# 目录

1	Hopcroft's algorithm	3
	1.0.1 algorithm	3
	1.0.2 Example	3
	1.1 Minimization by equivalence of states	27
	References	27

## Chapter 1

# Hopcroft's algorithm

## 1.0.1 algorithm

Member function min\_Hopcroft implements Hopcroft's  $n \log n$  minimization algorithm, as presented in [[WATSON94b], Algorithm 4.8].

The combination of the out-transitions of all of the States is stored in a  $\mathbf{CRSet}\ C$ .

Set L from the abstract algorithm is implemented as a mapping from States to int (an array of int is used).

Array L should be interpreted as follows: if State q a representative, then the following pairs still require processing (are still in abstract set L):

$$([q], C_0), \cdots, ([q], C_{L(q)-1})$$

The remaining pairs do not require processing:

$$([q], C_{L[q])}), \cdots, ([q], C_{|C|-1})$$

This implementation facilitates quick scanning of L for the next valid State-CharRange pair.

#### 1.0.2 Example

```
CRSet C; // the out labels of State's: { 'a' 'b' }
int L[|Q|]: // the index of L = q: 对应等价类[q]; L[q]表示正在处理[q]对
应的字符在C中的index,
0 1 2 3 4 5 6 7 8
0 0 0 0 0 0 0 0 0

Initialize partitions, E_0:
{ 0 1 2 3 4 5 }
{ 6 7 8 }
```

#### Algorithm 1 Hopcroft's minimization algorithm

19: end while

```
Input: \mathscr{A} = (Q, i, F)
Output: The equivalence classes of Q
1: \hat{P} \leftarrow \{F, F^c\}
                           ▶ The initial partition
 2: W \leftarrow \emptyset
                      ▶ The waiting set
3: for all a \in A do
       ADD((min(F,F^c),a),W)
                                            ▷ initialization of the waiting set
 5: end for
 6: while W \neq \emptyset do
 7:
        (W,a) \leftarrow TakeSome(W)
                                           ▶ Take and remove some splitter
        for all P \in \hat{P} do which is split by (W, a)
 8:
           P',P'' \leftarrow (W,a)|P|
                                     ▷ Compute the split
 9:
           Replace P by P' and P'' in \hat{P}
                                                ▶ Refine the partition
10:
        end for
11:
12:
        for all b \in A do
                                     ▷ Update the waiting set
           if (P,b) \in W then
13:
               Replace (P,b) by (P',b) and (P'') in W
14:
           else
15:
               ADD((min(P',P''),b),W)
16:
           end if
17:
        end for
18:
```

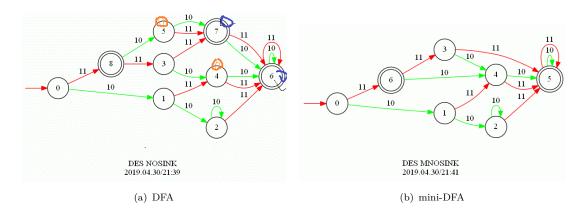


图 1.1: example minimization

```
Initialize L repr{F}: { 6 }
L:
0 1 2 3 4 5 6 7 8
0 0 0 0 0 0 2 0 0

--- while each [q], (split [p] w.r.t ([q],a))
L:
0 1 2 3 4 5 6 7 8
0 0 0 0 0 0 2 0 0
```

```
Pick one [q] in L, Processing [q] index of L = [6], 'b'
current all partitions:
\{ 0 \ 1 \ 2 \ 3 \ 4 \ 5 \}
{ 6 7 8 }
== for each [p], (split [p] w.r.t (index of L)[6], 'b')
==split [0] w.r.t (index of L) [6], 'b')
before split, partitions:
\{ 0 \ 1 \ 2 \ 3 \ 4 \ 5 \}
{ 6 7 8 }
new split of [0] is [1]
after split, partitions:
\{0\ 2\ 3\ 4\ 5\}
{ 1 }
{ 6 7 8 }
before L:
0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8
0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0
p and r are the new representatives. Now update L with the smallest of
   [0],[1]
using [r] = [1], L[r] = C. size();
affter L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
==split [6] w.r.t (index of L) [6], 'b')
before split, partitions:
\{ 0 \ 2 \ 3 \ 4 \ 5 \}
{ 1 }
{ 6 7 8 }
new split of [6] is [8]
after split, partitions:
\{0\ 2\ 3\ 4\ 5\}
{ 1 }
```

```
\{67\}
{ 8 }
before L:
0 1 2 3 4 5 6 7 8
0 \ 2 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0
p and r are the new representatives. Now update L with the smallest of
    [6], [8]
using [r] = [8], L[r] = C. size();
affter L:
0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8
0\ 2\ 0\ 0\ 0\ 0\ 1\ 0\ 2
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
0 \ 2 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [1], 'b'
current all partitions (eq. classes) repr:
\{0\ 1\ 6\ 8\}
current all partitions:
\{0\ 2\ 3\ 4\ 5\}
{ 1 }
\{67\}
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[1], 'b')
==split [0] w.r.t (index of L)[1], 'b')
before split, partitions:
StateEqRel
\{0\ 2\ 3\ 4\ 5\}
{ 1 }
{ 6 7 }
{ 8 }
new split of [0] is [-1]
==split[1] w.r.t (index of L)[1], 'b')
before split, partitions:
\{ 0 \ 2 \ 3 \ 4 \ 5 \}
```

```
{ 1 }
\{67\}
{ 8 }
new split of [1] is [-1]
==split[6] w.r.t (index of L)[1], 'b')
before split, partitions:
\{0\ 2\ 3\ 4\ 5\}
{ 1 }
{ 6 7 }
{ 8 }
new split of [6] is [-1]
==split [8] w.r.t (index of L)[1], 'b')
before split, partitions:
StateEqRel
\{ 0 \ 2 \ 3 \ 4 \ 5 \}
{ 1 }
\{67\}
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8
0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [1], 'a'
current all partitions (eq. classes) repr:
\{0\ 1\ 6\ 8\}
current all partitions:
\{0\ 2\ 3\ 4\ 5\}
{ 1 }
{ 6 7 }
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[1], 'a')
=split [0] w.r.t (index of L)[1], 'a')
before split, partitions:
\{0\ 2\ 3\ 4\ 5\}
{ 1 }
{ 6 7 }
```

```
{ 8 }
new split of [0] is [2]
after split, partitions:
StateEqRel
{ 0 }
{ 1 }
\{2 \quad 3 \quad 4 \quad 5 \}
{ 6 7 }
{ 8 }
before L:
L:
0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8
0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 2
p and r are the new representatives. Now update L with the smallest of
using [p] = [0], L[r]=L[p]; L[p]=C. size();
affter L:
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
2 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 2
==split[1] w.r.t (index of L)[1], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 4 5 }
{ 6 7 }
{ 8 }
new split of [1] is [-1]
==split[6] w.r.t (index of L)[1], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 4 5 }
{ 6 7 }
{ 8 }
```

```
new split of [6] is [-1]
==split [8] w.r.t (index of L)[1], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 4 5 }
{ 6 7 }
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0 1 2 3 4 5 6 7 8
2 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [0], 'b'
current all partitions (eq. classes) repr:
\{ 0 \quad 1 \quad 2 \quad 6 \quad 8 \}
current all partitions:
StateEqRel
\{0\}
{ 1 }
\{2 \quad 3 \quad 4 \quad 5\}
{ 6 7 }
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[0], 'b')
\Longrightarrow split [0] w.r.t (index of L) [0], 'b')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
\{2 \quad 3 \quad 4 \quad 5\}
\{67\}
{ 8 }
new split of [0] is [-1]
=split[1] w.r.t (index of L)[0], 'b')
before split, partitions:
StateEqRel
\{0\}
```

```
{ 1 }
{ 2 3 4 5 }
\{67\}
{ 8 }
new split of [1] is [-1]
=split [2] w.r.t (index of L) [0], 'b')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 3 4 5 }
\{67\}
{ 8 }
new split of [2] is [-1]
=split [6] w.r.t (index of L) [0], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
\{2 \quad 3 \quad 4 \quad 5\}
{ 6 7 }
{ 8 }
new split of [6] is [-1]
==split[8] w.r.t (index of L)[0], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
\{2 \quad 3 \quad 4 \quad 5\}
{ 6 7 }
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
1 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [0], 'a'
```

```
current all partitions (eq. classes) repr:
\{ 0 \ 1 \ 2 \ 6 \ 8 \}
current all partitions:
StateEqRel
{ 0 }
{ 1 }
\{2 \quad 3 \quad 4 \quad 5 \}
{ 6 7 }
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[0], 'a')
=split [0] w.r.t (index of L) [0], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
\{2 \quad 3 \quad 4 \quad 5\}
\{67\}
{ 8 }
new split of [0] is [-1]
==split[1] w.r.t (index of L)[0], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
\{2 \quad 3 \quad 4 \quad 5\}
{ 6 7 }
{ 8 }
new split of [1] is [-1]
==split[2] w.r.t (index of L)[0], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 3 4 5 }
\{67\}
{ 8 }
new split of [2] is [-1]
```

```
==split [6] w.r.t (index of L) [0], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
\{2 \quad 3 \quad 4 \quad 5 \}
\{67\}
{ 8 }
new split of [6] is [-1]
=split [8] w.r.t (index of L) [0], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
\{2 \quad 3 \quad 4 \quad 5 \}
\{67\}
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0 1 2 3 4 5 6 7 8
0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [6], 'a'
current all partitions (eq. classes) repr:
\{ 0 \ 1 \ 2 \ 6 \ 8 \}
current all partitions:
StateEqRel
\{0\}
{ 1 }
\{2 \quad 3 \quad 4 \quad 5\}
\{67\}
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[6], 'a')
==split [0] w.r.t (index of L) [6], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
```

```
\{2 \quad 3 \quad 4 \quad 5\}
\{67\}
{ 8 }
new split of [0] is [-1]
=split[1] w.r.t (index of L)[6], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 4 5 }
{ 6 7 }
{ 8 }
new split of [1] is [-1]
==split [2] w.r.t (index of L) [6], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 3 4 5 }
\{67\}
{ 8 }
new split of [2] is [4]
after split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
before L:
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 2
p and r are the new representatives. Now update L with the smallest of
   [2], [4]
using [p] = [2], L[r]=L[p]; L[p]=C. size();
```

```
affter L:
L:
0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8
0 \ 0 \ 2 \ 0 \ 0 \ 0 \ 0 \ 0 \ 2
==split[6] w.r.t (index of L)[6], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [6] is [-1]
==split [8] w.r.t (index of L) [6], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0 1 2 3 4 5 6 7 8
0 \ 0 \ 2 \ 0 \ 0 \ 0 \ 0 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [2], 'b'
current all partitions (eq. classes) repr:
{ 0 1 2 4 6 8 }
current all partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
\{67\}
{ 8 }
```

```
== for each [p], (split [p] w.r.t (index of L)[2], 'b')
=split [0] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [0] is [-1]
==split[1] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [1] is [-1]
=split[2] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [2] is [-1]
==split [4] w.r.t (index of L) [2], 'b')
before split, partitions:
{\bf State Eq Rel}
\{0\}
{ 1 }
{ 2 3 }
```

```
{ 4 5 }
\{67\}
{ 8 }
new split of [4] is [-1]
==split[6] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [6] is [-1]
==split[8] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
0 \ 0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [2], 'a'
current all partitions (eq. classes) repr:
\{ 0 \ 1 \ 2 \ 4 \ 6 \ 8 \}
current all partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
\{4 \quad 5 \}
\{67\}
```

```
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[2], 'a')
=split [0] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [0] is [-1]
=split[1] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [1] is [-1]
==split[2] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [2] is [3]
after split, partitions:
{\bf State Eq Rel}
\{0\}
{ 1 }
{ 2 }
```

```
{ 3 }
\{4 \quad 5\}
\{67\}
{ 8 }
before L:
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 2
p and r are the new representatives. Now update L with the smallest of
using [p] = [2], L[r]=L[p]; L[p]=C. size();
affter L:
L:
0\ 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8
0 \ 0 \ 2 \ 0 \ 0 \ 0 \ 0 \ 0 \ 2
胡
==split [4] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [4] is [-1]
==split[6] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [6] is [-1]
```

```
==split[8] w.r.t (index of L)[2], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0 1 2 3 4 5 6 7 8
0 \ 0 \ 1 \ 0 \ 0 \ 0 \ 0 \ 2
Pick one [q] in L, Processing [q] index of L = [2], 'a'
current all partitions (eq. classes) repr:
\{ 0 \ 1 \ 2 \ 3 \ 4 \ 6 \ 8 \}
current all partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[2], 'a')
=split [0] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
\{67\}
{ 8 }
```

```
new split of [0] is [-1]
==split[1] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
\{4 \quad 5\}
{ 6 7 }
{ 8 }
new split of [1] is [-1]
==split[2] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [2] is [-1]
==split[3] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [3] is [-1]
==split [4] w.r.t (index of L) [2], 'a')
before split, partitions:
StateEqRel
{ 0 }
```

```
{ 1 }
{ 2 }
{ 3 }
\{45\}
\{67\}
{ 8 }
new split of [4] is [-1]
==split[6] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [6] is [-1]
==split [8] w.r.t (index of L)[2], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 2
Pick one [q] in L, Processing [q] = index of L = [8], 'b'
current all partitions (eq. classes) repr:
{ 0 1 2 3 4 6 8 }
current all partitions:
StateEqRel
```

```
{ 0 }
{ 1 }
{ 2 }
{ 3 }
\{4 \quad 5 \}
{ 6 7 }
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[8], 'b')
=split [0] w.r.t (index of L)[8], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [0] is [-1]
==split[1] w.r.t (index of L)[8], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [1] is [-1]
==split[2] w.r.t (index of L)[8], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
```

```
\{67\}
{ 8 }
new split of [2] is [-1]
==split[3] w.r.t (index of L)[8], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
\{4 \quad 5 \}
\{67\}
{ 8 }
new split of [3] is [-1]
==split[4] w.r.t (index of L)[8], 'b')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [4] is [-1]
==split [6] w.r.t (index of L) [8], 'b')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [6] is [-1]
==split[8] w.r.t (index of L)[8], 'b')
```

```
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
\{4 \quad 5 \}
\{67\}
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0 1 2 3 4 5 6 7 8
0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1
Pick one [q] in L, Processing [q] = index of L = [8], 'a'
current all partitions (eq. classes) repr:
\{ 0 \ 1 \ 2 \ 3 \ 4 \ 6 \ 8 \}
current all partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
== for each [p], (split [p] w.r.t (index of L)[8], 'a')
==split [0] w.r.t (index of L)[8], 'a')
before split, partitions:
StateEqRel
\{0\}
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [0] is [-1]
```

```
==split[1] w.r.t (index of L)[8], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
\{67\}
{ 8 }
new split of [1] is [-1]
==split[2] w.r.t (index of L)[8], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [2] is [-1]
==split[3] w.r.t (index of L)[8], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
{ 4 5 }
{ 6 7 }
{ 8 }
new split of [3] is [-1]
==split [4] w.r.t (index of L) [8], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
```

```
\{2\}
{ 3 }
\{4 \quad 5 \}
\{67\}
{ 8 }
new split of [4] is [-1]
==split[6] w.r.t (index of L)[8], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
\{4 \quad 5\}
{ 6 7 }
{ 8 }
new split of [6] is [-1]
==split[8] w.r.t (index of L)[8], 'a')
before split, partitions:
StateEqRel
{ 0 }
{ 1 }
{ 2 }
{ 3 }
\{4 \quad 5 \}
{ 6 7 }
{ 8 }
new split of [8] is [-1]
--- while each [q], (split [p] w.r.t ([index of L]=[q],a))
L:
0\ \ 1\ \ 2\ \ 3\ \ 4\ \ 5\ \ 6\ \ 7\ \ 8
0 \;\; 0 \;\; 0 \;\; 0 \;\; 0 \;\; 0 \;\; 0 \;\; 0
```

References 27

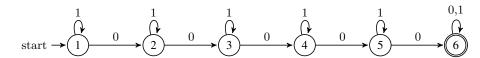


图 1.2: Minimizing example

{a,b},{d,e}is not equivalent states.

Sets of equivalent states:  $\{a,c\},\{b\},\{d\},\{e\}$ 

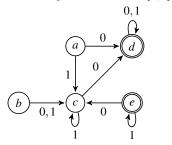


图 1.3: Finite state automaton

### 1.1 Minimization by equivalence of states

#### References

Hopcroft2008. John E. Hopcroft,Rajeev Motwani,Jeffrey D. Ullman 著, 孙家骕等译, 自动机理论、语言和计算机导论,Third Edition, 机械工业出版社,2008.7

WATSON93a. WATSON, B. W. A taxonomy of finite automata construction algorithms, Computing Science Note 93/43, Eindhoven University of Technology, The Netherlands, 1993. Available by ftp from ftp.win.tue.nl in pub/techreports/pi.

WATSON93b. WATSON, B. W. A taxonomy of finite automata minimization algorithms, Computing Science Note 93/44, Eindhoven University of Technology, The Netherlands, 1993. Available by ftp from ftp.win.tue.nl in pub/techreports/pi.

WATSON94a. WATSON, B. W. An introduction to the FIRE engine: A C++ toolkit for FInite automata and Regular Expressions, Computing Science Note 94/21, Eindhoven University of Technology, The Netherlands, 1994. Available by ftp from ftp.win.tue.nl in pub/techreports/pi

WATSON94b. WATSON, B.W. The design. and implementation of the FIRE engine: A C++ toolkit for FInite automata and Regular Expressions, Computing Science Note 94/22, Eindhoven University of Technology, The Netherlands, 1994. Available by ftp from ftp.win.tue.nl in pub/techreports/pi.

Christo<br/>3 G. Cassandras and Stéphane Lafortune, Introduction to Discrete Event Systems,<br/>Second Edition, New York, Springer, 2007

Wonham 2018. W. M. Wonham and Kai Cai, Supervisory Control of Discrete-Event Systems, Revised 2018.01.01

Jean<br/>2018. Jean-Éric Pin, Mathematical Foundations of Automata Theory,<br/>Version of June 15,2018

蒋宗礼 2013. 蒋宗礼, 姜守旭, 形式语言与自动机理论(第 3 版), 清华大学出版社,2013.05

Lipschutz 2007. S. Lipschutz and M. L. Lipson, Schaum's Outline of Theory and Problems of Discrete Mathematics, Third Edition, New York: McGraw-Hill, 2007.

Rosen 2007. K. H. Rosen, Discrete Mathematics and Its Applications, Seventh Edition, New York: McGraw-Hill, 2007.

R.Su and Wonham2004. R. Su and W. M. Wonham, Supervisor reduction for discrete-event systems, Discrete Event Dyn. Syst., vol. 14, no. 1, pp. 31–53, Jan. 2004.

Hopcroft71. Hopcroft, J.E. An n log n algorithm for minimizing states in a finite automaton, in The Theory of Machines and Computations (Z. Kohavi, ed.), pp.180-196, Academic Press, New York, 1971.

- Gries 73. Gries D. Describing an Algorithm by Hopcroft, Acta Inf. 2:97 109, 173. © by Springer-Verlag 1973
- Knuutila 2001. Knuutila, T. Re-describing an Algorithm by Hopcroft. Theoret. Computer Science 250 (2001) 333–363.
- Ratnesh<br/>95. Ratnesh Kumar, Modeling and Control of Logical Discrete Event Systems, © 1995 by Springer Science+Business Media New York.
- Jean<br/>2011. Jean Berstel, Luc Boasson, Olivier Carton, Isabelle Fagnot <br/>, *Minimization of automata*, Université Paris-Est Marne-la-Vallée 2010 Mathematics Subject Classification: 68Q45, 2011.
- Kenneth 2012. Kenneth H. Rosen 著, 徐六通译, 离散数学及其应用 Discrete Mathematics and Its Applications, seventh Edition, 2012, 机械工业出版社, 北京, 2014.