

## Learning Goals

During this lab, you will:

- review the Simplified Master Theorem
- solve recurrences by iteration and using the S.M.T.
- identify recurrences that can not be solved using the S.M.T.

## Simplified Master Theorem

The **master theorem** is a powerful tool in the analysis and classification of recurrences. It may be used to easily *classify* recurrences that might otherwise be very time-consuming!

### Simplified Master Theorem

Given a recurrence  $T(n)$  of the form,

$$T(n) = \begin{cases} c & n < c_1 \\ aT(n/b) + \Theta(n^i) & n \geq c_1 \end{cases}$$

Case I: If  $a > b^i$  then  $T(n) = \Theta(n^{\log_b a})$ .

Case II: If  $a = b^i$  then  $T(n) = \Theta(n^i \log_b n)$ .

Case III: If  $a < b^i$  then  $T(n) = \Theta(n^i)$ .

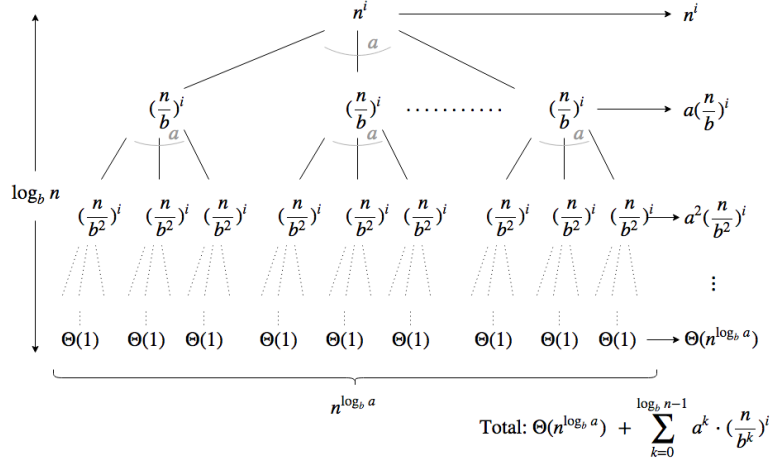
This probably seems very magical and hand-wavy. But we're computer scientists, so let's delve in and figure this out. Let us assume here that  $n$  is some power of  $b$ .



*You don't need to be a wizard to understand the master theorem!*

## Understanding the S.M.T.

As shown in the diagram on the next page, at the  $k$ th level of iteration, there are  $a^k$  subdivisions of  $(n/b^k)^i$  work. Therefore, at the bottom-most level there are  $a^{\log_b n}$  subdivisions of  $(n/b^{\log_b n})^i = 1$  work. Recall that by the properties of logarithms,  $a^{\log_b n} = n^{\log_b a}$  — so we can switch the base and the contents of the log.



Depicted above is a tree-representation of the work performed at each level of iteration in the recurrence.

We can therefore write the total amount of work represented by the recurrence  $T(n)$  as,

$$\text{Total Work: } \Theta(n^{\log_b a}) + \sum_{k=0}^{\log_b n - 1} a^k \cdot \left(\frac{n}{b^k}\right)^i$$

So what does this mean in the three cases shown in the simplified master theorem?

In the case where  $a > b^i$ , the work done at the leaves *heavily* outgrows that done at the root. That is,  $n^{\log_b a}$  is the dominating term in the sum and the total work is therefore  $\Theta(n^{\log_b a})$ .

In the case where  $a = b^i$ , the work done at each level is the same and the total work is just the height of the tree multiplied by the work at each level:  $\Theta(n^i \log_b n)$ .

In the case where  $a < b^i$ , then the work done at each subsequent level decreases with respect to the root, and the work done at the root dominates:  $\Theta(n^i)$ .

Tada! By drawing the recurrence tree and summing the total work performed at each level, we were able to find general expressions for the recurrence solutions for each case of the S.M.T.

## Problems

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