

Harmonic Analysis : Exercise Sheet 5 (1)

(i) Let's take $f, g \in \mathbb{C}^N$.

We have to prove that $C(g) + C(f) = C(f) + C(g)$
and $C(g)C(f) = C(f)C(g)$.

$C(g) + C(f) = C(f) + C(g) \rightarrow$ definition of matrix addition.

$C(f)C(g) = C(g)C(f)$

$$C(g) = \begin{pmatrix} f^T T_N g & T_N^2 g & \dots & T_N^{N-1} g \end{pmatrix} = \begin{pmatrix} R_N g^T \\ R_N^2 g^T \\ \vdots \\ R_N^N g^T \end{pmatrix}$$

$$C(f)C(g) = \begin{pmatrix} R_N f^T x g & R_N f^T T_N g & \dots & R_N f^T T_N^{N-1} g \\ R_N^2 f^T x g & & & \\ \vdots & & & \\ R_N^N f^T x g & \dots & \dots & R_N^N f^T T_N^{N-1} g \end{pmatrix} = \sum_{i=0}^{N-1} T_N^i g R_N^{i+1} f^T$$

We have to show that for $i \in \llbracket 1, N \rrbracket$ and $j \in \llbracket 0, N-1 \rrbracket$ $\sum R_N^i f^T T_N^{j+1} g = \sum R_N^i g T_N^{j+1} f^T$.

$$R_N^i(g) = (g_{i-1}, \dots, g_0, g_{N-1}, \dots, g_i)^T \text{ and } T_N^j(g) = (g_{N-j}, g_{N-j-1}, \dots, g_{N-1}, g_0, \dots, g_{N-j+1})$$

$$f \times R_N g^T = \begin{pmatrix} f_0 \\ \vdots \\ f_{N-1} \end{pmatrix} (g_0, g_{N-1}, \dots, g_1) = \begin{pmatrix} f_0 g_0 & f_0 g_{N-1} & \dots & f_0 g_1 \\ f_1 g_0 & & & \\ \vdots & & & \\ f_{N-1} g_0 & \dots & \dots & f_{N-1} g_1 \end{pmatrix}$$

$$g \times R_N f^T = \begin{pmatrix} g_0 & g_{N-1} & \dots & g_1 \\ \vdots & & & \end{pmatrix}$$

??

Units in $f = (1, 0, \dots, 0)$ so that $C(f) = \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & \dots & \dots & 1 \end{pmatrix} = I_N$.
Zeros

(ii)

$$\Delta^s(f * g) = (\Delta^s f) * g = f * (\Delta^s g).$$

$$\Delta f(n) = f(n+1) - f(n).$$

$\forall n \in \mathbb{Z}, \mathbb{N}$: By recurrence.

let say $s=0$

$$\Delta^0(f * g)(n) = (f * g)(n) = \begin{cases} (\Delta^0 f) * g(n) \\ f * (\Delta^0 g)(n) \end{cases}.$$

let suppose it is true for $s=k$, let show it's also true for $s=k+1$.

$$\begin{aligned} \Delta^{k+1}(f * g) &= \Delta(\Delta^k(f * g)) \\ &= \Delta((\Delta^k f) * g) \quad \text{and} = \Delta(f * (\Delta^k g)) \\ &= \Delta(\Delta^k f) * g \quad \text{and} = f * \Delta(\Delta^k g) \\ &= \Delta^{k+1} f * g \quad \text{and} = f * \Delta^{k+1} g. \end{aligned}$$

We then showed by recurrence that $\Delta^s(f * g) = (\Delta^s f) * g = f * \Delta^s g$
 $\forall n \in \mathbb{N}$.

(iii) ?

(5.2) see github link in comment (but I couldn't do the (iii)...) .

(5.3) $\lambda = 8 - 3$.

$$f_{ech} \geq 2 \times f_{max} = 2 \times 1633 = 3266 \text{ Hz and } 2048 < 3266 .$$