

CS253 Architectures II

Lecture 5

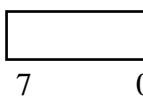
Assembly Language

(Data Types, Integers, Real and String)

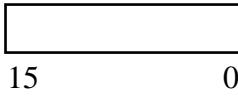
Charles Markham

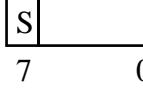
Data format (unsigned integer) in PC's

So far we have used the data stored in the registers in the following way, the formats are very similar to that used in C.

unsigned char  Unsigned, $0 \leq X \leq 255$
1 byte, 8bits, (al,dL,bh etc)

Reminder
of c

unsigned int  Unsigned, $0 \leq X \leq 65535$
2 bytes, 16bits, (ax,bx,dx etc)

signed char  Twos Complement, $-128 \leq X \leq 127$
1 byte, 8bits, (al,dL,bh etc)

signed int  Twos Complement, $-32768 \leq X \leq 32767$
2 bytes, 16bits, (ax,bx,dx etc)

Data format (signed integer) in 8086

There are other data formats that instructions assume, the following formats are used by the 80386 and 80387 coprocessor.

Word	<table border="1"><tr><td>S</td><td></td></tr><tr><td>15</td><td>0</td></tr></table>	S		15	0	Twos Complement, $-32768 < X < 32768$
S						
15	0					
Short int	<table border="1"><tr><td>S</td><td></td></tr><tr><td>31</td><td>0</td></tr></table>	S		31	0	Twos Complement, $-2 \times 10^9 < X < 2 \times 10^9$
S						
31	0					
Long int	<table border="1"><tr><td>S</td><td></td></tr><tr><td>63</td><td>0</td></tr></table>	S		63	0	Twos Complement, $-9 \times 10^{18} < X < 9 \times 10^{18}$
S						
63	0					

Data format (integer) in PC's

Packed Decimal	<table border="1"><tr><td>S</td><td>X</td><td>D₁₇</td><td>D₀</td></tr><tr><td>79</td><td>72</td><td></td><td>0</td></tr></table>	S	X	D ₁₇	D ₀	79	72		0	$-99\dots 99 < X < 99\dots 99$ (18 digits)
S	X	D ₁₇	D ₀							
79	72		0							

Each byte stores two BCD digits, each digit requires 4 bits

Twos Complement

You could use a single bit to set the sign of the number.
However if you do this then you need separate hardware/software for addition and subtraction. Two's complement avoids the problem.

By inverting each bit in a number and adding one to the result. A new number is formed that has the same effect when added to a number as subtracting the original number. If the MSB is 1 then the number is negative, twos complement, otherwise positive ordinary binary.

Twos Complement

Subtract 65 from 120 using the twos complement method.

120: 0 1 1 1 1 0 0 0

65: 0 1 0 0 0 0 0 1

Invert bits: 1 0 1 1 1 1 1 0

Add 1: 1 0 1 1 1 1 1 1

Add to 120 0 0 1 1 0 1 1 1 :55

128, 64, 32, 16, 8, 4, 2, 1

Negative numbers

The following code segment converts negative numbers to their absolute (positive) value. Add it to the start of the print algorithm to print signed numbers, it assumes the number is stored in AX.

```
opcode → mov cx,0  
          cmp ax,0  
          jge past ; if ax greater than or equal to zero  
          push ax  
          mov ah,02h  
          mov dl,'-' ] ; Print a minus sign  
          int 021h  
          pop ax  
          neg ax ; Twos complement ax=|ax|  
  
past:
```

Strings

A string is a group of ASCII characters used to describe text.

Two formats are used

String terminated with a NULL byte=0 character

Charles Markham /0

String using first byte to set length

/15 Charles Markham

ASCII: American standard code for information interchange.

String operations

A string is a set of bytes stored in successive memory locations.
The 8086 contains instructions for the manipulation of strings,
moving them and searching for text.

Registers SI (Source) and DI(Destination) are used for
string manipulation.

LEA SI,*Label* Load effective address in to SI

MOVSB Move string byte

CMPSB Compare string

movsb

Copies memory pointed to by DS:SI to the address specified in
ES:DI

If DF=0 SI,DI are incremented,

If DF=1 SI,DI are decreased.

Note a modifier ‘REP’ can be placed in front of any string
instruction and repeats the instruction CX times. Rep stands for
repeat string prefix.

CISC!!!

cmpsb

Compares memory byte pointed to by DS:SI to the memory byte pointed to by ES:DI. The command is normally used with a REPE modifier as part of a large string test loop.

REPE: Repeat while equal

RENE:Repeat while not equal

If DF=0 SI,DI are incremented after each comparison,

If DF=1 SI,DI are decreased after each comparison.

e.g. REPE CMPSB

 JNE Strings are not equal

.DATA

```
string1 db 'This is a test$'  
string2 db 'abcdefghijklmнопqrstuvwxyz$'
```

Repeat instruction until
CX=0, (modifier, vector?).

.CODE

.STARTUP

```
mov dx, OFFSET string2  
mov ah,09h  
int 021h
```

] BIOS Print until \$

```
mov ax,ds ; ES=DS  
mov es,ax  
LEA SI,string1 ; Offset  
LEA DI,string2 ;  
CLD ; Direction=0  
mov CX,15 ; No. Bytes  
REP MOVS B ; Copy
```

] Copy string

```
mov dx, OFFSET string2  
mov ah,09h  
int 021h
```

] BIOS Print until \$

.EXIT Output: abcdefghijklmнопqrstuvwxyz This is a test

Multiply

Many CISC processors include a multiply and divide function in the instruction set, many RISC processors don't. Can you write MASM code to multiply without using the MUL operation?

Decimal	Binary
1 2 3 4	1 1 0 1
x	x
12	101
$\frac{2468}{1234}$	$\frac{1101}{0000}$
+ 1 2 3 4 0 8	+ 1 1 0 1 0 0
1 4 8 0 8	1 0 0 0 0 0 1

Multiply (The Hard Way, RISC)

.STARTUP

```
mov ax,0
mov dx,100    ;Multiply dl by bl result in ax.
mov bx,123
mov cx,8

back: rcr dx,1      ;move lsb dx into c
       jnc over
       add ax,bx

over: shl bx,1      ;multiply bx by 2
       loop back     ;repeat 8 times

call Print
.EXIT          RISC: Reduced Instruction Set Computer
```

Multiply (The Easy Way, CISC)

```
.STARTUP  
  
mov dl,100    ;Multiply dl by bl result in ax.  
mov al,123  
  
mul dl        ;ax=dl*al  
  
call Print  
  
.EXIT
```

CISC: Complex Instruction Set

Divide

Decimal

$$\begin{array}{r} 0\ 1\ 1\ 2 \\ 1\ 1 \overline{\) 1\ 2\ 3\ 4} \\ 0 \\ 1\ 2 \\ 1\ 1 \\ 1\ 3 \\ 1\ 1 \\ 2\ 4 \\ 2\ 2 \\ 2 \end{array}$$

Binary

$$\begin{array}{r} 0\ 0\ 1\ 0 \\ 1\ 0\ 1 \overline{\) 1\ 1\ 0\ 1} \\ 1\ 0\ 1 \\ 0\ 0\ 1\ 1 \end{array}$$

Take the first digit (1) subtract 101, won't go, write down zero, take first and second digit (11) subtract 101, won't go write down zero, take first three digits (110) subtract 101, will go, write down a 1...., repeat until all digits are analysed.

Data format (float) in PC's

Short real	<table border="1"> <tr> <td>S</td><td>Exponent</td><td>Significand</td></tr> </table>	S	Exponent	Significand	-1.2x10 ⁻³⁸ < X < 3.4x10 ³⁸
S	Exponent	Significand			
	31 23 0				
Long real	<table border="1"> <tr> <td>S</td><td>Exponent</td><td>Significand</td></tr> </table>	S	Exponent	Significand	-2.3x10 ⁻³⁰⁸ < X < 1.7x10 ³⁰⁸
S	Exponent	Significand			
	63 52 0				
Temp real	<table border="1"> <tr> <td>S</td><td>Exponent</td><td>Significand</td></tr> </table>	S	Exponent	Significand	-3.4x10 ⁻⁴⁹³² < X < 1.1x10 ⁴⁹³²
S	Exponent	Significand			
	79 52 0				

Note: exponents are offset to speed up magnitude comparison.

Short exponent biased by +127 (7FH)

Note: Double precision division on the P90 was only calculated to single precision.

Floating point

Express 178.625 as a short real.

Use 4 bytes to represent the digits in the number.

178D=10110010B, Note: the binary always starts with a 1.

0.625D=0.101B, 4, 2, 1, 1/2, 1/4, 1/8,.....,

178.625D=10110010.101B

$$=1.0110010101 \times 2^7$$

Add bias gives 1.0110010101x2⁷⁺¹²⁷=1.0110010101x2¹³⁴

134D=10000110B

Result=0,10000110,0110010101000000000000000



In hex 4332A000

Floating Point Calculation

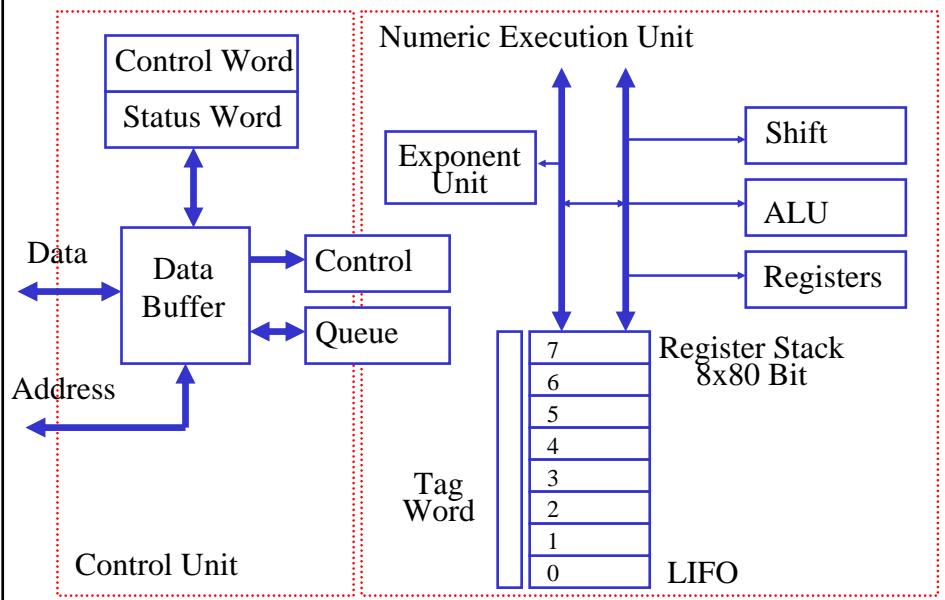
An optional coprocessor is available for many microprocessors, these chips are designed to speed up floating point calculation by as much as 100 times. From the 486 on, all Intel/PC Processors have contained an integrated co-processor.

A co-processor is a second processor, not just an enhancement to the existing processor.

The addition of a coprocessor gives an increased instruction set.

To separate coprocessor instructions from standard instructions each coprocessor instruction starts with 11011.....

8087 Architecture



The Status Word

15MSB

0LSB



B: Busy

C2-0: Condition bits

S2-0: Top of stack

IR: Interrupt request

PE: Precision

UE: Underflow

OE: Overflow

DE: Denormalized operand

IE: Invalid operation

The ControlWord

15MSB

0LSB



IC: Infinity control

R1-0: Rounding control

P1-0: Precision control

IR: Interrupt request

IE: Interrupt Enable Mask

PM: Precision

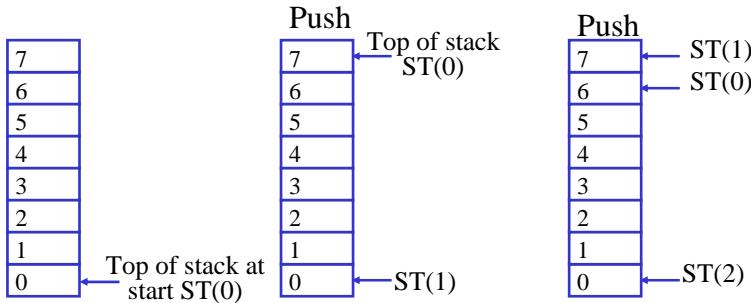
UM: Underflow

OM: Overflow

DM: Denormalized operand

IM: Invalid operation

8087 Stack



Each item on the stack is 10bytes, 80 bits, long double.

Stack pointer decreases on each push.

The stack can go around the clock, i.e. past 111,000.

In MASM ST(0) means the item at the top of the stack.

8087 Instructions

Addition

FADD //source/destination,*source*

Examples:

FADD ST(3),ST(0) ;Add ST(3) to ST(0) result in ST(3)

FADD ;Add ST(1) and ST(0) pop 1
number off the stack, result in ST(0)

FADD ST,ST(1) ;Add Stack top ST(0) to ST(1) result
placed in ST(0)

8087 Instructions

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placed in ST(0)

8087 Arithmetic Instructions

Addition/Subtract:

FADD //*s/d,s* : Add source to real at specified destination.

FSUB //*s/d,s* : Subtract real from specified source result in destination.

FSUBR //*s/d,s* : Subtract d-source from source, result in destination.

Multiplication:

FMUL //*s/d,s* : Multiply d-source by s-real..

Also FMULP, FIMUL

Division

FDIV //*s/d,s* : Divide d-source by s-real, result in destination.

Also FDIVR, FDIP, FIDIV etc.