



Regular Languages and Regular Expressions

Computation and Discrete Structures III

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1. Regular Languages

2. Regular Expressions

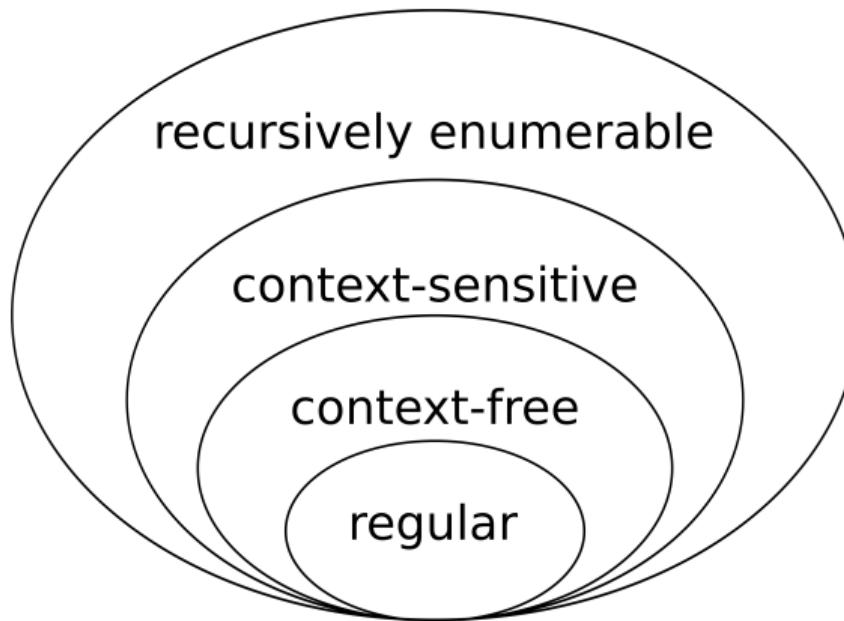
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Languages, Machine Types and Grammars

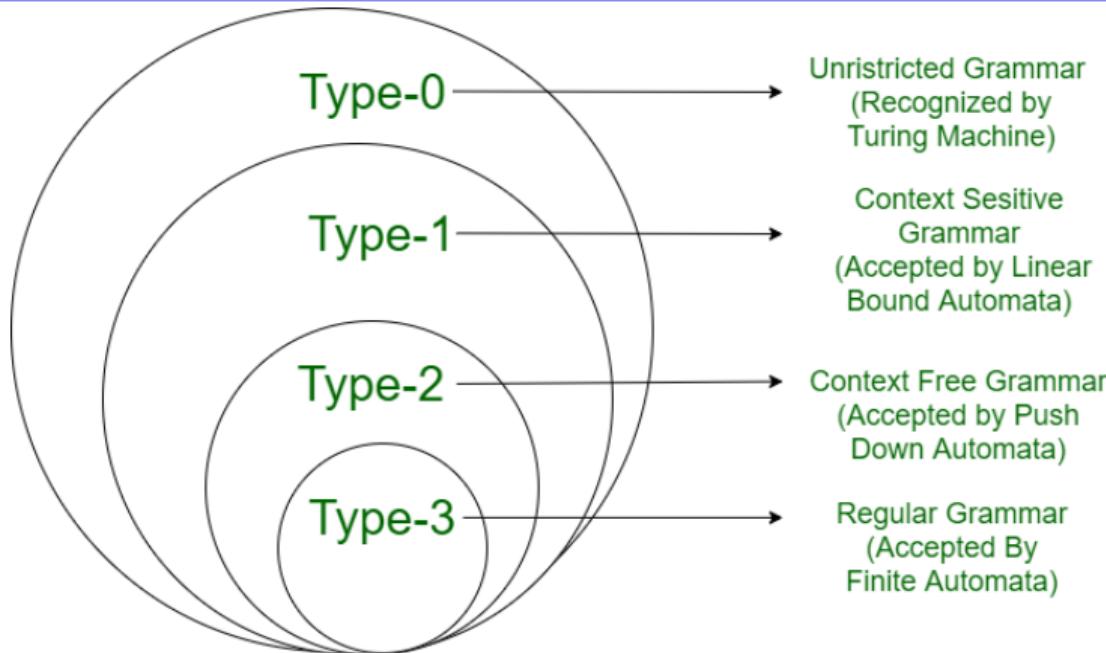
Type	Languages	Machine Type	Grammar Rules
0	Recursively enumerable	Turing machine	Unrestricted
1	Context-sensitive	Linear bounded automaton	$\alpha \rightarrow \beta, \quad \alpha \leq \beta $
2	Context-free	Pushdown automaton	$A \rightarrow \gamma$
3	Regular	Finite automaton	$A \rightarrow aB$ $A \rightarrow a$

Chomsky's Hierarchy



Source: <https://upload.wikimedia.org/wikipedia/commons/9/9a/Chomsky-hierarchy.svg>

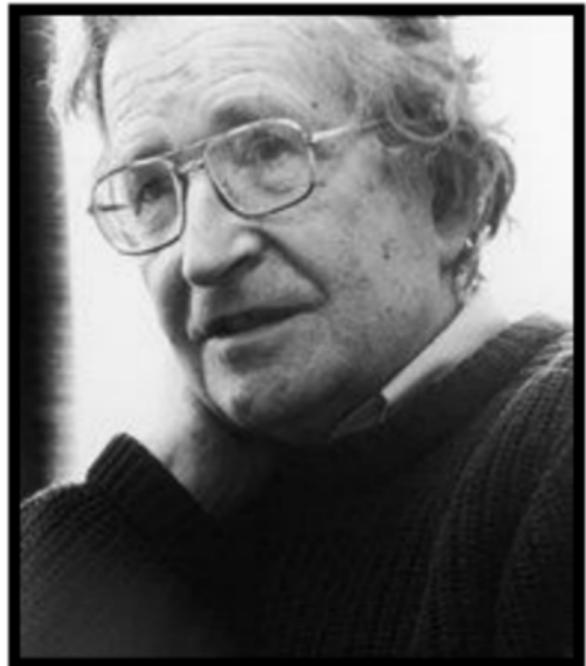
Chomsky's Hierarchy



Source: [https://www.geeksforgeeks.org/theory-of-computation/
chomsky-hierarchy-in-theory-of-computation/](https://www.geeksforgeeks.org/theory-of-computation/chomsky-hierarchy-in-theory-of-computation/)
Regular Languages and Regular Expressions

Noam Chomsky

- Defined context-free grammars
- Creator of the Chomsky hierarchy (1956)
- Defined Chomsky Normal Form (1979)



(1928 -)

Regular languages

Given an alphabet Σ , regular languages over that alphabet are defined recursively as:

- \emptyset is a regular language
- $\{\epsilon\}$ is a regular language
- For every symbol $a \in \Sigma$, $\{a\}$ is a regular language
- If A and B are regular languages, then
 - $A \cup B$, $A \cdot B$ and A^* are regular languages
- No other language is regular

Regular languages

Given $\Sigma = \{a, b\}$, the following statements are correct:

- \emptyset and $\{\epsilon\}$ are regular languages
- $\{a\}$ and $\{b\}$ are regular languages
- $\{a, b\}$ is regular because it is the union of $\{a\}$ and $\{b\}$
- $\{ab\}$ is regular because it is the concatenation of $\{a\}$ and $\{b\}$
- $\{a, ab, b\}$ is regular because it is the union of two regular languages
- $\{a^n \mid n \geq 0\}$ is regular
- $\{a^m b^n \mid m \geq 0 \wedge n \geq 0\}$ is regular
- $\{(ab)^n \mid n \geq 0\}$ is regular

Regular languages

Given $\Sigma = \{a, b, c\}$, indicate if the following languages are regular:

- $\{a\}^*$
- $\{a\}^* \cup \{b\}^*$
- $\{a\}^* \cdot \{b\}^*$
- $\{a, bc\}^*$
- $\{a\} \cdot \{b, c, ab\}$
- $\{a^n b^n \mid n \geq 0\}$
- $\{a^l b^m c^n \mid l \geq 0, m \geq 0, n \geq 0\}$
- $\{a^n b^{2n} \mid n \geq 0\}$

Regular languages

Given $\Sigma = \{a, b, c\}$, indicate if the following languages are regular:

- $\{a\}^*$
- $\{a\}^* \cup \{b\}^*$
- $\{a\}^* \cdot \{b\}^*$
- $\{a, bc\}^*$
- $\{a\} \cdot \{b, c, ab\}$
- $\{a^n b^n \mid n \geq 0\}$, not regular
- $\{a^l b^m c^n \mid l \geq 0, m \geq 0, n \geq 0\}$
- $\{a^n b^{2n} \mid n \geq 0\}$, not regular

$\{a^n b^n \mid n \geq 0\}$ is not regular

- $\{a^n \mid n \geq 0\} = \{\epsilon, a, aa, aaa, \dots\}$
- $\{b^n \mid n \geq 0\} = \{\epsilon, b, bb, bbb, \dots\}$

$aab \in \{\epsilon, a, aa, aaa, \dots\} \cdot \{\epsilon, b, bb, bbb, \dots\}$
but it does not satisfy $a^n b^n$

Regular languages

- Develop the language $L = \{abc, ab, a\}^+$
 $L = \{abc, ab, a, abcabc, abcab, abca, \dots\}$

Regular languages

- Develop the language $L = \{abc, ab, a\}^+$
 $L = \{abc, ab, a, abcabc, abcab, abca, \dots\}$
- Compare it with $\{abc, ab, a\}^*$

Regular languages

- Develop the language $L = \{abc, ab, a\}^+$
 $L = \{abc, ab, a, abcabc, abcab, abca, \dots\}$
- Compare it with $\{abc, ab, a\}^*$
 $\{abc, ab, a\}^* = \{\epsilon, abc, ab, a, abcabc, abcab, abca, \dots\}$
 $\{abc, ab, a\}^+ = \{abc, ab, a, abcabc, abcab, abca, \dots\}$

Regular languages

- Develop the language $L = \{abc, ab, a\}^+$
 $L = \{abc, ab, a, abcabc, abcab, abca, \dots\}$
- Compare it with $\{abc, ab, a\}^*$
 $\{abc, ab, a\}^* = \{\epsilon, abc, ab, a, abcabc, abcab, abca, \dots\}$
 $\{abc, ab, a\}^+ = \{abc, ab, a, abcabc, abcab, abca, \dots\}$
 $\{abc, ab, a\}^+ = \{abc, ab, a\}^* \cdot \{abc, ab, a\}$
- In general, $A^+ = A^* \cdot A$ holds

Regular languages

Indicate if the following languages are regular:

- $\{ab^n a \mid n \geq 0\}$
- $\{a^n b^m c^{n+m} \mid n, m \geq 0\}$
- $\{a^n b^m c^n \mid n, m \geq 0\}$
- $\{(abc)^n \mid n \geq 0\}$
- $\{(abc)^n \cdot (cba)^n \mid n \geq 0\}$
- $\{(ab)^n \cdot (cd)^m \mid m, n \geq 0\}$
- $\{wcw \mid w \in \{a, b\}^*\}$
- $\{w \in \{a, b\}^* \mid |w| = 2k, \text{ for } k \geq 0\}$
- $\{aa, ab, ba, bb\}^*$

Regular languages

Develop each of these regular languages:

- $\{a\}^*$
- $\{a\}^* \cup \{b\}^*$
- $\{a\}^* \cdot \{b\}^*$
- $\{a, bc\}^*$
- $\{abc, ab, a\}^+$
- $\{a\} \cdot \{b, c, ab\}$
- $\{(ab)^i \mid i \geq 0\}$
- $\{a^n b^m \mid n \geq 0, m \geq 0\}$
- $\{a^l b^m c^n \mid l \geq 0, m \geq 0, n \geq 0\}$

Regular languages

Develop each of these regular languages:

- $\{a\}^*$
- $\{a\}^* \cup \{b\}^*$, is $ab \in \{a\}^* \cup \{b\}^*?$
- $\{a\}^* \cdot \{b\}^*$
- $\{a, bc\}^*$
- $\{abc, ab, a\}^+$
- $\{a\} \cdot \{b, c, ab\}$
- $\{(ab)^i \mid i \geq 0\}$
- $\{a^n b^m \mid n \geq 0, m \geq 0\}$
- $\{a^l b^m c^n \mid l \geq 0, m \geq 0, n \geq 0\}$

Regular languages

Develop each of these regular languages:

- $\{a\}^*$
- $\{a\}^* \cup \{b\}^*$
- $\{a\}^* \cdot \{b\}^*$, is $bbb \in \{a\}^* \cdot \{b\}^*?$, is $baa \in \{a\}^* \cdot \{b\}^*?$
- $\{a, bc\}^*$
- $\{abc, ab, a\}^+$
- $\{a\} \cdot \{b, c, ab\}$
- $\{(ab)^i \mid i \geq 0\}$
- $\{a^n b^m \mid n \geq 0, m \geq 0\}$
- $\{a^l b^m c^n \mid l \geq 0, m \geq 0, n \geq 0\}$

Regular languages

Develop each of these regular languages:

- $\{a\}^*$
- $\{a\}^* \cup \{b\}^*$
- $\{a\}^* \cdot \{b\}^*$
- $\{a, bc\}^*$, is $bcbca \in \{a, bc\}^*$? , is $baaa \in \{a, bc\}^*$?
- $\{abc, ab, a\}^+$
- $\{a\} \cdot \{b, c, ab\}$
- $\{(ab)^i \mid i \geq 0\}$
- $\{a^n b^m \mid n \geq 0, m \geq 0\}$
- $\{a^l b^m c^n \mid l \geq 0, m \geq 0, n \geq 0\}$

Regular languages

Develop each of these regular languages:

- $\{a\}^* = \{\epsilon, a, aa, aaa, aaaa, \dots\}$
- $\{a\}^* \cup \{b\}^* = \{\epsilon, a, aa, aaa, \dots\} \cup \{\epsilon, b, bb, bbb, \dots\} = \{\epsilon, a, b, aa, bb, aaa, bbb, \dots\}$
- $\{a\}^* \cdot \{b\}^* = \{\epsilon, a, aa, aaa, \dots, ab, aab, aaab, \dots, b, b, bbb, \dots\}$
- $\{a, bc\}^* = \{\epsilon, a, bc, aa, abc, bca, bcbca, aaa, \dots\}$
- $\{abc, ab, a\}^+ = \{abc, ab, a, abcabc, abcab, abca, \dots\}$
- $\{a\} \cdot \{b, c, ab\} = \{ab, ac, aab\}$
- $\{(ab)^i \mid i \geq 0\} = \{\epsilon, ab, abab, ababab, \dots\}$
- $\{a^n b^m \mid n \geq 0, m \geq 0\} = \{\epsilon, a, b, ab, aab, abb, aaab, \dots\}$
- $\{a^l b^m c^n \mid l \geq 0, m \geq 0, n \geq 0\} = \{\epsilon, a, b, c, ab, bc, abc, aa, aab, aac, \dots\}$

Regular languages

Discuss the membership of the following strings given $L = \{a, bc\}^* \cup \{ad, d\}^*$

- Is $bcabc \in L$?
- Is $aabcad \in L$?
- Is $adbc \in L$?
- Is $adad \in L$?
- Is $adddd \in L$?

Regular languages

Discuss the membership of the following strings given $L = \{a, bc\}^* \cup \{ad, d\}^*$

- Is $bcabc \in L$? Yes
- Is $aabcad \in L$? No
- Is $adbc \in L$? No
- Is $adad \in L$? Yes
- Is $addd \in L$? Yes

Interpret regular language

Interpret the type of words that belong to the following language

$$L = \{a\}^* \cup \{b\}^*$$

Interpret regular language

Interpret the type of words that belong to the following language

$$L = \{a\}^* \cup \{b\}^*$$

Interpretation

Strings that have only a's or only b's. These symbols do not appear mixed

Interpret regular language

Interpret the type of words that belong to the following language

$$L = \{a\}^* \cdot \{b\}^*$$

Interpret regular language

Interpret the type of words that belong to the following language

$$L = \{a\}^* \cdot \{b\}^*$$

Interpretation

Strings that have zero or more a's followed by zero or more b's

Regular languages

Given a regular language in natural language, define language structure

Given $\Sigma = \{a, b\}$ define language A of all words that have exactly one a

Given a regular language in natural language, define language structure

Given $\Sigma = \{a, b\}$ define language A of all words that have exactly one a

Language structure

$$A = \{b\}^* \cdot \{a\} \cdot \{b\}^*$$

Regular languages

Given a regular language in natural language, define language structure

Given $\Sigma = \{a, b\}$ define language B of all words that begin with b

Given a regular language in natural language, define language structure

Given $\Sigma = \{a, b\}$ define language B of all words that begin with b

Language structure

$$B = \{b\} \cdot (\{a\} \cup \{b\})^*$$

Regular languages

Given a regular language in natural language, define language structure

Given $\Sigma = \{a, b\}$ define language C of all words that contain the string ba

Given a regular language in natural language, define language structure

Given $\Sigma = \{a, b\}$ define language C of all words that contain the string ba

Language structure

$$C = (\{a\} \cup \{b\})^* \cdot \{ba\} \cdot (\{a\} \cup \{b\})^*$$

1. Regular Languages

2. Regular Expressions

Regular expression

A regular expression is a simplified way of representing a regular language

Regular Language	Regular Expression
$\{ab\}$	ab
$\{a\}^*$	a^*
$\{a\}^+$	a^+
$\{a\} \cup \{b\}$	$a \cup b$

Regular expressions

Some regular expressions:

- b^*
- $b(a \cup b)^*$
- $(a \cup b)^*ba(a \cup b)^*$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that begin with b and end with a

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that begin with b and end with a

Regular expression

$$b(a \cup b)^*a$$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have exactly two a's

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have exactly two a's

Regular expression

$$b^*ab^*ab^*$$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have an even number of a's

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have an even number of a's

Regular expression

$$(b^*ab^*ab^*)^*$$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have even length

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have even length

Regular expression

$$(aa \cup ab \cup ba \cup bb)^*$$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have odd length

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have odd length

Regular expression

$$(aa \cup ab \cup ba \cup bb)^*(a \cup b)$$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have at least one b

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ that have at least one b

Regular expression

$$(a \cup b)^* b (a \cup b)^*$$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ where the second-to-last symbol is an a

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ where the second-to-last symbol is an a

Regular expression

$$(a \cup b)^* a (a \cup b)$$

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ where the third-to-last symbol is an a

Natural language description of a regular expression

Indicate the **regular expression** that denotes the language of all words over $\Sigma = \{a, b\}$ where the third-to-last symbol is an a

Regular expression

$$(a \cup b)^* a (a \cup b) (a \cup b)$$

Exercises for the reader

Indicate the **regular expression** for each of the following languages:

- All strings over $\{a, b\}$ that begin with aa and end with bb
- All strings over $\{a, b\}$ that have exactly three b 's
- All strings over $\{a, b\}$ that begin and end with different symbols

Equivalent regular expressions

1. $r \cup s = s \cup r$
2. $r \cup \emptyset = r = \emptyset \cup r$
3. $r \cup r = r$
4. $(r \cup s) \cup t = r \cup (s \cup t)$
5. $r \cdot \epsilon = \epsilon \cdot r = r$
6. $r \cdot \emptyset = \emptyset \cdot r = \emptyset$
7. $(rs)t = r(st)$
8. $r(s \cup t) = rs \cup rt$
9. $r^* = (r^*)^* = r^*r^* = (\epsilon \cup r)^*$
10. $(r \cup s)^* = (r^* \cup s^*)^* = (r^*s^*)^*$

Equivalent regular expressions

11. $r(sr)^* = (rs)^*r$
12. $(r^*s)^* = \epsilon \cup (r \cup s)^*s$
13. $(rs^*)^* = \epsilon \cup r(r \cup s)^*$
14. $s(r \cup \epsilon)^*(r \cup \epsilon) \cup s = sr^*$
15. $rr^* = r^*r$
16. $(r \cup \epsilon)^* = r^*$

References



- Kozen, D. C. (2007)
Automata and computability
Springer Science & Business Media. Lectures 3–4.