

ARITHMETIC ON CHURCH NUMERALS USING A NOTATIONAL TRICK

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Abstract

Church naturals allow us to represent numbers in pure lambda calculus. In this short note we'll see how to define addition, multiplication, and exponentiation on Church nats using a cute notational trick. We'll also see how to define predecessor and fast growing functions.

1 Addition, multiplication, and exponentiation

Church represents a natural number n as a higher order function, which I'll denote \mathbf{n} . The function \mathbf{n} takes another function f and composes f with itself n times:

$$\mathbf{n} f = \underbrace{f \circ f \cdots \circ f}_{n \text{ times}} = f^n$$

We can convert a Church nat \mathbf{n} back to an ordinary nat by applying it to the ordinary successor function $S : \mathbb{N} \rightarrow \mathbb{N}$ given by $S n = n + 1$: then $\mathbf{n} S 0$ gives us back an ordinary natural number n because $\mathbf{n} S 0$ is the n -fold application of the successor function to the number 0, which just increments it n times.

The first few Church natural numbers are:

$$\begin{aligned} \mathbf{0} &\triangleq \lambda f. \lambda z. z \\ \mathbf{1} &\triangleq \lambda f. \lambda z. f z \\ \mathbf{2} &\triangleq \lambda f. \lambda z. f (f z) \\ \mathbf{3} &\triangleq \lambda f. \lambda z. f (f (f z)) \end{aligned}$$

Many descriptions of Church nats will view them in that way: as a function that takes *two* arguments f and z that computes $f(f(\dots(fz)\dots))$, but this point of view gets incredibly confusing when you try to define arithmetic on them, particularly multiplication and exponentiation. So think about $\mathbf{n} f = f^n$ as performing n -fold function composition.

It will be helpful to introduce an alternative notation for function application:

$$x^f \equiv f(x)$$

This may seem strange at first, but consider that using this notation we can *define* the first few Church natural numbers as:

$$\begin{aligned} f^0 &\triangleq \text{id} \\ f^1 &\triangleq f \\ f^2 &\triangleq f \circ f \\ f^3 &\triangleq f \circ f \circ f \end{aligned}$$

Note that on the left hand side, we are really defining $\mathbf{3}$ as $\mathbf{3}(f) \triangleq f \circ f \circ f$.

The advantage of this notation is apparent when defining addition and multiplication on Church nats:

$$f^{\mathbf{n+m}} \triangleq f^{\mathbf{n}} \circ f^{\mathbf{m}} \qquad f^{\mathbf{n \cdot m}} \triangleq (f^{\mathbf{n}})^{\mathbf{m}}$$

Exponentiation of Church nats is even better: our notation already makes $\mathbf{n}^{\mathbf{m}}$ do the right thing:

$$\mathbf{n}^{\mathbf{m}} \equiv \mathbf{m}(\mathbf{n}) \quad (\text{already does the right thing!})$$

The notation makes the proofs that this does arithmetic correctly a triviality: if $[n]$ is the Church nat corresponding to an ordinary natural number $n \in \mathbb{N}$, then

$$f^{[n]+[m]} = f^{[n]} \circ f^{[m]} = f^n \circ f^m = f^{n+m} = f^{[n+m]}$$

where $f^{[m]} \equiv [m](f)$ according to our notation, and f^n for ordinary natural number $n \in \mathbb{N}$ is n -fold function composition. The proofs for multiplication and exponentiation are analogous.

2 Predecessor

Surprisingly, defining the predecessor on Church nats is the most difficult. I think this is due to Curry.

We define the function $f : \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N} \times \mathbb{N}$:

$$f((a, b)) = (s(a), a)$$

If we start with $(0, x)$ and keep applying f we get the following sequence:

$$(0, x) \rightarrow (1, 0) \rightarrow (2, 1) \rightarrow (3, 2) \rightarrow (4, 3) \rightarrow \dots$$

So

$$\begin{aligned} f^n((0, x))_1 &= n \\ f^n((0, x))_2 &= \begin{cases} x & \text{if } n = 0 \\ n - 1 & \text{if } n > 0 \end{cases} \end{aligned}$$

So we can define the predecessor function:

$$p = \lambda n. n \, f \, (0, 0)$$

So that $p(0) = 0$ and $p(n) = n - 1$ for $n > 0$.

2.1 Pairs

We made use of pairs to define the predecessor, so to use pure lambda calculus we need to define pairs in terms of lambda. We represent a pair (a, b) as:

$$(a, b) = \lambda f. f \, a \, b$$

We can extract the components by passing in the function f :

$$\begin{aligned} \text{fst} &= \lambda x. x \, (\lambda a. \lambda b. a) \\ \text{snd} &= \lambda x. x \, (\lambda a. \lambda b. b) \end{aligned}$$

2.2 Disjoint union

Another way to define the predecessor is with disjoint unions. We take:

$$\text{inl}(a) = \lambda f. \lambda g. f a$$

$$\text{inr}(a) = \lambda f. \lambda g. g a$$

Then we can define:

$$f(\text{inl}(a)) = \text{inr}(a)$$

$$f(\text{inr}(a)) = \text{inr}(s(a))$$

We can do this pattern match on an inl/inr by calling it with the two branches as arguments:

$$f(x) = x (\lambda a. \text{inr}(a)) (\lambda a. \text{inr}(s(a)))$$

And we can define:

$$p(n) = (n f \text{ inl}(0)) (\lambda x. x) (\lambda x. x)$$

3 Fast growing functions

Given any function $g : N \rightarrow N$ we can define a series of ever faster growing functions as follows:

$$f_0(n) = g(n)$$

$$f_{k+1}(n) = f_k^n(n)$$

We can define this function using Church naturals:

$$f_k = k (\lambda f. \lambda n. n f n) g$$

If we take $g = s$ then,

$$f_0(n) = n + 1$$

$$f_1(n) = 2n$$

$$f_2(n) = 2^n \cdot n$$

The function $A(n) = f_n(n)$ grows pretty quickly. We can play the same game again, by putting $g = A$, obtaining a sequence:

$$h_0(n) = A(n)$$

$$h_{k+1}(n) = h_k^n(n)$$

To get a feeling for how fast this grows, consider h_1 :

$$h_1(n) = h_0^n(n)$$

$$= A(A(A(\dots A(A(n)))))$$

$$= A(A(A(\dots A(f_n(n)))))$$

$$= A(A(A(\dots f_{f_n(n)}(f_n(n)))))$$

An expression like $h_3(3)$ gives us a relatively short lambda term that will normalise to a huge term. We might as well start with $g(n) = n^n$ since that's even easier to write using Church naturals:

$$\begin{aligned} g &= \lambda a. a \ a \\ A &= \lambda k. k \ (\lambda f. \lambda n. n f n) \ g \ k \\ h &= \lambda k. k \ (\lambda f. \lambda n. n f n) \ A \ k \\ 3 &= \lambda f. \lambda z. f(f(f \ z)) \\ X &= h \ 3 \end{aligned}$$

You can't write down anything close to the number X even if you were to write a hundred pages of towers of exponentials. Of course, we can continue this game, and define a sequence

$$\begin{aligned} g_0 &= \lambda a. a \ a \\ g_1 &= \lambda k. k \ (\lambda f. \lambda n. n f n) \ g_0 \ k \\ g_2 &= \lambda k. k \ (\lambda f. \lambda n. n f n) \ g_1 \ k \\ &\dots \end{aligned}$$

Which can be generalised as:

$$\begin{aligned} f(g) &= \lambda k. k \ (\lambda f. \lambda n. n f n) \ g \ k \\ g_n &= f^n(g_0) \end{aligned}$$

So we get an even more compact, yet much larger number with:

$$\begin{aligned} f &= \lambda g. \lambda k. k \ (\lambda f. \lambda n. n f n) \ g \ k \\ Y &= (3 \ f) \ (\lambda a. a a) \ 3 \end{aligned}$$

Of course, you can easily define much faster growing functions. But here's a challenge: what's the shortest lambda term that normalises, but takes more than the age of the universe to normalise? Or: what's the largest Church natural you can write down in less than 30 symbols?

Please let me know of any mistakes. I haven't checked for mistakes at all :)