## Practice Exam W24 - Results



## Attempt 1 of Unlimited

Written Apr 19, 2024 10:14 PM - Apr 19, 2024 10:14 PM

Attempt Score 0 %

Overall Grade (Highest Attempt) 0 %

Question 1 0 / 1 point

The execution mode of overlapping: (select all that apply)

- x can occur on a multiprocessor system
  - ✓ can occur on a single-processor system
- x is an example of concurrent processing
  - ✓ is the same as interleaving

Question 2 0 / 1 point

When processes use and update shared data without reference to other processes, but know other processes may access the same data, and cooperate to ensure the data is managed properly, they exhibit this relationship:

- competition
- cooperation by sharing
  - cooperation by communication
  - cooperation by contention

Question 3 0 / 1 point

Consider function P1 of Peterson's algorithm for mutual exclusion, shown below. What should replace XX?

```
void P1() {
    while (true) {
        flag [1] = true;
        turn = 0;
        while (XX) /* do nothing */;
        /* critical section */
        flag [1] = false;
        /* remainder */
}
```

```
flag [0] && turn == 0

flag [0] && turn == 1

flag [1] && turn == 0
```

flag [1] && turn == 1

Question 4 0 / 1 point

Operations wait and signal may be performed on a semaphore. The wait operation:

	always increments the semaphore's value
Þ	always decrements the semaphore's value
	always blocks the process

only returns when a corresponding signal is executed

Question 5 0 / 1 point

Consider the following solution to the bounded-buffer Producer/Consumer problem. What is the purpose of semaphore e?

```
semaphore s=1, n=0, e=sizeofbuffer
void main() {
   parbegin (producer, consumer);
}
```

```
void producer() {
  while (true) {
    produce();
    semWait(e);
    semWait(s);
    append();
    semSignal(s);
    semSignal(n);
}
```

```
void consumer() {
    while (true) {
        semWait(n);
        semWait(s);
        take();
        semSignal(s);
        semSignal(e);
        consume();
}
```

to ensure the producer can't add data to a full buffer
 to ensure the consumer can't remove data from an empty buffer
 to ensure mutual exclusion on the buffer
 to ensure no race conditions on semaphore s

Question 6 0 / 1 point

Which are true about traditional Hoare monitors? (select all that apply)

- multiple processes at a time may be executing in the monitor
- it must have only global variables; no local variables are allowed
- $\Rightarrow$   $\times$   $\bigcirc$  a process enters the monitor by invoking one of its procedures
- $\Rightarrow$  x synchronization is provided by the use of condition variables

Question 7 0 / 1 point

When processes use a mailbox to send/receive messages:

- they use direct addressing
- they use indirect addressing
  - they perform a rendezvous
  - the mailbox must be owned by the OS

Question 8 0 / 1 point

Which is true of the following solution to the Readers/Writers problem. (select all that apply)

```
int readcount;
semaphore x=1, wsem=1;

void main () {
  readcount=0;
  parbegin (reader, writer);
}
```

```
void writer {
  while (true) {
    semWait (wsem);
    WRITEUNIT();
    semSignal (wsem);
} }
```

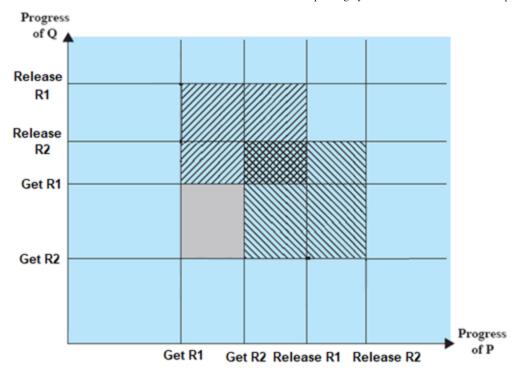
```
void reader {
    while (true) {
        semWait (x);
        readcount++;
        if (readcount==1) semWait (wsem);
        semSignal(x);
        READUNIT();
        semWait(x);
        readcount--;
        if (readcount==0) semSignal (wsem);
        semSignal (x);
}
```

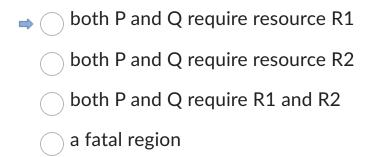
- \Rightarrow 💢 🦳 readers have priority
  - ✓ writers have priority
  - multiple writers may perform WRITEUNIT at the same time
- multiple readers may perform READUNIT at the same time

Question 9 0 / 1 point

Consider the following Joint Progress Diagram for processes P and Q below. What does the area with upward slanted lines indicate?







Question 10 0 / 1 point

There are 3 conditions that are necessary, but not sufficient, for a deadlock to exist. Actual deadlock comprises these 3 conditions along with a circular wait. These 3 conditions are mutual exclusion, hold-and-wait, and:

pre-emption

no pre-emption

resource allocation

a linear ordering of resources

Question 11 0 / 1 point

A system contains 15 units of resource R. Maximum Claims and Current Allocations for processes P1-P5 are given below (e.g., P1 currently holds 2 units of R, although its maximum claim is 3 units). Suppose P1 and P2 execute concurrently, while P5, P4 and P3 request initiation, in that order. Using the Process Initiation Denial approach, the OS will allow which processes to begin execution?

	P1	P2	P3	P4	P5	P1	P2
	3	2	5	5	5	2	1
•	Maximum Claim				Cur	rent	
					Alloc	ation	

- none may begin
- only P5
- only P5 and P4
  - P5, P4, and P3

Question 12 0 / 1 point

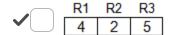
The Claim and Allocation matrices below describe the state of a system consisting of 4 processes and 3 resources. The state would be **safe** for which of these Resource vectors? (select all that apply)

	R1	R2	R3	
P1	2	2	3	
P2	1	1	3	
P3	1	1	1	
P4	1	2	5	
Claim matrix C				

	R1	R2	R3		
P1	1	0	1		
P2	1	1	1		
P3	1	1	1		
P4	0	0	2		
	A 11 41 4 1 A				

Allocation matrix A

	R1	R2	R3	
<b>~</b>	3	2	5	



Question 13 0 / 1 point

A major advantage of Deadlock Avoidance algorithms are that they:

- are less restrictive than deadlock prevention
  - although they do rollbacks, there are less rollbacks than for deadlock detection
  - although they do pre-emptions, there are less pre-emptions than for deadlock detection
  - allow deadlock, but recover from it

Question 14 0 / 1 point

Which are true about Linux unnamed pipes? (select all that apply)

- - may be deleted using rm from a bash shell
- may be created from a bash shell
  - OS provides no mutual exclusion for them

Question 15 0 / 1 point

Which true about Linux kernel spinlocks? (select all that apply)

- typically used when wait time for acquiring a lock is expected to be very long
- 🛶 \chi 🦳 uses busy-waiting
- $\Rightarrow$  x most common technique for protecting critical sections in Linux

one spinlock may be acquired by multiple processes simultaneously
Question 16 0 / 1 point
What is true of the memory management scheme of Dynamic Partitioning? (select all that apply)
⇒ 🗶 🔃 a process is allocated exactly as much memory as it requires
it has internal fragmentation
⇒ 🗙 🦳 it has external fragmentation
small jobs will not utilize partition space efficiently
Question 17 0 / 1 point
Consider simple paging, with no virtual memory, as outlined in Chapter 7 of our text. Assume 16-bit addresses, and page size of $4K = 4096 = 2^{12}$ bytes. A process can consist of a maximum of X pages of $4K$ bytes. What is X?
<u>4</u>
<u> </u>
<u> </u>
→ 16
Question 18 0 / 1 point
Which are true when virtual memory is used? (select all that apply)
⇒ 🗙 🔲 a process can be larger than all of main memory
program-generated addresses are translated automatically to corresponding machine addresses
the size of virtual storage is limited by the number of frames of main memory
the size of virtual storage is proportional to the size of disk

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Question 19	0 / 1 poin

For virtual memory to be practical and effective, 2 ingredients are needed: (1) the OS must include software for managing the movement of pages/segments between secondary memory and main memory, and (2) hardware support for

The page placement policy: (select all that apply)	
Question 20	0 / 1 point
semaphores	
fixed partitioning with variable size partitions	
dynamic partitioning	
paging and/or segmentation	

The page placement policy: (select all that apply)

- x determines where a process is to reside in real memory
  - determines where a process is to reside in secondary memory
  - ✓ is relevant in pure paging systems
- x is relevant for NUMA systems

Question 21 0 / 1 point

Suppose a process does the following sequence of page references and is allocated a fixed 3 frames. How many page faults occur using the LRU Policy after the frame allocation is initially filled? 7,8,7,6,9,8,7

- $\bigcirc$  0
- $\bigcirc$  1
- $\bigcirc$  2
- **⇒** 3

Question 22 0 / 1 point

With this type of allocation policy, whenever a page fault occurs in the execution of a process, one of the pages of that process must be replaced by the needed page.

⇒ fixed allo	cation	
variable a	allocation	
resident s	set allocation	
odynamic a	allocation	
Question 23		0 / 1 point
A cleaning polic	cy: (select all that apply)	
<b>~</b> ~	cerned with when a modified page dary memory	e should be written out to
⇒ 🗶 📄 is the o	opposite of a fetch policy	
	cerned with where on secondary n I be written	nemory a modified page
is conc placed	cerned with where in main memor	y a modified page should be
Question 24		0 / 1 point
Lab06 used POS apply)	SIX semaphores. What was true a	bout them? (select all that
⇒    ✓ named	l semaphores persist, even when r	no processes are using them
✓ an unn of a prediction of a prediction of a prediction.	named semaphore cannot be share ocess	ed among (pthread) threads
the init sem_o	tial value of a named semaphore o pen	can be set by function
	equire library <semaphore.h></semaphore.h>	

Question 25 0 / 1 point

Which are true about a System V message queue, as used in kirk.c & spock.c in LabO8? (select all that apply)

<b>⇒</b> ×	the first member of the message structure must have type	long (	int)
<b>⇒</b> ×	msgsnd blocks when the message queue is full		

✓ its permissions must always be exactly 600

it is automatically deleted by Linux when no process is using it

Done