Proof of Brooks' theorem. Suppose the theorem is false and choose a counterexample G minimizing |G|. Let  $\{A_1, A_2\}$  be a separation of G minimizing  $k := |A_1 \cap A_2|$ . If  $A_1 \cap A_2$  is a clique, then by minimality of |G|, we have  $\Delta$ -colorings of each  $G[A_i]$  which use k colors on  $A_1 \cap A_2$ . By permuting color names if necessary we can combine these to get a  $\Delta(G)$ -coloring of G, a contradiction. In particular,  $k \geq 2$  and if k = 2, then  $A_1 \cap A_2$  is independent.

Suppose k=2 and say  $A_1 \cap A_2 = \{u,v\}$ . By symmetry we may assume that every  $\Delta$ -coloring of  $G[A_1]$  gives u and v different colors and every  $\Delta$ -coloring of  $G[A_2]$  gives u and v the same color (otherwise we could combine the colorings into a  $\Delta$ -coloring of G as above). By minimality of k, each of u and v have a neighbor on both sides of the separation. Hence u and v must each have  $\Delta - 1$  neighbors in  $A_1$  and 1 neighbor in  $A_2$ . But then since  $\Delta \geq 3$ , any  $\Delta$ -coloring of  $G[A_2 - \{u,v\}]$  can be extended to a  $\Delta$ -coloring of  $G[A_2]$  where u and v receive the same color, a contradiction.

Thence  $k \geq 3$ . Since G is not a disjoint union of cliques, we have an induced  $P_3$  xyz in G. Since  $k \geq 3$ , G - x - z is connected and hence we may order V(G) as  $x, z, v_1, \ldots, v_n, y$  so that each  $v_i$  has a neighbor to the right. But this is a contradiction since greedily coloring in this order uses at most  $\Delta$  colors as x and z get the same color, each  $v_i$  has at most  $\Delta - 1$  neighbors to the left and y has two neighbors (x and z) colored the same.