

# A simple proof of Kirchoff's theorem, and some other combinatorial graph determinants

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## Abstract

Kirchoff's matrix tree theorem states that the number of spanning trees in a graph is  $\det L'$ , where  $L'$  is the Laplacian matrix  $L$  of the graph, with any row and column deleted. We give a direct proof of the fact that  $\det(xI + L)$  is the generating function of spanning forests with  $k$  roots. Our proof does not rely on the Cauchy-Binet formula, yet is arguably simpler than the standard proof, and applies to directed graphs as well.

The proof is based on a lemma that if  $A$  is the adjacency matrix of a graph where each vertex has at most one outgoing edge, then  $\det(I - A) = 1$  if  $A$  is a forest and 0 otherwise. We also generalize this lemma to any graph, in which case  $\det(I - xA)^{-1}$  is shown to be the generating function of *upward routes*.

Lastly, we generalize Kirchoff's theorem to a theorem about  $\det(A + L)$  where  $A$  is the adjacency matrix of a second graph, which reduces to counting spanning forests when  $A = xI$  and to the fact that  $\det(A)$  counts signed cycle covers when  $L = 0$ . The all-minors matrix tree theorem also follows as a corollary. For instance, the fact that  $\det(L')$  counts spanning forests, where  $L'$  is  $L$  with any column  $i$  and row  $j$  deleted, follows by picking  $A$  to be the matrix where  $A_{ij} = 1$  and zero elsewhere.

## 1 Introduction

Given finite sets of numbers  $S_1, \dots, S_n \subset \mathbb{R}$ , we have the identity:

$$\sum_{x_1 \in S_1, \dots, x_n \in S_n} x_1 x_2 \cdots x_n = \left( \sum_{x_1 \in S_1} x_1 \right) \cdots \left( \sum_{x_n \in S_n} x_n \right)$$

On the left hand side, we get one term for every way of choosing  $(x_1, \dots, x_n) \in S_1 \times \cdots \times S_n$ . This identity is useful in combinatorics, for instance to show that the coefficient of  $x^k$  of  $(x + x^2 + \cdots + x^6)^n$  counts the number of ways of obtaining  $k$  as the sum of  $n$  dice.

Similarly, given finite sets of vectors  $S_1, \dots, S_n \subset \mathbb{R}^k$ , we have the following identity, by multilinearity of the determinant:

$$\sum_{v_1 \in S_1, \dots, v_n \in S_n} \det(v_1, \dots, v_n) = \det \left( \sum_{v_1 \in S_1} v_1, \dots, \sum_{v_n \in S_n} v_n \right)$$

Where  $\det(v_1, \dots, v_n)$  is the determinant of a matrix with those columns. We shall see that Kirchoff's theorem can be obtained from this identity, but we first need some definitions and a lemma.

**Definition 1.1.** We define the concepts *01-graph*, *forest*, *root*, and *tree*:

- A 01-graph is a directed graph where each vertex has 0 or 1 outgoing edges.
- A *forest* is a 01-graph with no cycles.

- The *roots* are vertices with 0 outgoing edges.
- A *tree* is a forest with one root.

If we start at a vertex  $v$  in a 01-graph and keep following the unique outgoing edge, then we either loop in a cycle, or we end up at a vertex with no outgoing edges. Therefore a general 01-graph consists of roots and cycles, plus trees converging onto the roots and cycles.

**Definition 1.2.** The adjacency matrix  $A$  of a graph  $G$  is defined as follows, where  $e_j \in \mathbb{R}^n$  are the basis vectors:

$$A_i = \sum_{(i \rightarrow j) \in G} e_j$$

## 2 The matrix-tree theorem

We start with a lemma that gives us an indicator function for forests.

**Lemma 2.1.** *Let  $A$  be the adjacency matrix of a 01-graph  $G$ , then*

$$\det(I - A) = \begin{cases} 1 & \text{if } G \text{ is a forest} \\ 0 & \text{if } G \text{ has a cycle} \end{cases}$$

*Proof.* In a 01-graph, the vector  $A^i e_k$  follows the path out of  $k$  for  $i = 0, 1, \dots$  (note that this path is unique because each vertex has at most one outgoing edge).

- If  $G$  is forest then  $A^n e_k = 0$  for all  $k$ , where  $n$  is the number of vertices of  $G$ . So  $A^n = 0$ , so all eigenvalues of  $A$  are zero, so all eigenvalues of  $I - A$  are one, so  $\det(I - A) = 1$ .
- If  $G$  has a cycle consisting of vertices  $C$ , take  $v = \sum_{i \in C} e_i$ . Then  $Av = v$ , so  $(I - A)v = 0$ , so  $\det(I - A) = 0$ .

Thus, for 01-graphs,  $\det(I - A)$  indicates whether the graph is a forest or not. □

If using eigenvalues to prove a finitary lemma feels wrong, here is an alternative proof:

*Proof.* In a 01-graph, the  $i$ -th column of  $A$  consists of only the diagonal 1 when  $i$  is a root, and has one other entry equal to  $-1$  if  $i$  is not a root.

We may therefore do a Laplace expansion along the column of a root, which deletes column  $i$  and row  $i$  from the matrix. On the graph side, this corresponds to deleting root  $i$  and all its incoming edges.

Repeating this process, we end up deleting the entire graph if the graph was a forest, in which case the determinant is one<sup>1</sup>. If the graph had a cycle, this we will be left with a matrix that has a 1 and a  $-1$  in each column, so the sum of the rows is 0 so  $\det(I - A) = 0$ . □

**Definition 2.1.** The Laplacian matrix  $L$  of a graph  $G$  is defined as:

$$L_i = \sum_{(i \rightarrow j) \in G} (e_i - e_j)$$

**Theorem 2.2.** (Kirchoff, Tutte) *The determinant  $\det(I + L)$  gives the number of ways to choose a forest as a subgraph of  $G$ .*

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<sup>1</sup>The determinant of a  $0 \times 0$  matrix is 1, but one can also stop deleting when the matrix is  $1 \times 1$ , and use  $\det([1]) = 1$ .

*Proof.* The strategy is to first consider all possible ways of choosing a 01-graph as a subgraph of  $G$ , and then summing  $\det(I - A)$  over those, which counts only the forests. To choose a 01-graph in  $G$ , we choose for each vertex  $i$  whether we make it a root, or whether we give it an outgoing edge from  $G$ . Therefore:

$$\det(I + L) = \det(e_1 + L_1, \dots, e_n + L_n) = \sum_{\text{01-graph } A \subseteq G} \det(I - A) = \sum_{\text{forest } A \subseteq G} 1$$

In the second step, we have expanded the determinant by multilinearity in each column

$$e_i + L_i = e_i + \sum_{(i \rightarrow j) \in G} (e_i - e_j)$$

In each column we either pick  $e_i$  or we pick one of the terms  $(e_i - e_j)$  in the sum over the outgoing edges<sup>2</sup>. The former corresponds to picking  $e_i$  as a root, and the latter corresponds to picking  $i \rightarrow j$  as the outgoing edge for vertex  $i$ .  $\square$

The same proof obtains a weighted version of the theorem. Let  $A$  be a matrix of weights with  $A_{ij}$  corresponding to edge  $i \rightarrow j$ , and define the Laplacian  $L_i = \sum_j A_{ij}(e_i - e_j)$ . Let  $D$  be a diagonal matrix of weights with  $D_{ii}$  corresponding to vertex  $i$  as a root. Let the weight of a forest be

$$w(F) = \prod_{i \in \text{roots}(F)} D_{ii} \prod_{(i \rightarrow j) \in \text{edges}(F)} A_{ij}$$

Then we have the following weighted version of the previous theorem.

**Theorem 2.3.** (Kirchoff, Tutte) *The determinant  $\det(D + L)$  is the weighted sum of all forests on  $n$ -vertices:*

$$\det(D + L) = \sum_{\text{forest } G} w(G)$$

*Proof.* The  $i$ -th column of the matrix is  $(D + L)_i = D_{ii}e_i + \sum_j A_{ij}(e_i - e_j)$ , so

$$\det(D + L) = \det \left( \begin{array}{c|c|c} \begin{matrix} D_{11}e_1 \\ + \\ A_{11}(e_1 - e_1) \\ + \\ A_{12}(e_1 - e_2) \\ \vdots \\ A_{1n}(e_1 - e_n) \end{matrix} & \begin{matrix} D_{22}e_2 \\ + \\ A_{21}(e_2 - e_1) \\ + \\ A_{22}(e_2 - e_2) \\ \vdots \\ A_{2n}(e_2 - e_n) \end{matrix} & \cdots & \begin{matrix} D_{nn}e_n \\ + \\ A_{n1}(e_n - e_1) \\ + \\ A_{n2}(e_n - e_2) \\ \vdots \\ A_{nn}(e_n - e_n) \end{matrix} \end{array} \right) = \sum_{\text{01-graph } G} w(G) \det(I - A_G) = \sum_{\text{forest } G} w(G)$$

In the middle step we expand the determinant by multilinearity: we sum over all possible ways to pick one term in each column. To choose a 01-graph on  $n$  vertices, for each vertex  $i$  (column  $i$ ) we choose whether to make  $i$  a root (term  $D_{ii}e_i$ ) or to give  $i$  an outgoing edge  $i \rightarrow j$  (term  $A_{ij}(e_i - e_j)$ ). After taking the weights  $D_{ii}$  and  $A_{ij}$  out of the determinant, we're left with a sum of  $w(G) \det(I - A_G)$  over all 01-graphs  $G$  on  $n$  vertices, which is equal to the sum of  $w(G)$  over all forests, by the lemma.  $\square$

<sup>2</sup>We do not further expand  $(e_i - e_j)$  into two separate terms.

### 3 Upwards routes

We now know that when  $A$  is the adjacency matrix of a 01-graph, then  $\det(I - A) = 1$  if  $G$  is a forest and  $\det(I - A) = 0$  if  $G$  has a cycle. One naturally wonders about the value of  $\det(I - A)$  when  $A$  is an arbitrary adjacency matrix.

**Definition 3.1.** Given a directed graph  $G$  with an order on the vertices, we define *(strict) upwards loops* and *(strict) upwards routes*:

- An *upwards loop* at vertex  $i$  is a path from  $i$  to  $i$  that does not visit vertices lower than  $i$ .
- A *strictly upwards loop* at vertex  $i$  is a path from  $i$  to  $i$  that only visits vertices higher than  $i$  (except at the start/endpoint of the path, where it does visit  $i$  itself).
- A *(strictly) upwards route* is a choice of (strictly) upwards loop at each vertex.

Let  $f_i(x)$  be the generating function of strictly upwards loops of length  $k$  at vertex  $i$ . Then

$$\bar{f}_i(x) = (1 - f_i(x))^{-1}$$

is the generating function of upwards loops of length  $k$  at vertex  $i$ , because an upwards loop of length  $k$  splits up uniquely into a sequence of strictly upwards loops. Furthermore, the generating functions  $f(x)$  and  $\bar{f}(x)$  of (strictly) upward routes of  $k$  edges are given by:

$$f(x) = \prod_{i=1}^n f_i(x) \qquad \bar{f}(x) = \prod_{i=1}^n \bar{f}_i(x)$$

Recall Cramer's rule:

**Theorem 3.1.** (Cramer's rule) Let  $A$  be a matrix and let  $A_{[i,j]}$  be the same with column  $i$  and row  $j$  deleted, then:

$$A_{ij}^{-1} = \frac{\det(A_{[i,j]})}{\det(A)}$$

From this, we get the following lemma that allows us to calculate  $\det(A)^{-1}$  in terms of entries of inverses of submatrices of  $A$ :

**Lemma 3.2.** Given an invertible matrix  $A$ ,

$$\det(A)^{-1} = \prod_{i=0}^{n-1} (A_{[1\dots i, 1\dots i]})_{11}^{-1}$$

Where  $A_{[1\dots i, 1\dots i]}$  is the matrix  $A$  with the first  $i$  rows and columns deleted.

*Proof.* Cramer's rule implies:

$$\det(A)^{-1} = A_{1,1}^{-1} \cdot \det(A_{[1,1]})^{-1}$$

Continuing this by induction, we get:

$$\begin{aligned} \det(A)^{-1} &= A_{1,1}^{-1} \cdot \det(A_{[1,1]})^{-1} \\ &= A_{1,1}^{-1} \cdot (A_{[1,1]})_{11}^{-1} \cdot \det(A_{[1..2, 1..2]})^{-1} \\ &= \dots \\ &= A_{1,1}^{-1} \cdot (A_{[1,1]})_{11}^{-1} \cdot (A_{[1..2, 1..2]})_{1,1}^{-1} \cdots (A_{[1..n-1, 1..n-1]})_{1,1}^{-1} \cdot 1 \end{aligned}$$

□

We apply this lemma to the matrix  $I - xA$ , to obtain:

**Lemma 3.3.** *The generating function of upwards routes with  $k$  edges is  $\det(I - xA)^{-1}$ .*

*Proof.* Apply the preceding lemma:

$$\det(I - xA)^{-1} = \prod_{i=0}^{n-1} ((I - xA)_{[1\dots i, 1\dots i]})_{11}^{-1}$$

Thus, for each  $i$  we first obtain a subgraph by deleting vertices with lower number than  $i$ , and then  $((I - xA)_{[1\dots i, 1\dots i]})_{11}^{-1}$  is the generating function of loops from vertex  $i$  to  $i$  in the resulting graph. Thus, in terms of the original graph, these are loops that do not visit vertices with lower number than  $i$ . Multiplying this over each vertex  $i$  in the original graph, we obtain the result.  $\square$

**Corollary 3.3.1.** The number of (strictly) upwards routes of  $k$  edges does not depend on the order of the vertices.

*Proof.* If we permute the order of the vertices by a permutation  $P$ , the generating function

$$\det(I - xPAP^{-1}) = \det(P(I - xA)P^{-1}) = \det(I - xA)$$

stays the same.  $\square$

**Lemma 3.4.** *For an arbitrary adjacency matrix  $A$ ,*

$$\det(I - xA) = \prod_{i=1}^n (1 - f_i(x))$$

Where  $f_i(x)$  is the generating function of strictly upwards loops at vertex  $i$ .

*Proof.* We have the following relationship between the generating functions:

$$\det(I - xA) = \bar{f}(x)^{-1} = \prod_{i=1}^n \bar{f}_i(x)^{-1} = \prod_{i=1}^n (1 - f_i(x))$$

$\square$

This is kind of interesting, because  $\det(I - xA)$  is a polynomial, whereas  $f_i(x)$  is a power series, so many terms cancel on the right hand side. The original lemma follows as a corollary, which gives us a third proof of the lemma:

**Corollary 3.4.1.** For a 01-graph,  $\det(I - A) = 1$  if  $G$  is a forest, and 0 if  $G$  has a cycle.

*Proof.* If  $G$  is a forest, then it has no strictly upward loops, so  $f_i(x) = 0$ , so  $\det(I - xA) = 1$ . If  $G$  has a cycle, let  $i$  be the lowest vertex on the cycle. Then  $f_i(x) = x^k$  where  $k$  is the length of the cycle. Now substitute  $x = 1$  to obtain  $\det(I - A) = 0$ .  $\square$

It goes without saying that weighted versions of the preceding lemmas hold too.

## 4 Kirchoff's theorem with cycles

Let  $G_A$  be a graph with adjacency matrix  $A$  and let  $G_L$  be a graph with Laplacian  $L$ . We shall generalize Kirchoff's theorem from  $\det(I + L)$  to  $\det(A + L)$ . In order to do this we need to define 1-graphs.

**Definition 4.1.** A 1-graph is a directed graph where each vertex has exactly one outgoing edge.

Thus, at each vertex we can continue following a unique path indefinitely. In a finite graph that path must eventually cycle. So a general 1-graph looks like a bunch of disjoint cycles and a bunch of trees converging onto those cycles.

In our 1-graphs, some edges will be selected from  $G_A$  and some will be selected from  $G_L$ . We define the weight function:

$$w(F) = \begin{cases} \dots & \dots \\ 0 & \text{otherwise} \end{cases}$$

**Theorem 4.1.** *Kirchoff's theorem with cycles.*

$$\det(A + L) = \sum_{1\text{-graph } F \subseteq (G_A + G_L)} w(F)$$

*Proof.* ... Main idea: generalize the lemma to account for cycles. Each time we Laplace expand a column with one entry (from  $G_A$ ), which is now not necessarily in diagonal position, we obtain a sign. □

Bunch of corollaries:

**Corollary 4.1.1.**  $\det(A)$  is the number of signed cycle covers.

**Corollary 4.1.2.**  $\det(I + L)$  is the number of spanning forests.

**Corollary 4.1.3.**  $\det(L_{[i,j]})$  spanning trees, for all  $i, j$ .

**Corollary 4.1.4.** All-minor matrix tree theorem.

**Corollary 4.1.5.** Undirected matrix tree theorem.