ ,  $L' := V(\mathcal{L}) \setminus W$  and  $H' := \{v \in V(\mathcal{H}) : d_G(v) = k\}$ . By Lemma 4.2, it will be sufficient to prove that

$$S := \sigma_k(G) + \sum_{y \in H'} d_{\mathcal{H}}(y) \geq c |H'|.$$

Let  $\mathcal{D}$  be the components of  $\mathcal{L}$  containing  $K_{k-1}$ . By Lemma 3.6, we have

$$\sigma_k(G) \geq 2|\mathcal{D}| + \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) \sum_{v \in L'} (k-1-d_{\mathcal{L}}(v)).$$

Now we define an auxiliary bipartite graph  $F$  with parts  $A$  and  $B$  where:

1.  $B = H'$  and  $A$  is the disjoint union of the following sets  $A_1, A_2$  and  $A_3$ ,
2.  $A_1 = \mathcal{D}$  and each  $T \in \mathcal{D}$  is adjacent to all  $y \in H'$  where  $N(y) \cap W^k(T) \neq \emptyset$ ,
3. For each  $v \in L'$ , let  $A_2(v)$  be a set of  $|N(v) \cap H'|$  vertices connected to  $N(v) \cap H'$  by a matching in  $F$ . Let  $A_2$  be the disjoint union of the  $A_2(v)$  for  $v \in L'$ ,
4. For each  $y \in H'$ , let  $A_3(y)$  be a set of  $d_{\mathcal{H}}(y)$  vertices which are all joined to  $y$  in  $F$ . Let  $A_3$  be the disjoint union of the  $A_3(y)$  for  $y \in H'$ .

**Case 1.**  $\delta \geq 6$ .

Define  $f: V(F) \rightarrow \mathbb{N}$  by  $f(v) = 1$  for all  $v \in A_2 \cup A_3$  and  $f(v) = 2$  for all  $v \in B \cup A_1$ . First, suppose  $\|F\| > \sum_{v \in V(F)} f(v)$ . Then by Lemma 4.3,  $F$  has an induced subgraph  $Q$  such that  $d_Q(v) > f(v)$  for each  $v \in V(Q)$ . In particular,  $V(Q) \subseteq B \cup A_1$  and  $\delta(Q) \geq 3$ . Put  $Y := B \cap V(Q)$  and let  $X$  be  $\bigcup_{T \in V(Q) \cap A_1} V(T)$ . Now  $H := G[X \cup Y]$  satisfies the hypotheses of Lemma 2.2, so  $H$  has an induced subgraph  $G'$  that is  $f$ -AT where  $f(y) = d_{G'}(y) - 1$  for  $y \in Y$  and  $f(v) = d_{G'}(v)$  for  $v \in X$ . Since  $Y \subseteq H'$  and  $X \subseteq \mathcal{L}$ , we have  $f(v) = \delta(G) + d_{G'}(v) - d_G(v)$  for all  $v \in V(G')$ . Hence,  $G$  is AT-reducible to  $G'$ , a contradiction.

Therefore  $\|F\| \leq \sum_{v \in V(F)} f(v) = 2(|H'| + |\mathcal{D}|) + |A_2| + |A_3|$ . By Lemma 2.1, for each  $y \in B$  we have  $d_F(y) \geq k-1$ . Hence  $\|F\| \geq (k-1)|H'|$ . This gives  $(k-3)|H'| \leq 2|\mathcal{D}| + |A_2| + |A_3|$ . By our above estimate we have  $S \geq 2|\mathcal{D}| + \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) \sum_{v \in L'} (k-1-d_{\mathcal{L}}(v)) + \sum_{y \in H'} d_{\mathcal{H}}(y) = 2|\mathcal{D}| + \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) |A_2| + |A_3| \geq \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) (2|\mathcal{D}| + |A_2| + |A_3|)$ . Hence  $S \geq \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) (k-3)|H'|$ . Thus our desired bound holds by Lemma 4.2.

**Case 2.**  $\delta \in \{4, 5\}$ .

Define  $f: V(F) \rightarrow \mathbb{N}$  by  $f(v) = 1$  for all  $v \in A_1 \cup A_2 \cup A_3$  and  $f(v) = 3$  for all  $v \in B$ . First, suppose  $\|F\| > \sum_{v \in V(F)} f(v)$ . Then by Lemma 4.3,  $F$  has an induced subgraph  $Q$  such that  $d_Q(v) > f(v)$  for each  $v \in V(Q)$ . In particular,  $V(Q) \subseteq B \cup A_1$  and  $d_Q(v) \geq 4$  for  $v \in B \cap V(Q)$  and  $d_Q(v) \geq 2$  for  $v \in A_1 \cap V(Q)$ . Put  $Y := B \cap V(Q)$  and let  $X$  be  $\bigcup_{T \in V(Q) \cap A_1} V(T)$ . Now  $H := G[X \cup Y]$  satisfies the hypotheses of Lemma 2.3, so  $H$  has an induced subgraph  $G'$  that is  $f$ -AT where  $f(y) = d_{G'}(y) - 1$  for  $y \in Y$  and  $f(v) = d_{G'}(v)$  for  $v \in X$ . Since  $Y \subseteq H'$  and  $X \subseteq \mathcal{L}$ , we have  $f(v) = \delta(G) + d_{G'}(v) - d_G(v)$  for all  $v \in V(G')$ . Hence,  $G$  is AT-reducible to  $G'$ , a contradiction.

Therefore  $\|F\| \leq \sum_{v \in V(F)} f(v) = 3|H'| + |\mathcal{D}| + |A_2| + |A_3|$ . By Lemma 2.1, for each  $y \in B$  we have  $d_F(y) \geq k-1$ . Hence  $\|F\| \geq (k-1)|H'|$ . This gives  $(k-4)|H'| \leq |\mathcal{D}| + |A_2| + |A_3|$ . By our above estimate we have  $S \geq 2|\mathcal{D}| + \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) \sum_{v \in L'} (k-1-d_{\mathcal{L}}(v)) + \sum_{y \in H'} d_{\mathcal{H}}(y) = 2|\mathcal{D}| + \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) |A_2| + |A_3| \geq \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) (|\mathcal{D}| + |A_2| + |A_3|)$ . Hence  $S \geq \left( \frac{1}{2} + \frac{k-5}{3(k-2)(k-3)} \right) (k-4)|H'|$ . Thus our desired bound holds by Lemma 4.2.  $\square$



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