M1 Computers Science Parallelism



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# Lab work - n°3 Processus and monitor synchronization

#### **Documentation**

The concept of "monitor" can be used to synchronize processes using conditions (in the **multiprocessing** module). See the documentation available in Moodle.

### Exercice 1 - Reader-Writers - FIFO version

We consider parallel activities (processes) that simulate readers and writers having read or write access to a common file. Reads can be done in parallel but writes can only be done in mutual exclusion.

The behaviors of the processes are as follows:

```
A Reader

Loop on {
Loop on {
...
start_read()
Read the shared file;
end_read()
...
}

A Writer

Loop on {
...
Modify the shared file;
end_write()
...
}
```

Ensuring monitor synchronization, write the **start\_\*** and **end\_\*** operations so that processes run in order of arrival, and readers arriving before the first waiting writer run in parallel.

To do this, implement an **ExtendedCondition** class (using classic **Condition**s internally) that offers the following capabilities:

- Ability to handle high priority (i.e. **wait(0)**) in addition to normal priority (i.e. **wait(1)** or **wait()**). We limit ourselves here to the case of conditions with 2 priority levels.
- Possibility to check if the waiting list is empty (i.e. .empty())

Remark: Start from the code skeleton provided in the file tp2\_lectred\_base.py

Remember that in Python a function can take a default value for its arguments:

## def wait(priority = 1):

If **wait** is called with no argument then **priority** will be 1, otherwise **priority** will take the value of the argument.

## [Code to upload on moodle]

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## Exercice 2 - Management of access to voting booths

We want to simulate the behavior of NBV voters sharing access to NBP polling booths (with NBV much higher than NBP). The behavior of a voter consists in arriving at the polling station, entering a polling booth, preparing his envelope and coming out of the booth to place it in the ballot box.

The constraints are as follows:

- A voting booth can only be used by one voter at a time.
- Some voters have priority access to these voting booths (for example, disabled access). This allows them, in the event of a surge, to move ahead of "non-priority" voters.
- 1. Propose a specification for a monitor that manages competing accesses of NBV voters to these NBP booths.
- 2. Propose blocking and unblocking conditions as well as shared variables and conditions associated with the monitor.
- 3. Implement this monitor using the **ExtendedCondition** class defined in the previous exercise.
- 4. Write an application in which NBE voter processes use the operations of this monitor to synchronize their accesses to existing NBI booths. We will consider that 1/ratio voters have priority (for example by considering that voters for which id\_voter % ratio == 0 have priority)

## [Code to upload on moodle]

**Reminder**: If the displays are too fast, it is possible to delay the execution of a process for a few microseconds or nanoseconds using the primitives:

time.sleep(secondes)

See the online manual for their use:

https://docs.python.org/fr/3/library/time.html#time.sleep

We can use a randomly generated value (see the functions random.rand and random.seed) to vary the waiting times from one process to another.

https://docs.python.org/3/library/random.html

But, **be careful**, the timeout is not there to solve the problems of concurrent access to shared variables. In other words: any execution of a parallel application must give a consistent result without timeout!