Tutorial 2 solutions

in a floating point system: $x=\frac{1}{2}\left(\frac{dx}{dx}\right)\beta^{E}$

· Consider a floating point system with B=10, p=4, and U=3

The largest number we can represent is 0.999 × 103 (we want to make each digit in the mantissa as large as possible, and make the exponent as large as possible too)

· Therefore, for a general floating point system, the largest number, xx, must satisfy the following:

1) dk=B-1 4k (so each digit in the mantissa is as large as possible

2) sign = +

3) E=V (ic. exponent is as large as possible)

Substituting these properties into our expression for x, we get:

$$X_{k} = \sum_{k=0}^{p-1} \left(\frac{\beta-1}{\beta^{k}} \right) \beta^{V}$$

• Now we can use the following equality: $\sum_{n=0}^{N} \frac{\alpha^{-1}}{\alpha^n} = \alpha - \alpha^{-1}$, where N=p-1 n=k $\alpha=B$

$$:= \times_{*} = (\beta - \beta^{-(p-1)}) \beta^{U}$$

: the largest number
we can represent

is $x_* = (1-\beta^{-p})\beta^{\nu+1}$

FIVE STAR.

2) Recall that an n-bit binary warriable can take on 2n unique values.

Now, if our floating point system can replesent in unique values for the mantissa, and e unique values for the exponent, the total number of unique floating point numbers it can represent is:

M.e (ignoring the sign)

If we use 8 bits to store the exponent, we can have 28 unique values for the exponent. Using 23 bits for the mantissa means it can take on 23 unique values.

.. we can represent 28.223 = 231 unique floating point numbers

If we use q bits for the exponent and 22 for the mantissa, we can represent 29.22=231 unique floating point numbers.

This means, both systems can represent the same number of Floating point numbers. However, the question asks about normalized floats.

Recall that if the exponent is all zeroes, the number is de-normalized. Therefore, we should exclude this exponent from our calculations

: 8-bit exponent & 23-bit mantissa can represent $(2^8-1)(2^3)=2^3-2^3$ unique normalized floats

9-bit exponent & 22-bit mantissa can represent (29-1)(22)=231-222 Normalized floats

this system has more normalized floating point numbers

 $0=5.659\times10^4$ $C=9.337\times10^3$ B=10 P=4 L=-8 U=8 $b=5-629\times10^4$ $d=7.529\times10^{-1}$ (3-) 1. atb = 11.288×10 which gets rounded to (1.129×105) Using A "round to even" 2. A-b= 0.03 ×10 which is represented as 3.000×102) 3. C/d = 1240.138133 - rounds to (1.240 × 10) 4. $b \times d = 42380.741 \longrightarrow (4.238 \times 10^3)$ 5. atc = 57523.7 -> 57520 6+d = 56290.7529 ---> 56290 (n+c) - (b+d) = 57520 - 56290 = 1230 - (1.230 × 10³)Better way of doing culculation: (a-b) + (c-d)a-b => 3.000 x 102 (4ee part 2) btd = 932.9471 - 9 9.329 x102 :. $(a-b) + (c-b) = 1232.9 \rightarrow (1.233 \times 10^{3})$ notice that this is closer to the true value of 1.2324471 x103 than this



Rather than doing x2-y2, do (x+y)(x-y)

Why is this better?

Consider the following case, where x and y are similar in size:

Y=x+2, where EKXx, but there is still enough precision to represent y exactly

ie. the difference between x and y (E) is much smaller than x or y, but the floating point system has just enough precision to represent y exactly (so x+& \pm x)

If we calculate y', we get y2 - (x+8)2 = x2 +2x8 + 62

since Exxx, E2xxx so there is not enough plecision to include it, so y2 gets represented as x2+1x8

Aside: For example, if x=1.0, &= 1.0x10 and p=2. $X+\xi^2=1.01$ which can be regresented exactly $x^2+\xi^2=1.01$ which is represented as 1.0 since $\rho=2$,

Back to our question: Y'is represented as x2+2x8 , so x2-Y2 = -2x8

Now let's do (x+y)(x-y) instead: $x+y = x + (x+\xi) \qquad x-y = x - (x+\xi)$ $= 2x + \xi \qquad = -\xi$

- 82 and 2x2 are much closer in size than & and x2, so there =. (x+y)(x-y)=-2x8-82 W is (hopefully) enough precision to retain the 62 term

2xE = 2.0 x10-1 : 2x2+ 2= 2.1 ×10 which can be represented exactly so the 2 information is not lost

ex. going back to our example here:

82 = 1.0 ×10-2