

**ESTRUCTURA DE DATOS 1**  
Código ST0245

## Laboratory practice 1:

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Member	Date	Done	Doing	Todo list
tomas calle	23/02/2021	codingBat 2	Make the document, and organize everything so we can start to make the final paper.	Volume algorithm or 1.2
Juan Camilo salazar	23/02/2021	3 exercices of codingBat 2	finish recursion 1 of codingBat	make changes to the Genome algo.
tomas calle	25/02/2021	Volume 2 algo Finished	helping to debug the Genome algo	Start and finish the test Practice.
Juan Camilo Salazar	25/02/2021	.5 Test Practice	.5 test practice	make changes to the Genome algo

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## ESTRUCTURA DE DATOS 1

### Código ST0245

#### 1.1

```
public static int genoma(String cadena, String subCadena){
    if (cadena.length() == 0 || subCadena.length() == 0) {
        return 0;
    } else if (cadena.charAt(0) == subCadena.charAt(0)) {
        return 1 + genoma(cadena.substring(1), subCadena.substring(1));
    }
    int llamado1 = genoma(cadena.substring(1), subCadena);
    int llamado2 = genoma(cadena, subCadena.substring(1));
    if (llamado1 > llamado2) {
        return llamado1;
    } else {
        return llamado2;
    }
}
```

#### 1.2

```
public static int fit1x2(int nrec ,int tiles){
    if(nrec*2 == tiles){
        return 1;
    } else if(nrec*2 < tiles){
        return 0;
    }else{
        return fit1x2(nrec, tiles + 2) + fit1x2(nrec, tiles + 4);
    }
}
```

#### 2.1

```
1) public boolean array220(int[] nums, int index) {

    if(index>=nums.length-1){
        return false;
    }else if(nums[index]*10==nums[index+1]){
        return true;
    }
    else return array220(nums,index+1); // F(n-1)
}

O(n-1) n being the length of the static array.
```

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```
2) public boolean nestParen(String str) {

    if(str.length()==0){
        return true;
    }
    else if(str.length()==1) {
        return false;
    }
    else if((str.charAt(0)!='(' && str.charAt(str.length()-1)==')')
    || (str.charAt(0)=='(' && str.charAt(str.length()-1)!='(')
    || (str.charAt(0)!='(' && str.charAt(str.length()-1)!='('))){
        return false;
    }
    return nestParen(str.substring(1,str.length()-1));
}

O(n-2) = O(n) n being the length of the String
```

```
3) public String parenBit(String str) {

    if(str.length()==0) return "";
    else if(str.charAt(0) == '(' && str.charAt(str.length()-1)==')') {
        return str;
    } else if(str.charAt(0) == '(' && str.charAt(str.length()-1)!='') {
        return parenBit(str.substring(0,str.length()-1));
    } else if(str.charAt(0) != '(' && str.charAt(str.length()-1)==')') {
        return parenBit(str.substring(1));
    }
    else{
        return parenBit(str.substring(1,str.length()-1));
    }
}

O(n-2) = O(n) n being the length of the String
```

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```
4) public int countAbc(String str) {

    if(str.length()<3){
        return 0;
    }

    if(str.substring(0,3).equals("abc")||str.substring(0,3).equals("aba")){
        return 1 + countAbc(str.substring(1));
    }else return countAbc(str.substring(1));
}

O(n-2) = O(n) n being the length of the String
```

```
5) public int strCount(String str, String sub) {

    if(str.length()==0) return 0;
    else if(str.length()<sub.length()) return 0;
    else if(str.substring(0,sub.length()).equals(sub)){
        return 1+strCount(str.substring(sub.length()),sub);
    }else return strCount(str.substring(1),sub);
}

O(n-2) = O(n) n being the length of the String
```

## 2.2

```
public boolean groupSum6(int start, int[] nums, int target) {

    if(start>=nums.length){ // 3C_1
        return target == 0; // 2C_2
    }else if(nums[start] == 6){ // 4C_3
        return groupSum6(start+1,nums,target-6); // 4C_4 + F(n-1)
    }

    return groupSum6(start+1,nums,target-nums[start])||groupSum6(start+1,nums,target);
    // 7C_5 + F(n-1) + F(n-1)
}

// F(n) = c_1 + 2^(n - 1) (where c_1 is an arbitrary parameter)
// O(2^(n))

n = the number of elements in the array nums.

public boolean groupSum(int index,int[] nums,int target){

    if(index == nums.length){ // 3C_1
        return target == 0; // // 2C_2
```

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```

    }else{
        return groupSum(index + 1,nums,target-nums[index]) || groupSum(index + 1,nums,target);
        // 7C_3 + F(n-1) + F(n-1)

        // F(n) = c_1 2^(n - 1) (where c_1 is an arbitrary parameter)
        // O(2^(n)) donde n es la longitud del arreglo

        n = the number of elements in the array nums.
    }
}

public static boolean groupNoAdj(int index, int[] nums, int target){
    if(index >= nums.length){ // 3C_1
        return target == 0; // 3C_2
    }

    return groupNoAdj(index + 2,nums,target-nums[index]) || groupNoAdj(index + 1,nums,target);
    // 7C_3 + F(n-1) + F(n-1)
    // F(n) = c_1 2^(n - 1) (where c_1 is an arbitrary parameter)
    O(2^(n))

    n = the number of elements in the array nums.
}

public static boolean groupSum5(int start, int[] nums, int target){

    if(start >= nums.length){
        return target == 0;
    }

    if(nums[start]%5 == 0){
        if(start < nums.length-1){
            if(nums[start+1] == 1){
                return groupSum5(start+2,nums,target-nums[start]);
            }
            return groupSum5(start+1,nums,target-nums[start]);
        }
    }

    return groupSum5(start+1,nums,target-nums[start]) || groupSum5(start+1,nums,target);

    O(2^(n))

```

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```

n = the number of elements in the array nums.

}

public static boolean groupSumClump(int start, int[] nums, int target) {
    if(start >= nums.length){ // 3C_1
        return target == 0; // 2C_1
    }
    int countadorValor = 0; // C_2
    int countadorIndex = 0; // C_3
    boolean seguidos = true; // C_4

    for(int tt = start; tt < nums.length; tt++){ // 2C_5 * F(n-1)
        seguidos = (nums[tt] == nums[start]);

        if(seguidos){
            countadorValor = countadorValor + nums[tt];
            countadorIndex++;
        }
    }
    return (
        groupSumClump(start+countadorIndex, nums, target-countadorValor) F(n-1)
        ||
        groupSumClump(start+countadorIndex, nums, target); F(n-1)
    ) // F(n) = F(n-1) + F(n-1) + F(n-1)  $\longleftrightarrow$   $O(3^n)$ 

    n = the number of elements in the array nums.

```

Input:

$$f(n) = f(n-1) + f(n-1) + f(n-1)$$

Result:

$$f(n) = 3 f(n-1)$$

Recurrence equation solution:

$$f(n) = c_1 3^{n-1} \text{ (where } c_1 \text{ is an arbitrary parameter)}$$

### 2 3.1. Calculen la complejidad, para el peor de los casos, del ejercicio 1.1

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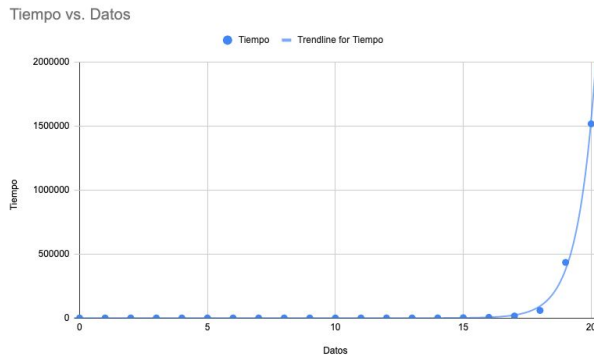
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- $O(2^n)$  al tener dos llamados recursivos

### 3.2. Analyzing the algorithm with 20 different cases.



Donde los datos es la cantidad de caracteres en cada string

### 3.3. ¿La complejidad del algoritmo del ejercicio 1.1 es apropiada para encontrar la subsecuencia común más larga entre ADNs mitocondriales como los de los datasets?

No es apropiada porque tiene complejidad de  $O(2^n)$  entonces se demorara varios minutos en analizar tanta información. Esto significa que el algoritmo es bastante ineficiente en términos de tiempo.

### 3.5. Calculate the complexity of the following exercises: 2.1, 2.2, 3.6 and explain in your own words the meaning of n or m depending on the exercise.

The analysis of the complexities are in the photos above. (2.1, 2.3, 3.6)

## 4. Simulacro de Parcial

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1. Completa la línea 6 .....

- ☐ a. s.substring(0, i)
- ☐ b. s.substring(0, n)
- ☐ c. s.substring(i, n)

2. Completa la línea 9 .....

- ☐ a. false
- ☐ b. s.substring(0, i)
- ☒ c. true

3. Completa la línea 11 .....

- ☒ a. solve(t, s.substring(i), n - i)
- ☐ b. solve(pfx, t), n - i)
- ☐ c. solve(t, s.substring(n), I - n)

1. Completa la línea 9 .....

Y Completa la línea 10 .....

2. Completa la línea 11 .....

Y Completa la línea 12 .....

3. ¿Cuál es la ecuación de recurrencia que representa la complejidad, en el tiempo, para el peor de los casos, en términos de  $p=n+m$ ?

$T(p) = \dots\dots\dots$

FloodFill Util(screen, x+1, y+1,prevC,newC,N,M);

FloodFill Util(screen, x-1, y-1,prevC,newC,N,M);

FloodFill Util(screen, x+1, y-1,prevC,newC,N,M);

FloodFill Util(screen, x-1, y+1,prevC,newC,N,M);

2-5 :  $C_1$

9-12 :  $C_2 + f(m-1)^2 \cdot f(n-1)^2$

12-16 :  $C_3 + f(m-1)^2 \cdot f(n-1)^2$

2-16 :  $C_1 + C_2 + C_3 + (2 \cdot F(m-1) + F(n-1)) + (2F(m-1) + 2 \cdot F(n-1))$

$= 4F(m-1) + 4F(n-1)$

$F(m) = 4 \cdot F(m-1)$

$F(m) = 4^m$

$F(n) = 4 \cdot F(n-1)$

$F(n) = 4^n$

$O(4^m + 4^n)$

La complejidad de la función *mystery* es

- ☐ a.  $T(n,m) = C \times n \times m$
- ☒ b.  $T(n,m) = C \times n \times m^2$
- ☐ c.  $T(n,m) = C \times n \times m \cdot \log m$
- ☐ d.  $T(n,m) = C \times n \times m^3$

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```
01 static boolean isPal(String s) {  
02     if(s.length() == 0 || s.length() == 1)  
03         return true;  
04     if ( s.charAt(0) == s.charAt(s.length()-1) )  
05         return isPal(s.substring(1, s.length()-1));  
06     //else  
07     return false;  
08 }
```

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