Twitter thread by @judijasa

Title: The logarithmic chain complex

Slide Nr: 001

The conventional sum, +, and multiplication,  $\times$ , are abelian products with distributivity. For example, if,  $\alpha, \beta, \gamma \in (\mathbb{R}, +, \times)$ , then,

$$\alpha \times (\beta + \gamma) = \alpha \times \beta + \alpha \times \gamma. \tag{1}$$

Formally, these two products form a *ring*. As aspiring goldsmiths we shall try to make chains out of these rings. Let,  $\overset{0}{\times}$ , and,  $\overset{1}{\times}$ , denote two products that behave like the conventional sum and multiplication, respectively. The set of products<sup>1</sup>,

$$\alpha \stackrel{n+1}{\times} \beta \doteq \log^{-1} \left( \log(\alpha) \stackrel{n}{\times} \log(\beta) \right),$$
 (2)

for,  $n \in \mathbb{Z}$ , forms a sequence of abelian groups and homomorphisms. Note that,  $\overset{n+1}{\times}$ , is distributive in,  $\overset{n}{\times}$ . Proof by induction: Assume,

$$\alpha \stackrel{n}{\times} (\beta \stackrel{n-1}{\times} \gamma) = (\alpha \stackrel{n}{\times} \beta) \stackrel{n-1}{\times} (\alpha \stackrel{n}{\times} \gamma), \tag{3}$$

then,

$$\alpha \stackrel{n+1}{\times} (\beta \stackrel{n}{\times} \gamma) = \log^{-1} \{ \log(\alpha) \stackrel{n}{\times} \log(\beta \stackrel{n}{\times} \gamma) \}$$

$$= \log^{-1} \{ \log(\alpha) \stackrel{n}{\times} [\log(\beta) \stackrel{n-1}{\times} \log(\gamma)] \}$$

$$\stackrel{(3)}{=} \log^{-1} \{ [\log(\alpha) \stackrel{n}{\times} \log(\beta)] \stackrel{n-1}{\times} [\log(\alpha) \stackrel{n}{\times} \log(\gamma)] \}$$

$$= \log^{-1} \{ \log(\alpha \stackrel{n+1}{\times} \beta) \stackrel{n}{\times} \log(\alpha \stackrel{n+1}{\times} \gamma) \}$$

$$= (\alpha \stackrel{n+1}{\times} \beta) \stackrel{n}{\times} (\alpha \stackrel{n+1}{\times} \gamma). \tag{4}$$

This covers the proof for,  $n \in \mathbb{Z}^+$ . Left to the reader is the proof for,  $n \in \mathbb{Z}^-$ .

In future posts, I will share more trivia about this object and discuss possible applications.

<sup>&</sup>lt;sup>1</sup>Details such as the base of the log function and its algebraic role are discussed later on.

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Slide Nr: 002

In the latest slide we defined a sequence of abelian products, based on the log function. A complete definition requires a generalisation of the log function to include negative reals as input. Before doing this, let us appreciate the behaviour of these products in,  $\mathbb{R}^+$ . The definition in Eq. (2) makes it clear these products are commutative. As observed in Coya's law, this is not always obvious, e.g., the equivalent definition of,  $\stackrel{2}{\times}$ :

$$\alpha \stackrel{2}{\times} \beta = \alpha^{\log \beta}. \tag{5}$$

The product,  $\overset{2}{\times}$ , is interesting, together with,  $\overset{-1}{\times}$ , are the closest relatives to the conventional products,  $\overset{0}{\times}$ , and,  $\overset{1}{\times}$ , for which we hold more intuition. The product,  $\overset{2}{\times}$ , is distributive in the conventional multiplication ( $\overset{1}{\times}$ ). This makes,  $\overset{2}{\times}$ , the abelian version of the power product, which is also distributive in the multiplication, i.e.,  $(\alpha \times \beta)^{\gamma} = \alpha^{\gamma} \times \beta^{\gamma}$ .

It is convenient to move into a more abstract characterisation of the sequence of products. Some definitions and conventions:

$$\overset{n}{\alpha} \overset{n}{\times} \overset{n}{\beta} = \overset{n}{\beta} \overset{n}{\times} \overset{n}{\alpha}, \tag{6a}$$

$$\stackrel{n}{\alpha} \stackrel{n}{\times} \stackrel{n}{I_0} = \stackrel{n}{I_0} \stackrel{n}{\times} \stackrel{n}{\alpha} = \stackrel{n}{\alpha}, \tag{6b}$$

$$\stackrel{\scriptscriptstyle n}{\alpha} \stackrel{\scriptscriptstyle n}{\times} \stackrel{\scriptscriptstyle n}{\rm I}_1 = \stackrel{\scriptscriptstyle n}{\rm I}_1 \stackrel{\scriptscriptstyle n}{\times} \stackrel{\scriptscriptstyle n}{\alpha} = \stackrel{\scriptscriptstyle n}{\rm I}_1. \tag{6c}$$

$$\left(\tilde{\mathcal{T}}\left(\tilde{\mathcal{T}}^{n}\alpha\right)\right) = \tilde{\alpha} \tag{7a}$$

$$\stackrel{n}{\alpha} \stackrel{n}{\times} (\stackrel{n}{\mathcal{T}} \stackrel{n}{\alpha}) = (\stackrel{n}{\mathcal{T}} \stackrel{n}{\alpha}) \stackrel{n}{\times} \stackrel{n}{\alpha} = \stackrel{n}{\mathrm{I}_0}, \tag{7b}$$

$$\mathcal{J}^{n} I_{0}^{n} = I_{0}^{n} \tag{7c}$$

$$\stackrel{n}{\alpha} \stackrel{n}{\times} (\stackrel{n}{\mathcal{T}} \stackrel{n}{I_1}) = (\stackrel{n}{\mathcal{T}} \stackrel{n}{I_1}) \stackrel{n}{\times} \stackrel{n}{\alpha} = \stackrel{n}{\mathcal{T}} \stackrel{n}{I_1}. \tag{7d}$$

Abusing of notation on behalf of intuition we keep,  $\log(.)$ , to denote the generator of homomorphisms. We use,  $\log^{+1}(.) = \log(.)$ , and,  $\log^{-1}(\log(\tilde{\alpha})) = \log(\log^{-1}(\tilde{\alpha})) = \tilde{\alpha}$ . The generator of homomorphisms connects nearby groups along the sequence, i.e.,

$$\stackrel{n-1}{\alpha} = \log(\stackrel{n}{\alpha}) \in \stackrel{n}{\mathcal{T}} \stackrel{n}{\mathrm{I}_1} \tag{8a}$$

$$\stackrel{n+1}{\alpha} = \log^{-1}(\stackrel{n}{\alpha}) \in \stackrel{n}{I_1}. \tag{8b}$$

I imagine the elements,  $\overset{n}{\alpha}$ ,  $\overset{n+1}{\alpha}$ , ..., as different scales of a physical degree of freedom. Formally,  $\overset{n}{I_1} \cup \overset{n}{\mathcal{F}} \overset{n}{I_1}$ , is the *ideal* of the *n*-group, but it feels as the horizont of too large/small elements, seen from the perspective of the *n*-scale. In the next slide we identify these abstract elements in the group,  $(\mathbb{R}^+, \times)$ .