

About the reading group

- 4~5 meetings the rest of the semester
- Topics:
 - Obviously strategy-proof mechanisms (Alon, next Wednesday)
 - Gali's series of works (Auction and regret quantal response)
 - level-k reasoning
 - Endowment Effect
 - Planning
- We are open to any suggestions. Whoever wants to lead a session is very invited (and can bring her/his own topic).

Expectation-based Reference-dependent Utility

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Outline

- Paradoxes under EU model
- Reference-dependent models
- **Expectations-based models and solution concepts**
- Applications
 - **Overbid in Auction**

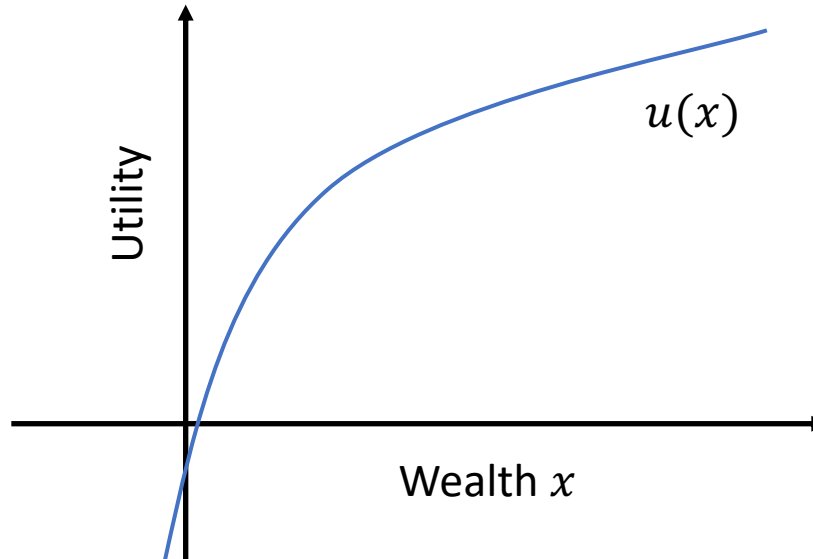
Preliminaries

- Payoff: a lottery $L = (x_1, p_1; x_2, p_2; \dots; x_n, p_n)$
- A decision problem: choice one out of a set of lotteries

Preliminaries

- Payoff: a lottery $L = (x_1, p_1; x_2, p_2; \dots; x_n, p_n)$
- A decision problem: choice one out of a set of lotteries
- Expected Utility model (EU):
 - $U(L) = \sum_{i=1}^n p_i u(x_i)$
 - Completeness, continuity, transitivity and **independence** axioms

Risk-averse:
 $u(x)$ is concave



Paradoxes under Expected Utility Model

1. Small-stake risk aversion implies absurd risk aversion at large stake
 - Turn down (\$110, 0.5; -\$100, 0.5) at all wealth levels \Rightarrow Turn down ($+\$ \infty$, 0.5; \$-1000, 0.5) Rabin (2000)
2. Violation of independence axiom
 - Axiom: $L \succeq L' \Leftrightarrow \alpha L + (1 - \alpha)L'' \succeq \alpha L' + (1 - \alpha)L'', \forall L, L', L'', \alpha \in (0,1)$
 - Samuelson (1963) anecdote:
 - A colleague will turn down $L = (\$200, 0.5; -\$100, 0.5)$ but will accept 100 independently draw L .
 - Allais' paradoxes, Allais (1954)

Reference-dependent Models

- Prospect Theory (Kahneman and Tversky, 1979)
 - $U(L|r) = \sum_{i=1}^n p_i v(x_i - r)$, solving Paradox 1
 - Probability weighing $U(L|r) = \sum_{i=1}^n w(p_i) v(x_i - r)$, solving both paradoxes

Exogenous reference: “status quo”

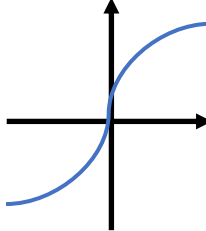
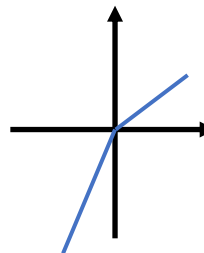
- Expectations-based reference-dependent model (EBRD, Koszegi and Rabin 2006, 2007, 2009)

Endogenous reference: “expectation”

Reference-dependent Models

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- Expectations-based reference-dependent model (EBRD, Koszegi and Rabin 2006, 2007, 2009)

$$\bullet \quad v(x_i - r) = \begin{cases} v(0) = 0 \\ \text{Diminishing sensitivity: } v''(x) \cdot x < 0 \\ \text{Loss aversion: } v(x) < -v(-x); v'(x) < v'(-x), \forall x > 0 \end{cases}$$

Why use expectation as reference point

- When facing a lottery choice problem, one forms a prior expectation of its outcome, and then after the uncertainty is resolved one experiences elation or disappointment (“Disappointment aversion”, Loomes and Sugden, 1986).
- Impose some discipline on models of reference dependence.

EBRD Framework

(This slide: given the reference, next slide: how to set the reference)

- Reference point: $R = (r_1, q_1; r_2, q_2; \dots; r_m, q_m)$
- Utility:

$$U(L|R) = EU(L) + V(L|R)$$

Monetary/Intrinsic/
Instrumental utility

Gain-loss utility

- Gain-loss utility: $V(L|R) = \sum_{i=1}^n p_i v(x_i|R)$
 - **KR approach:** $v(x_i|R) = \sum_{j=1}^m q_j v(u(x_i) - u(r_j))$
 - DA approach: $v(x_i|R) = v(u(x_i) - EU(R))$

Solution Concepts

- Given a lottery set $S = \{L, L', L'', \dots\}$, which one will be chosen?

- Choice-acclimating Personal Equilibrium (CPE, KR 2007):

$$L^* = \operatorname{argmax}_{L \in S} \{U(L|L)\}$$

Most tractable one!

- Personal Equilibrium (PE, KR 2006)

$$L^* \in \{L \in S \mid U(L|L) \geq U(L'|L), \forall L' \in S\}$$

- Preferred Personal Equilibrium (PPE, KR 2007)

$$L^* = \operatorname{argmax}_{L \in \operatorname{PE}(S)} \{U(L|L)\}$$

CPE

- $v(x) = \begin{cases} \eta x, & x \geq 0 \\ -\eta\lambda x, & x < 0 \end{cases} \quad \lambda > 1 \text{ (Loss aversion)}$

- Maximize

$$\begin{aligned} U(L|L) &= \sum_{i=1}^n p_i u_i + \sum_{i=1}^n \sum_{j=1}^n p_i p_j v(u_i - u_j) \\ &= \sum_{i=1}^n p_i u_i - \sum_{i=1}^n \sum_{j=i+1}^n p_i p_j \Lambda (u_j - u_i) \end{aligned}$$

$$\Lambda = \eta(\lambda - 1)$$

- For two outcomes: $U(L|L) = EU - \Lambda p_1(1 - p_1)(u_2 - u_1)$

CPE

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$$\text{Var} = \sum_{i=1}^n \sum_{j=1}^n p_i p_j \frac{1}{2} (u_j - u_i)^2$$

- [Solved] Samuelson (1963) anecdote “A colleague will turn down $L = (\$200, 0.5; -\$100, 0.5)$ but will accept 100 independently draw L .”

CPE

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CPE and first-order stochastic dominance

- First-order stochastic dominance (FOSD):
 - $L' = (\$10, 0.9, \$11, 0.1) \succeq L = (\$10, 1)$
- CPE violates FOSD, but EU, PE and PPE not.
- Experimental result: people violate FOSD when it is not obvious.

Applications

- Overbid in auction (Theory: Lange and Ratan 2010; Experiments: Banerji and Gupta)

First-price auction

- One item, n agents
- Intrinsic valuation: CDF $H: [\underline{w}, \bar{w}] \rightarrow [0,1]$
- Gain-loss utility function:
 - Two dimensions: $v_k(x) = \begin{cases} 0, & x \geq 0 \\ -\lambda_k x, & x < 0 \end{cases}, k = 0 \text{ for money, } 1 \text{ for item}$
 - $v = v_0 + v_1$
- Solution concept: Symmetric increasing Bayesian Nash equilibrium under EBRU

First-price auction

- $U(L|L)$:

$$\Pi^{1,CA}(b^i, w^i) = \underbrace{f(b^i)(w^i - b^i)}_{\text{Intrinsic utility}} - \underbrace{f(b^i)(1 - f(b^i))[\lambda_0 b^i + \lambda_1 w^i]}_{\text{Gain-loss utility}}$$

$f(b^i)$ -- Winning probability under bid b^i

Solving the equilibrium

- First-order derivative equation:

$$f'(b^i)[w^i - b^i - (1 - 2f^i)(\lambda_0 b^i + \lambda_1 w^i)] = f^i[1 + (1 - f^i)\lambda_0]$$

-variable $b(w^i)$

$$-f(b(w^i)) = H^{n-1}(b^{-1}(b(w^i)))$$

- Meanwhile, at the equilibrium, $b^{-1}(b(w^i)) = w^i$
- Therefore, $f'(b^i) = (H^{n-1}(w^i))' / b'(w^i)$

$$(H^{n-1})'(w^i) w^i [1 - (1 - 2H^{n-1}(w^i))\lambda_1] = [H^{n-1}(w^i)(1 + (1 - H^{n-1}(w^i))\lambda_0)b(w^i)]'$$

Equilibrium Bid

Proposition 1 (Commodity first-price auction – CA). The unique monotonically increasing symmetric Bayesian Nash equilibrium bidding function for commodity auctions is given by

$$b^{1,CA}(w) = \begin{cases} w \frac{1-\lambda_1(1-H^{n-1}(w))}{1+\lambda_0(1-H^{n-1}(w))} - \frac{\int_{w_L}^w H^{n-1}(z)[1-\lambda_1(1-H^{n-1}(z))]dz}{H^{n-1}(w)[1+\lambda_0(1-H^{n-1}(w))]} & \text{if } w \geq w_L^{CA} \\ 0 & \text{if } w < w_L^{CA} \end{cases} \quad (7)$$

- For large w , an agent overbids w.r.t. BNE
- For small w (some w with winning probability < 0.5), an agent underbids w.r.t. BNE.
- Intuition: People tend to reduce the uncertainty.

Equilibrium Bid (induced-value)

- One-dimension gain-loss utility v

Corollary 1 (*Induced-value first-price auction – IV*). The unique continuous monotonically increasing symmetric Bayesian Nash equilibrium bidding function for induced-value auctions is given by

$$b^{1,IV}(w) = \begin{cases} w - \frac{\int_{w_L^{IV}}^w H^{n-1}(z)[1-\lambda_0(1-H^{n-1}(z))]dz}{H^{n-1}(w)[1-\lambda_0(1-H^{n-1}(w))]} & \text{if } w \geq w_L^{IV} \\ w & \text{if } w < w_L^{IV} \end{cases} \quad (9)$$

- Induced-value, lab experiments
- People should always overbid w.r.t. BNE

Second-price auction

- Commodity auction:
 - For small w (some w with winning probability < 0.5), an agent underbids w.r.t. BNE.
- Induced-value auction:
 - Bid truthfully (=BNE)
- Support by (Banerji and Gupta, 2014)
- Might contradict with (Noti et al 2014): Position auction, induced-value auction, small w agent overbids.

Applications

- Endowment Effects
- Labor Supply (Farber, 2005)
- ...

Q&A