

RtToy

Here we present two variants of Toy for experimenting with mixing call-time choice and run-time choice parameter passing mechanisms. One of them starts with the regular call-time choice framework of Toy and provides a primitive to express run-time choice [1]. In the other the default call-time choice mechanism of Toy has been substituted to provide run-time choice by default and a primitive `where` is used to define local bindings to express sharing of values and therefore call-time choice [2].

Error: Macro TOC(None) failed

```
'NoneType' object has no attribute 'endswith'
```

Call-time choice based variant

The main ideas behind this prototype can be found in [1].

Download

A first prototype can be found here

Basic usage

To use the prototype:

0. Unzip the file `RtToy.zip`

This will create a directory *rttoy* with many files and some folders.

1. Start a Sicstus Prolog session (3.12 or earlier).
2. Change the Sicstus working directory to the created *rttoy*. This can be done by invoking:

```
?- use_module(library(system)),working_directory(Old,<write here the path to 'rttoy'>).
```

For instance:

```
?- use_module(library(system)),working_directory(Old,'f:/systems/rttoy').
```

3. Invoke

```
?- compile(rttoy).
```

This starts a Toy session.

4. You can move to the folder with examples by executing the Toy command

```
Toy> /cd(bancopruebas)
```

5. You can compile the program `ej-runtime.toy` by executing the Toy command

```
/run(ej-runtime)
```

6. You can solve goals. In particular, to evaluate one expression *e*, you execute:

```
Toy> e==X
```

and this will return the value(s) of e as binding(s) for X

Example: Consider the program *ej-runtime.toy* included with the prototype

```
/******
Some examples to test RTToy, an extension of Toy to support combination
of call-time choice and run-time choice
******/

/****** First, some general functions *****/

infixr 5 //
X // Y = X
X // Y = Y

infixr 6 ++
[] ++ Ys = Ys
[X|Xs] ++ Ys = [X|Xs++Ys]

reverse [] = []
reverse [X|Xs] = reverse Xs ++ [X]

take N [] = []
take N [X|Xs] = if N<=0 then [] else [X|take (N-1) Xs]

/****** Example 1: string generation by grammar rules *****/

letter = "a" // "b" // "c"

star X = [] // X ++ star X

word = star (rt letter) % or equivalently, rt (star letter)

palindrome = palaux word

palaux X = X++([]//letter)++reverse X

/****** Example 2: variations upon coin *****/

coin = 0
coin = 1

double X = X+X

f X = g X coin
g X Y = (X,X,Y,Y)

a = double (rt coin) % or equivalently, rt (double coin)
b = double coin
c = f (rt coin) % or equivalently, rt (f coin)
a' = rt (double coin)

/****** Example 3: numbers and repetitions *****/

repeat X = [X|repeat X]

numberCt N = take N (repeat (0//1//2))
numberRt N = take N (repeat (rt (0//1//2)))
```

Some goal examples with this program follow:

```

Toy> a == L
  { L -> 0 }
  Elapsed time: 0 ms.
sol.1, more solutions (y/n/d/a) [y]?
  { L -> 1 }
  Elapsed time: 0 ms.
sol.2, more solutions (y/n/d/a) [y]?
  { L -> 1 }
  Elapsed time: 0 ms.
sol.3, more solutions (y/n/d/a) [y]?
  { L -> 2 }
  Elapsed time: 0 ms.
sol.4, more solutions (y/n/d/a) [y]?
  no
  Elapsed time: 0 ms.

Toy> c == X
  { X -> (0, 0, 0, 0) }
sol.1, more solutions (y/n/d/a) [y]?
  { X -> (0, 0, 1, 1) }
sol.2, more solutions (y/n/d/a) [y]?
  { X -> (0, 1, 0, 0) }
sol.3, more solutions (y/n/d/a) [y]?
  { X -> (0, 1, 1, 1) }
sol.4, more solutions (y/n/d/a) [y]?
  { X -> (1, 0, 0, 0) }
sol.5, more solutions (y/n/d/a) [y]?
  { X -> (1, 0, 1, 1) }
sol.6, more solutions (y/n/d/a) [y]?
  { X -> (1, 1, 0, 0) }
sol.7, more solutions (y/n/d/a) [y]?
  { X -> (1, 1, 1, 1) }

```

7. The use of *rt* annotations is also allowed in goals. For instance:

```

Toy> rt (double coin) == L
  { L -> 0 }
sol.1, more solutions (y/n/d/a) [y]?
  { L -> 1 }
sol.2, more solutions (y/n/d/a) [y]?
  { L -> 1 }
sol.3, more solutions (y/n/d/a) [y]?
  { L -> 2 }
sol.4, more solutions (y/n/d/a) [y]?
  no

```

Run-time choice based variant

The main ideas behind this prototype can be found in [2].

- Excepting the use of sharing bindings (see below) the programs behave like term rewriting systems. In other words, the default behaviour is run-time choice.
- The express that the evaluation of a certain expression is subject to sharing explicit bindings are used, written using the `where` construction. The reason to use `where` instead of `let`, as it is done in [2] has been that it makes the implementation easier, as the `where` construction was already present in the syntax of Toy.

Download

A first prototype can be found here

Basic usage

1. Open a Sicstus Prolog session (the Sicstus version must be previous to 4)

2. Execute

```
?- use_module(library(system)),working_directory(Old,<rttoy directory>).
```

For example

```
use_module(library(system)),working_directory(Old,'f:/0.trabajo/0.juanrhortala/svn/wlpe2008-1/rttoyprolog')
```

3. Execute

```
?- compile(rttoy).
```

4. Now, inside the Toy session, execute

```
/prolog(assert(option(rt_all)))
```

As a result now we have a run-time choice behaviour by default.

5. Call-time choice behaviour can be achieved by using sharing bindings by means of the where syntax. For example:

```
double X = Y+Y where Y=X.
```

In [2] can be found a precise formal specification of the bindings which must be introduced in a program in order to get call-time choice behaviour for the whole program.

6. To find some example: inside the Toy session, move to:

```
/cd(bancopruebas)  
/run(ej-runtime-pepm)
```

Other resources

References

1. ^ F.J. López-Fraguas, J. Rodríguez-Hortalá and J. Sánchez-Hernández , A Lightweight Combination of Semantics for Non-deterministic Functions, In Proc. WLPE'08, 2008, <http://gpd.sip.ucm.es/fraguas/papers/WLPE08.pdf>
2. ^ F.J. López-Fraguas, J. Rodríguez-Hortalá and J. Sánchez-Hernández , A Flexible Framework for Programming with Non-deterministic Functions, In Proc. PEPM'09, ACM Press, 2009, <http://gpd.sip.ucm.es/fraguas/papers/PEPM09.pdf>

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