

# Math 321: Relations and functions

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Spring 2021

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These **functions** and **relations** are themselves objects of mathematical study.

# Relations

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The standard in mathematics is to take the extension of a relation as *the* definition of relations.

- A **binary relation** on a set  $A$  is a subset of  $A^2$ , the set of ordered pairs from  $A$ .
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(You can talk about relations between more than two objects, but binary relations are used the most.)

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- $\equiv \text{ mod } 3$  on  $\mathbb{Z}$  (that is,  $a$  and  $b$  have the same remainder when divided by 3)

This is called **equivalence modulo 3**.

# Some properties a relation can have

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- $\star$  is **reflexive** if  $a \star a$  for all  $a \in A$ .
- $\star$  is **symmetric** if  $a \star b$  implies  $b \star a$  for all  $a, b \in A$ .
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# Reflexivity, symmetry, and transitivity

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- (RST) We checked earlier that  $=$  on  $\mathbb{R}$  has all three properties.
- (RT) We checked earlier that  $\leq$  on  $\mathbb{R}$  is reflexive and transitive but not symmetric.

- (T) We checked earlier that  $<$  on  $\mathbb{R}$  is transitive but neither reflexive nor symmetric.
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We have examples for all eight cases, completing the proof of the theorem.

# Closures of a relation

Let  $\star$  be a relation on a set  $A$ . We can add new instances to  $\star$  to make it satisfy these properties.

- The **reflexive closure** of  $\star$  is the smallest reflexive relation on  $A$  which contains  $\star$ , i.e. as a subset.
- The **symmetric closure** of  $\star$  is the smallest symmetric relation on  $A$  which contains  $\star$ .
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Examples:

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- What is the reflexive closure of the empty relation on  $\mathbb{R}$ ?

## Another way to define closures

Let  $\dagger$  be a binary relation on a set  $A$ .

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- Start with  $\dagger_0 = \dagger$ .
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- Then the transitive closure of  $\dagger$  is

$$\begin{aligned}\bar{\dagger} &= \bigcup_{n=0}^{\infty} \dagger_n \\ &= \dagger_0 \cup \dagger_1 \cup \dagger_2 \cup \dots\end{aligned}$$

# Equivalence relations

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- The **rearrangement relation** on finite lists of integers: two lists are equivalent if they are rearrangements of each other.  
This is the equivalence relation used in the statement of the fundamental theorem of arithmetic: when we proved that any two prime factorizations of  $n$  must be the same, what we meant is that the two lists were related in this way.

# Equivalence classes and partitions

Let  $\sim$  be an equivalence relation on  $A$ . Then  $\sim$  partitions  $A$  into equivalence classes.

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Notation: Mathematicians write  $A/\sim$  for the family of  $\sim$ -equivalence classes.