

1.) Prove:

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

• assumption 1: $A \cup (B \cap C)$

- if $x \in A$, then $x \in A \cup B$ and $x \in A \cup C$
 $\equiv x \in (A \cup B) \cap (A \cup C)$

- if $x \in B \cap C$, then $x \in B$, thus $x \in A \cup B$

- if $x \in B \cap C$, then $x \in C$, thus $x \in A \cup C$
 $\equiv x \in (A \cup B) \cap (A \cup C)$

Based on such statements, we can draw out:

if $A \cup (B \cap C)$, then $x \in (A \cup B) \cap (A \cup C)$

• assumption 2: $(A \cup B) \cap (A \cup C)$

- if $x \in A$, then $x \in A \cup B$ and $x \in A \cup C$,
thus we can conclude $x \in A \cup (B \cap C)$

- if $x \notin A$, then $x \in B$ and $x \in C$, thus
 $x \in B \cap C$, meaning we can introduce
 $x \in A \cup (B \cap C)$.

Based on such statements, we can draw out:

if $(A \cup B) \cap (A \cup C)$, then $x \notin A \cup (B \cap C)$

⇒ Combine each conclusion from both assumption 1 and assumption 2:

if $A \cup (B \cap C)$, then $x \in (A \cup B) \cap (A \cup C)$

+

if $(A \cup B) \cap (A \cup C)$, then $x \notin A \cup (B \cap C)$

so we get

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

2.) Suppose $A \subseteq B$ and $C \subseteq D$,

Show $A \times C \subseteq B \times D$

Cartesian product:

$$A \times B = \{(x, y) \mid x \in A \wedge y \in B\}$$

meaning for any $(x, y) \in A \times C =$
 $x \in A \wedge y \in C$

Since $A \subseteq B$ and $C \subseteq D$, and $x \in A$, so $x \in B$

Since $A \subseteq B$ and $C \subseteq D$, and $y \in C$, so $y \in D$

Meaning, if $(x, y) \in A \times C$, then $(x, y) \in B \times D$

thus, $A \times C \subseteq B \times D$ given $A \subseteq B$ and $C \subseteq D$

3.) Prove that $A - (B \cap C) = (A - B) \cup (A - C)$

• assumption 1: $A - (B \cap C)$

- if $x \in A - (B \cap C)$, then $x \in A \wedge x \notin B \cap C$
meaning $x \in A \wedge (x \notin B \vee x \notin C)$.

Therefore, $\Rightarrow x \in A$ and $x \notin B$, so $x \in A - B$
 $x \in A$ and $x \notin C$, so $x \in A - C$

when put together, $x \in (A - B) \cup (A - C)$

in conclusion, if $A - (B \cap C)$, then $(A - B) \cup (A - C)$

• assumption 2: $(A - B) \cup (A - C)$

- if $x \in (A - B) \cup (A - C)$, then it means:

if $x \in (A - B)$ \vee if $x \in (A - C)$

$x \in A$, but $x \notin B$

$x \in A$, but $x \notin C$

\Rightarrow $x \in A$, but $x \notin (B \text{ and } C)$

$\equiv x \in A$ and $x \notin (B \cap C)$

thus $x \in A - (B \cap C)$

in conclusion, if $(A - B) \cup (A - C)$,
then $A - (B \cap C)$

put conclusion from assumption 1 and 2 together, we have

if $A - (B \cap C)$, then $(A - B) \cup (A - C)$
+

in conclusion, if $(A - B) \cup (A - C)$,
then $A - (B \cap C)$

means that $A - (B \cap C) = (A - B) \cup (A - C)$

$$A - (B \cap C) = (A - B) \cup (A - C)$$

4.) An ordered pair (a, b) can be defined as
the set $\{\{a\}, \{a, b\}\}$.

Show that $(a, b) = (c, d)$ iff
 $a = c$ and $b = d$

Based on the definition of ordered pair,
 $\text{Pair}(a, b)$ makes the set $\{\{a\}, \{a, b\}\}$.
 $\text{Pair}(c, d)$ makes the set $\{\{c\}, \{c, d\}\}$.

So if $(a, b) = (c, d)$, then

$$\{\{a\}, \{a, b\}\} = \{\{c\}, \{c, d\}\}$$

Given (a, d) ,

\Rightarrow if $a = b$, then it makes $\{ \{b\}, \{b, d\} \}$

\Rightarrow if $a = d$, then it makes $\{ \{d\}, \{b, d\} \}$

\Rightarrow if $a = c$, then it makes $\{ \{c\}, \{c, d\} \}$

meaning, ONLY if $a=c, \{ \{c\}, \{c, d\} \}$

given (c, b) ,

\Rightarrow if $b = a$, then it makes $\{ \{c\}, \{a\} \}$

\Rightarrow if $b = c$, then it makes $\{ \{c\}, \{c\} \}$

\Rightarrow if $b = d$, then it makes $\{ \{c\}, \{c, d\} \}$

meaning, ONLY if $b=d, \{ \{c\}, \{c, d\} \}$

Since ONLY if $a=c, \{ \{c\}, \{c, d\} \}$

+

ONLY if $b=d, \{ \{c\}, \{c, d\} \}$

\Rightarrow if and only if $a=c$ and $b=d, \{ \{c\}, \{c, d\} \}$

5.) Given an arbitrary relation R ,
suppose we compute two relations:

- R_1 , the reflexive closure of the transitive closure of R
- R_2 , the transitive closure of the reflexive closure of R

Prove or disprove: $R_1 = R_2$ for all R

Since both reflexive closure and transitive closure adds minimum amount of pairs required to satisfy its respective requirements to the set by definition, doing a closure of a closure on the same set R would produce the same result no matter of which closure occurs first.

For example,

assuming

$$R = \{(1,2), (2,3), (3,4)\},$$

- the transitive closure of R is

$$R^t = \{(1,2)(2,3)(3,4), (1,3), (2,4), (1,4)\}$$

thus R_1 (reflexive closure of R^t) is

$$R_1 = \{(1,1)(2,2), (3,3), (4,4), (1,2), (2,3), (3,4)\}$$

$\{(1,3), (2,4), (1,4)\}$

- the reflexive closure of R is

$R^R = \{(1,1), (2,2), (3,3), (4,4), (1,2), (2,3), (3,4)\}$

thus R_2 (transitive closure of R^R) is

$R_2 = \{(1,1), (2,2), (3,3), (4,4), (1,2), (2,3), (3,4), (1,3), (2,4), (1,4)\}$

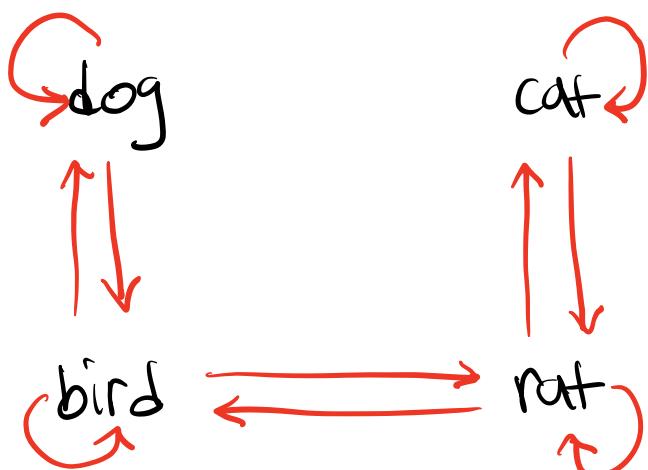
thus $R_1 = R_2$ for all R

6.) Let $A = \{\text{cat, dog, bird, rat}\}$

and R be a relation on A defined by

$\{(x,y) \mid x \text{ and } y \text{ have at least one letter in common}\}$

a.)



b.)

$$R = \{ (\underline{\text{cat}}, \underline{\text{cat}}), (\underline{\text{dog}}, \underline{\text{dog}}), (\underline{\text{bird}}, \underline{\text{bird}}), \\ (\underline{\text{rat}}, \underline{\text{rat}}), (\underline{\text{cat}}, \underline{\text{rat}}), (\underline{\text{rat}}, \underline{\text{cat}}), \\ (\underline{\text{dog}}, \underline{\text{bird}}), (\underline{\text{bird}}, \underline{\text{dog}}), (\underline{\text{bird}}, \underline{\text{rat}}) \\ (\underline{\text{rat}}, \underline{\text{bird}}) \}$$

- Reflexive: since $\forall x (x, x) \in R$, R is reflexive.
- Symmetric: since $(x, y) \in R \rightarrow (y, x) \in R$,
 R is symmetric.
- Transitive:
by definition $(a, b) \in R \wedge (b, c) \in R \rightarrow$
 $(a, c) \in R$

since $(\text{cat}, \text{rat}) \wedge (\text{rat}, \text{bird}) \in R$
but $(\text{cat}, \text{bird}) \notin R$,
 R is not transitive.

7.) Given a relation R on set A ,
 Prove that if R is transitive,
 so is R^{-1}

- Definition of inverse relations:

If R is from x to y , then R^{-1}
 is from y to x , and

$$R^{-1} = \{(y, x) : (x, y) \in R\}$$

So, assume $x, y, z \in A$ and
 $(x, y) \in R^{-1} \wedge (y, z) \in R^{-1}$.

By definition of inverse relations,

$$(x, y) \in R^{-1} \rightarrow (y, x) \in R$$

$$(y, z) \in R^{-1} \rightarrow (z, y) \in R$$

Given $(y, x) \in R$ and $(z, y) \in R$,
 and R is transitive, $(z, x) \in R$

Since we have $(z, x) \in R$
 and by definition of inverse relations,
 we have $(x, z) \in R^{-1}$.
 meaning,

$$\underline{(z,x) \in R} \longrightarrow (x,z) \in R^{-1}$$

Since we know

$$(x,y) \in R^{-1} \text{ and } (y,z) \in R^{-1}$$

from assumption and we have

$$(x,z) \in R^{-1},$$

by definition of transitivity,

$$(x,y) \in R^{-1} \wedge (y,z) \in R^{-1} \rightarrow (x,z) \in R^{-2}$$

R^{-1} is transitive.

therefore, R^{-1} is transitive if R is transitive.



8.) Suppose R and S are symmetric relations on a set A . Prove that $R \circ S$ is symmetric iff $R \circ S = S \circ R$

assume: $a, b, c \in A$ and $R \circ S = S \circ R$

- if $(a,c) \in R \circ S$, then
 $\exists b (a,b) \in R \quad (b,c) \in S$

Knowing R and S are symmetric, we know

$$(c, b) \in S \text{ and } (b, a) \in R$$

meaning,

$$(c, a) \in S \circ R = (c, a) R \circ S$$

and since we have $(c, a) R \circ S$

$$\underline{(a, c) \in R \circ S} \rightarrow (c, a) R \circ S$$

which satisfies $R \circ S$ symmetric requirement.

- Knowing $R \circ S$ is symmetric,

If we assume

$$(c, a) \in R \circ S \leftrightarrow (a, c) \in R \circ S$$

then

$$(a, c) \in R \circ S \leftrightarrow (c, a) \in S \circ R$$

which concludes

$$R \circ S \subseteq S \circ R$$

on the other hand, knowing $R \circ S$ is symmetric

if we assume

$$(a, c) \in S \circ R \leftrightarrow (c, a) \in R \circ S$$

then

$$(c, a) \in R \circ S \leftrightarrow (a, c) \in R \circ S$$

which concludes

$$S \circ R \subseteq R \circ S$$

Since we have

$R \circ S \subseteq S \circ R$ and $S \circ R \subseteq R \circ S$,
we now know

$$R \circ S = S \circ R$$

q.)

Symmetry : if $(a, b) \in R \rightarrow (b, a) \in R$

Anti-Symmetric : if $(a, b) \in R \wedge (b, a) \in R \rightarrow a = b$

if we have some $R = \{(x, y) \mid x = y\}$,

since we know $x = y$,

if $(x, y) \in R$, then $(y, x) \in R$.

So by definition, R is symmetric since

$$\underline{(x, y) \in R} \rightarrow \underline{(y, x) \in R}$$

at the same time,

given $x = y$ means $\underline{(x, y) \in R} \wedge \underline{(y, x) \in R}$

R is antisymmetric as

$$(x, y) \in R \wedge (y, x) \in R \rightarrow x = y$$

10.) Find all equivalence relations on $\{1, 2, 3\}$

given $A = \{1, 2, 3\}$

and its cartesian product ($A \times A$) to make R ,

$$R = \{(1, 1), (2, 2), (3, 3), (1, 2), (2, 1), (2, 3), (3, 2), (1, 3), (3, 1)\}$$

equivalence relation requires the relation to be reflexive, symmetric, and transitive.

$$R_1 = \{(1, 1), (2, 2), (3, 3), (1, 2), (2, 1), (1, 3), (3, 1), (2, 3), (3, 2)\}$$

satisfies equivalence ✓

$$R_2 = \{(1, 1), (2, 2), (3, 3), (2, 3), (3, 2)\}$$

satisfies equivalence ✓

$$R_3 = \{(1, 1), (2, 2), (3, 3), (1, 3), (3, 1)\}$$

satisfies equivalence ✓

$$R_4 = \{(1, 1), (2, 2), (3, 3), (1, 2), (2, 1)\}$$

satisfies equivalence ✓

$$R_5 = \{(1, 1), (2, 2), (3, 3)\}$$

satisfies equivalence ✓

all other combinations of sets on R does not satisfy equivalence, thus there are 5 equivalence relations.