COMP108 Data Structures and Algorithms

Algorithm Efficiency (Part II - Exercises)

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Determine the order of growth of the following functions.

1.
$$n^3 + 3n^2 + 3$$





2.
$$4n^2 \log n + n^3 + 5n^2 + n$$

- 3. $2n^2 + n^2 \log n$
- 4. $6n^2 + 2^n$

Determine the order of growth of the following functions.

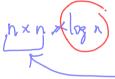
1.
$$n^3 + 3n^2 + 3$$

2.
$$4n^2 \log n + n^3 + 5n^2 + n$$

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4.
$$6n^2 + 2^n$$

 $O(n^3)$







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$$O(n^3)$$

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 $O(n^2 \log n)$

4.
$$6n^2 + 2^{n^2}$$

$$O(2^n)$$

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$$O(n^3)$$

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$$4n^2 \log n + n^3 + 5n^2 + n$$

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$$2n^2 + n^2 \log n$$

$$O(n^2 \log n)$$

4.
$$6n^2 + 2^n$$

$$O(2^n)$$

1.
$$n^2 \log n + (n^3 + 3n^2 + 3)$$

2.
$$n + 1000$$

3.
$$6n^{20} + 2^n$$

4.
$$n^3 + 5n^2 \log n + n$$

$$O(n^2 \log n)$$
? $\bigvee O(N^3)$

$$O(n^{20})$$
?

$$O(n^2 \log n)$$
?

Are the following correct?

1.
$$n^2 \log n + n^3 + 3n^2 + 3$$

 $O(n^2 \log n)$?

$$O(n^3)$$

2.
$$n + 1000$$

3.
$$6n^{20} + 2^n$$

$$O(n^{20})$$
?

4.
$$n^3 + 5n^2 \log n + n$$

$$O(n^2 \log n)$$
?

1.
$$n^2 \log n + n^3 + 3n^2 + 3$$

$$O(n^2 \log n)$$
?

$$O(n^3)$$

2.
$$n + 1000$$

3.
$$6n^{20} + 2^n$$

$$O(n^{20})$$
?

4.
$$n^3 + 5n^2 \log n + n$$

$$O(n^2 \log n)$$
?

1.
$$n^2 \log n + n^3 + 3n^2 + 3$$

$$O(n^2 \log n)$$
?

$$O(n^3)$$

2.
$$n + 1000$$

3.
$$6n^{20} + 2^n$$

$$O(n^{20})$$
?

$$O(2^n)$$

4.
$$n^3 + 5n^2 \log n + n$$

$$O(n^2 \log n)$$
?

1.
$$n^2 \log n + n^3 + 3n^2 + 3$$

$$O(n^2 \log n)$$
?

$$O(n^3)$$

2.
$$n + 1000$$

3.
$$6n^{20} + 2^n$$

$$O(n^{20})$$
?

$$O(2^n)$$

4.
$$n^3 + 5n^2 \log n + n$$

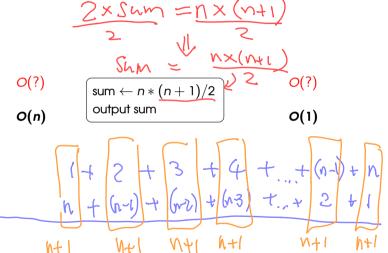
$$O(n^2 \log n)$$
?

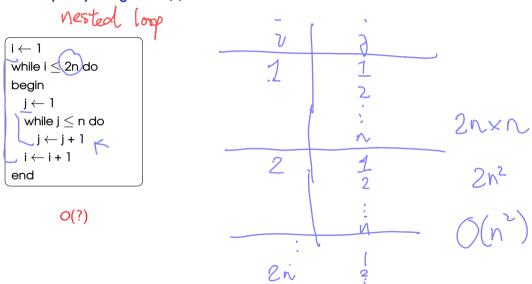
$$O(n^3)$$

```
\begin{cases} \operatorname{sum} \leftarrow 0, i \leftarrow 1 \\ \operatorname{while} i \leq n \text{ do} \\ \operatorname{begin} \\ \operatorname{sum} \leftarrow \operatorname{sum} + i \\ i \leftarrow i + 1 \\ \operatorname{end} \\ \operatorname{output} \operatorname{sum} \end{cases}
```

 $\begin{array}{c} \text{sum} \leftarrow 0, \text{i} \leftarrow 1 \\ \text{while i} \leq \text{n do} \\ \text{begin} \\ \text{sum} \leftarrow \text{sum + i} \\ \text{i} \leftarrow \text{i + 1} \\ \text{end} \\ \text{output sum} \end{array} \qquad \begin{array}{c} \text{O(?)} \\ \text{sum} \leftarrow n*(n+1)/2 \\ \text{output sum} \end{array}$

sum \leftarrow 0, i \leftarrow 1 while $i \le n$ do O(?)begin $sum \leftarrow sum + i$ O(n) $i \leftarrow i + 1$ end output sum

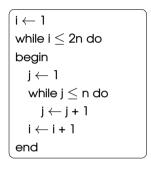




```
\begin{aligned} \mathbf{i} &\leftarrow \mathbf{1} \\ \text{while } \mathbf{i} &\leq 2\mathbf{n} \text{ do} \\ \text{begin} \\ \mathbf{j} &\leftarrow \mathbf{1} \\ \text{while } \mathbf{j} &\leq \mathbf{n} \text{ do} \\ \mathbf{j} &\leftarrow \mathbf{j} + \mathbf{1} \\ \mathbf{i} &\leftarrow \mathbf{i} + \mathbf{1} \\ \text{end} \end{aligned}
```

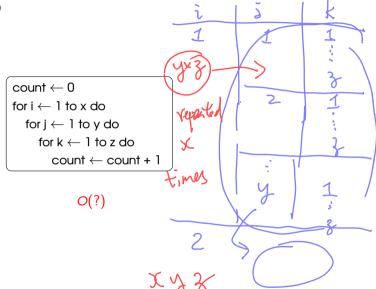
0(?

 $O(n^2)$



0(?)

 $O(n^2)$



```
i \leftarrow 1
while i < 2n do
                                            count \leftarrow 0
begin
                                            for i \leftarrow 1 to x do
  j ← 1
                                               for i \leftarrow 1 to y do
   while j \leq n do
                                                  for k \leftarrow 1 to z do
      j \leftarrow j + 1
                                                     count \leftarrow count + 1
   i \leftarrow i + 1
end
                                                            O(?)
            O(?)
                                                          O(xyz)
           O(n^2)
```

 $i \leftarrow 1$ while i < 2n do begin $j \leftarrow 1$ while $j \leq n$ do $j \leftarrow j + 1$ $i \leftarrow i + 1$ end

O(?)0(n2) hththt...th

while i < n docount \leftarrow 0 begin for $i \leftarrow 1$ to x do for $j \leftarrow 1$ to y do while $j \leq n$ do for $k \leftarrow 1$ to z do $j \leftarrow j + 1$ $count \leftarrow count + 1$ $i \leftarrow i + 1$ end O(?)O(xyz)

 $i \leftarrow 1$

```
i \leftarrow 1
                                                                                                  i \leftarrow 1
while i < 2n do
                                                                                                  while i \le n do
                                             count \leftarrow 0
begin
                                                                                                  begin
                                             for i \leftarrow 1 to x do
  j ← 1
                                                                                                     j \leftarrow i
                                                for i \leftarrow 1 to y do
  while j \leq n do
                                                                                                     while j \leq n do
                                                    for k \leftarrow 1 to z do
      j \leftarrow j + 1
                                                                                                        j \leftarrow j + 1
                                                       count \leftarrow count + 1
   i \leftarrow i + 1
                                                                                                     i \leftarrow i + 1
end
                                                                                                  end
                                                             O(?)
            O(?)
                                                                                                               O(?)
                                                           O(xyz)
           O(n^2)
                                                                                                              O(n^2)
```

```
\begin{split} \mathbf{i} \leftarrow \mathbf{1} \\ \text{count} \leftarrow \mathbf{0} \\ \text{while } \mathbf{i} < \mathbf{n} \text{ do} \\ \text{begin} \\ \mathbf{i} \leftarrow \mathbf{i} * \mathbf{2} \\ \text{count} \leftarrow \text{count} + \mathbf{1} \\ \text{end} \\ \text{output count} \end{split}
```

O(?)

$$\begin{split} & i \leftarrow 1 \\ & \text{count} \leftarrow 0 \\ & \text{while } i < n \text{ do} \\ & \text{begin} \\ & i \leftarrow i * 2 \\ & \text{count} \leftarrow \text{count} + 1 \\ & \text{end} \\ & \text{output count} \end{split}$$

0(?)

suppose n = 8

	•		
iteration	i before	i after	count
before loop		1	0
1	1	2	1
2	2	4	2
3	4	8	3

$$\begin{split} \mathbf{i} \leftarrow \mathbf{1} \\ & \text{count} \leftarrow \mathbf{0} \\ & \text{while } \mathbf{i} < \mathbf{n} \text{ do} \\ & \text{begin} \\ & \mathbf{i} \leftarrow \mathbf{i} * \mathbf{2} \\ & \text{count} \leftarrow \text{count} + \mathbf{1} \\ & \text{end} \\ & \text{output count} \end{split}$$

O(?)

logn 2 = n

suppose n = 8

30pp030 11 -	0		
iteration	i before	i after	count
before loop		1	0
1	1	2	1
2	2	4	2
3	4	8	3

suppose n = 32

iteration	i before	i after	count
before loop		1	0
1	1	2	1
2	2	4	2
3	4	8	3
4	8	16	4
5	16	32	5



$$\begin{split} & i \leftarrow 1 \\ & \text{count} \leftarrow 0 \\ & \text{while } i < n \text{ do} \\ & \text{begin} \\ & i \leftarrow i * 2 \\ & \text{count} \leftarrow \text{count} + 1 \\ & \text{end} \\ & \text{output count} \end{split}$$

$$O(?)$$
 $O(\log n)$

suppose n = 8

iteration	i before	i after	count
before loop		1	0
1	1	2	1
2	2	4	2
3	4	8	3

suppose n = 32

iteration	i before	i after	count
before loop		1	0
1	1	2	1
2	2	4	2
3	4	8	3
4	8	16	4
5	16	32	5

COMP108-03-Efficiency-02

Summary: Measuring algorithm efficiency

Next: Data structures - arrays

For note taking