

PROBLEMAS DE FLUJO EN REDES

Tecnología Digital V: Diseño de Algoritmos
Universidad Torcuato Di Tella

Algoritmo de Ford y Fulkerson



Lester Ford
(1927–2017)



Delbert Fulkerson
(1924–1976)

- El algoritmo de Ford y Fulkerson (1956) obtiene un flujo máximo con complejidad $O(nmU)$, donde $U = \max_{ij \in A} u_{ij}$.

Algoritmo de Ford y Fulkerson

Definir un flujo inicial en N (por ejemplo, $x = 0$)
mientras existe $P :=$ camino de aumento en $R(G, x)$ **hacer**
 para cada arco $ij \in P$ **hacer**
 si $ij \in A$ **entonces**
 $x_{ij} := x_{ij} + \Delta(P)$
 si no ($ji \in A$)
 $x_{ji} := x_{ji} - \Delta(P)$
 fin si
 fin para
fin mientras

Flujo en redes - Camino de aumento

- Dada una red $G = (N, A)$ con función de capacidad u y un flujo factible x , definimos la **red residual** $R(G, x) = (N, A_R)$, donde:
 1. $ij \in A_R$ si $x_{ij} < u_{ij}$,
 2. $ji \in A_R$ si $x_{ij} > 0$.
- Un **camino de aumento** es un camino orientado de s a t en $R(G, x)$.

Flujo en redes - Camino de aumento

- Dado un camino de aumento P , para cada arco $ij \in P$ definimos

$$\Delta(ij) = \begin{cases} u_{ij} - x_{ij} & \text{si } ij \in A \\ x_{ji} & \text{si } ji \in A \end{cases}$$

- Definimos además $\Delta(P) = \min_{ij \in P} \{\Delta(ij)\}$.
- Podemos encontrar un camino de aumento P en la red residual en $O(m)$, y calculamos $\Delta(P)$ en $O(n)$.

Proposición

Sea x un flujo definido sobre una red N con valor F y sea P un camino de aumento en $R(G, x)$. Entonces el flujo \bar{x} , definido por

$$\bar{x}(ij) = \begin{cases} x_{ij} & \text{si } ij \notin P \\ x_{ij} + \Delta(P) & \text{si } ij \in P \\ x_{ij} - \Delta(P) & \text{si } ji \in P \end{cases}$$

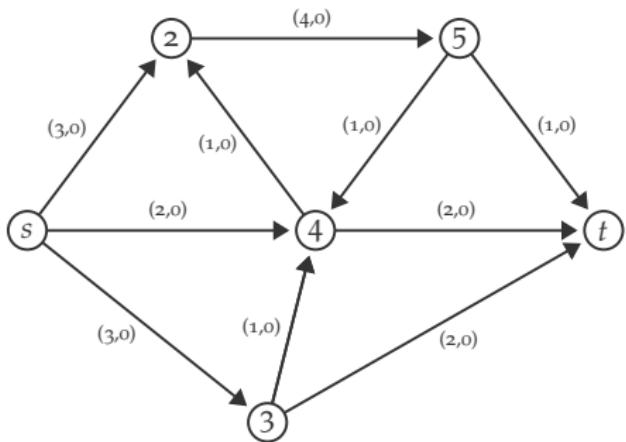
es un flujo factible sobre N con valor $\bar{F} = F + \Delta(P)$.

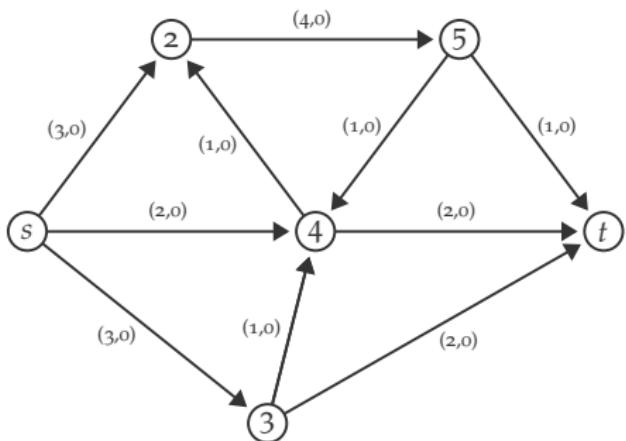
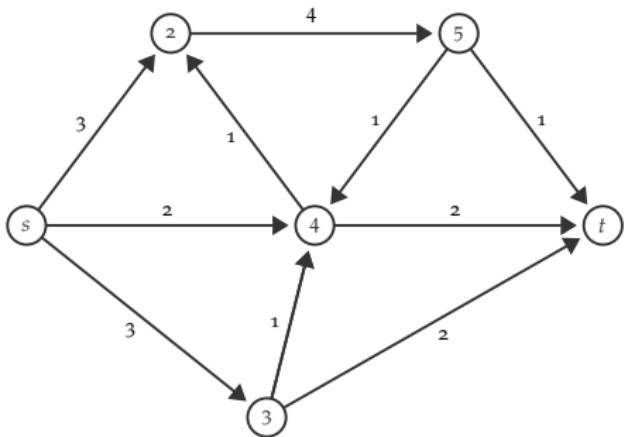
Teorema

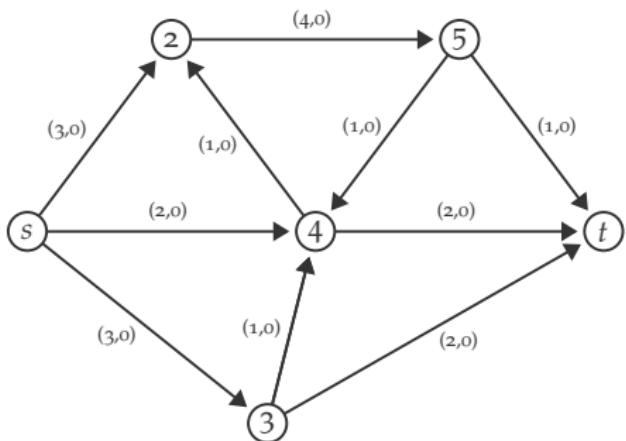
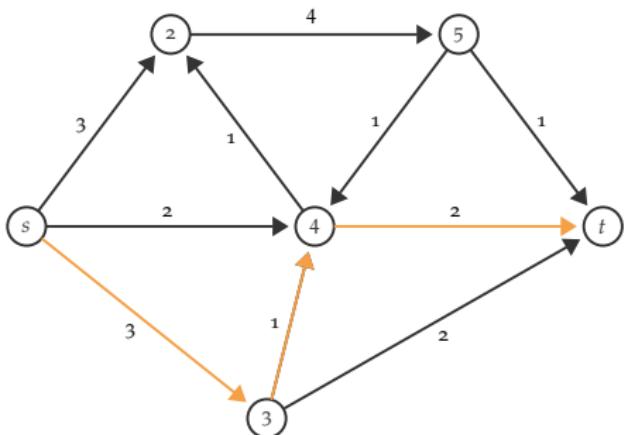
Sea x un flujo definido sobre una red N . Entonces x es un flujo máximo
 \iff no existe camino de aumento en $R(G, x)$.

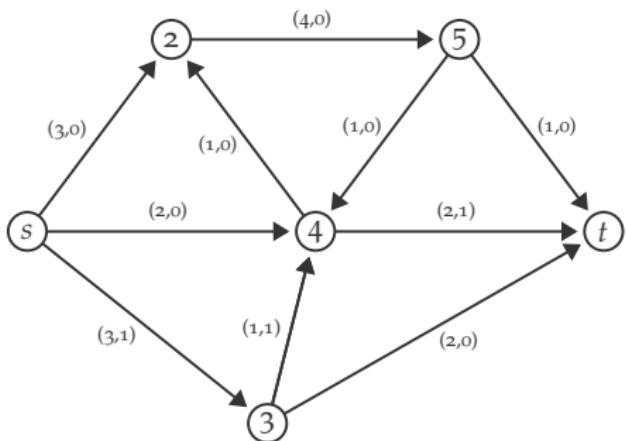
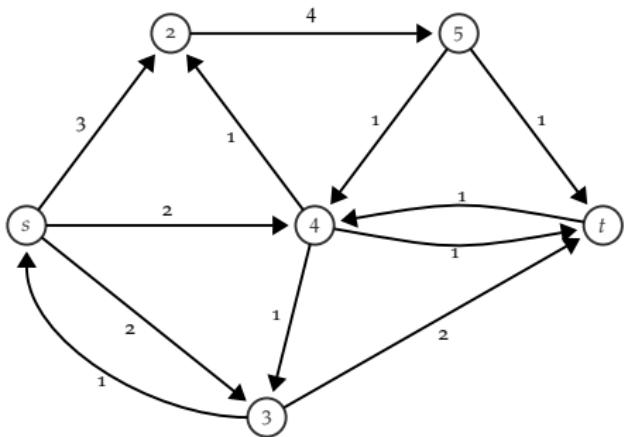
Teorema (max flow-min cut)

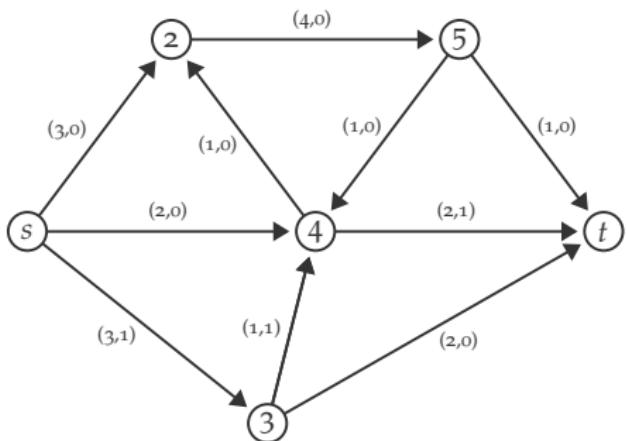
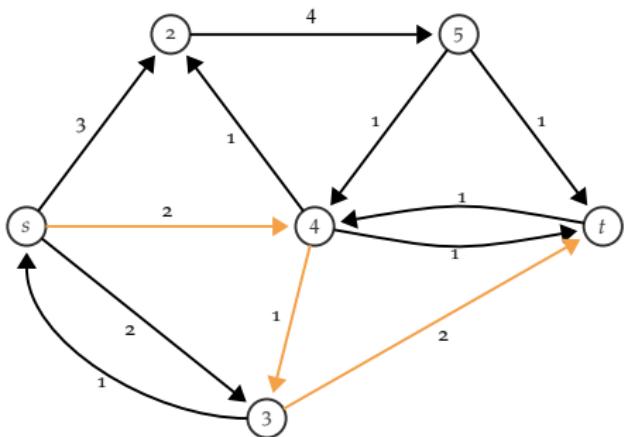
Dada una red N , el valor del **flujo máximo** es igual a la capacidad del **corte mínimo**.

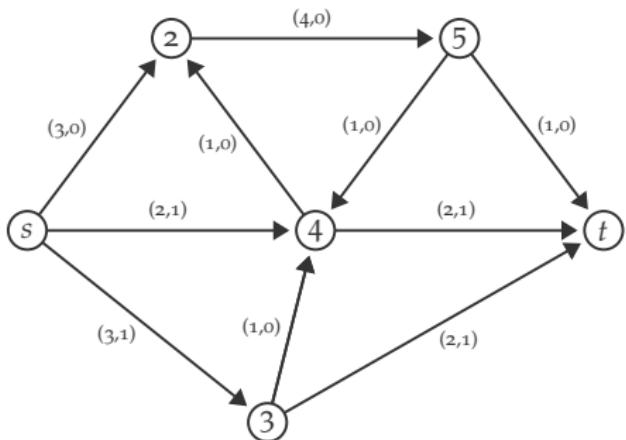
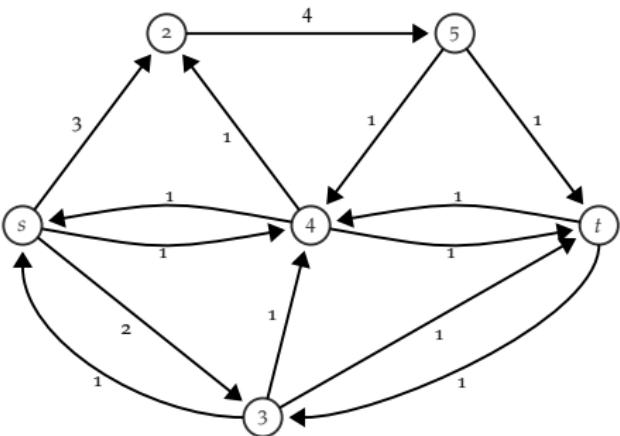


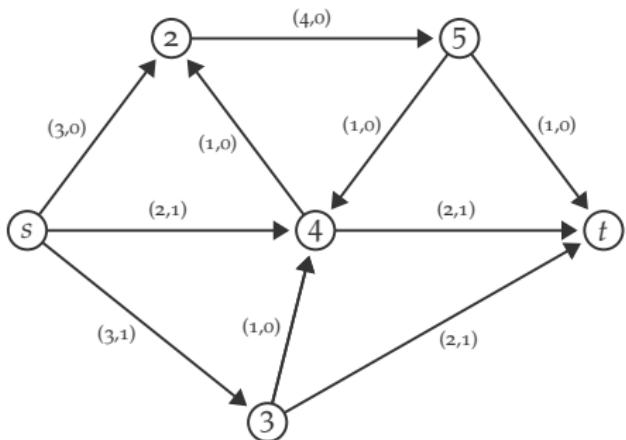
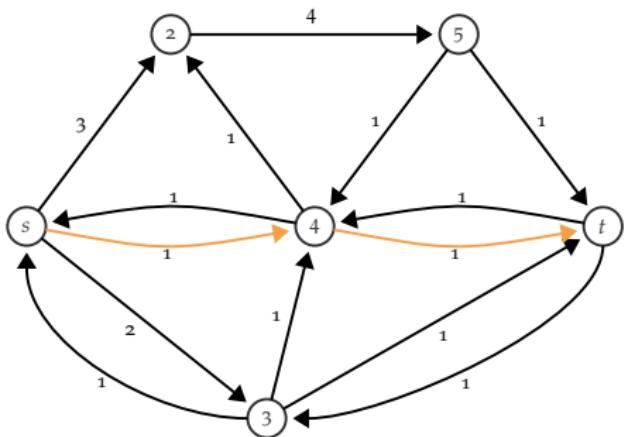

 $R(G, x)$


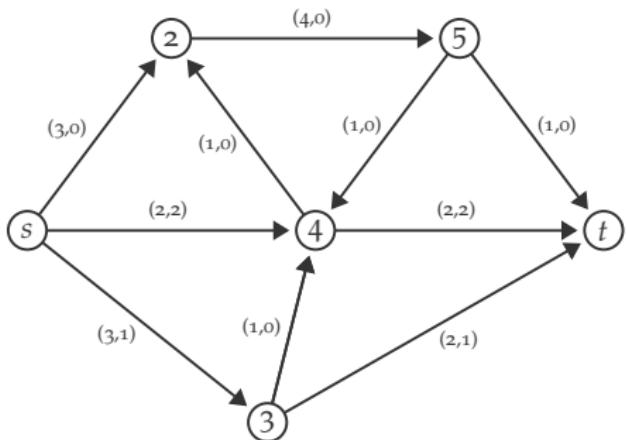
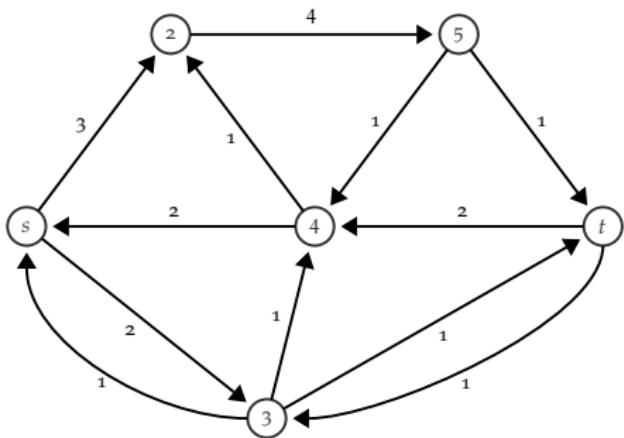

 $R(G, x)$


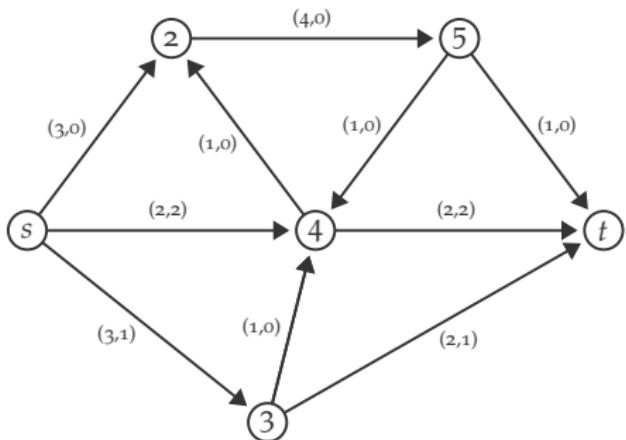
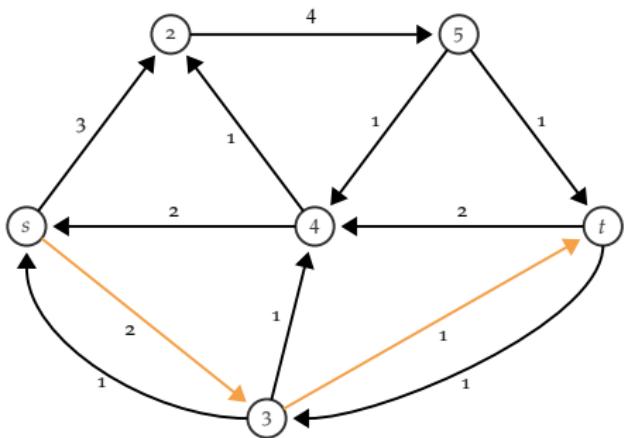

 $R(G, x)$


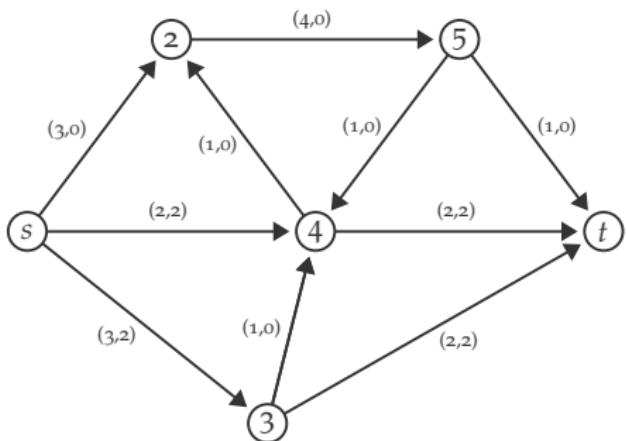
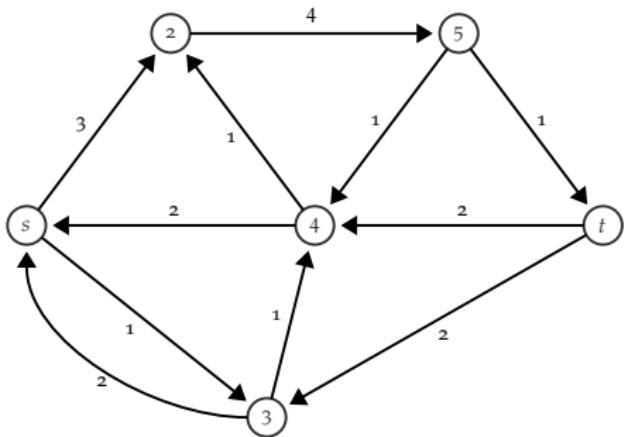

 $R(G, x)$


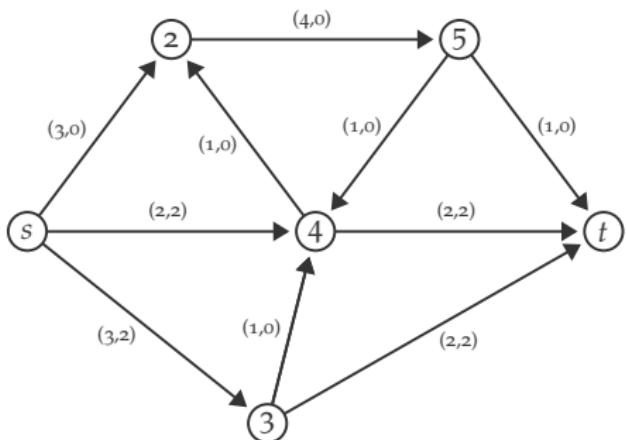
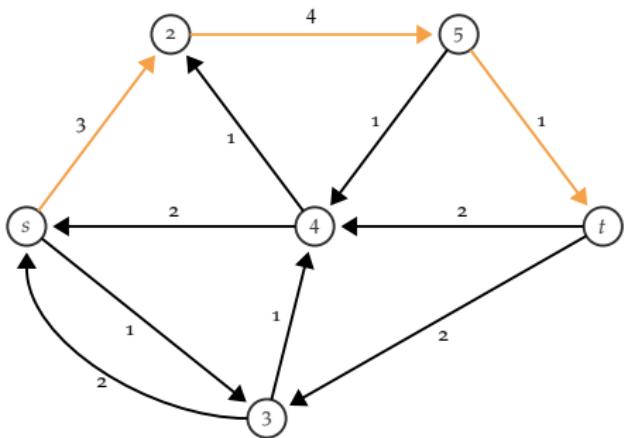

 $R(G, x)$


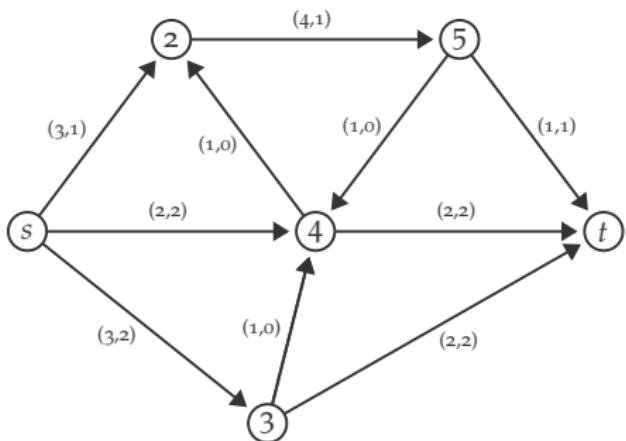
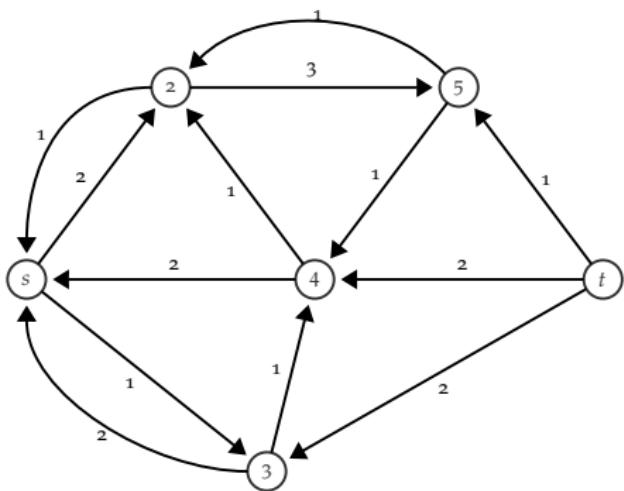

 $R(G, x)$


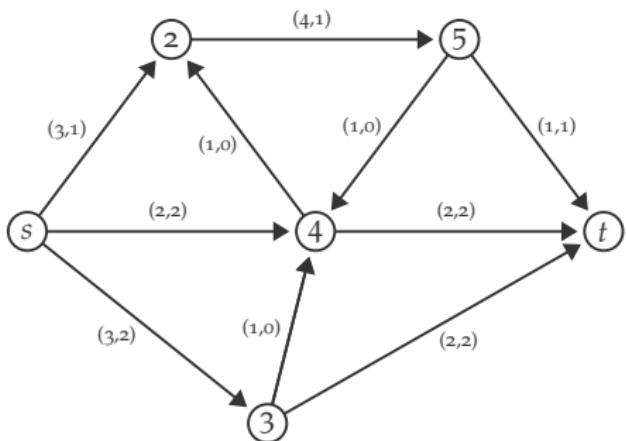

 $R(G, x)$



 $R(G, x)$



 $R(G, x)$



 $R(G, x)$



 $R(G, x)$



 $R(G, x)$
