

### More than query: Morel, SQL and the evolution of data languages

Wednesday, October 29, 2025 2:45 PM - 3:30 PM

What is the difference between a query language and a general-purpose programming language? Can SQL be extended to support streaming, incremental computation, data engineering, and general-purpose programming? How well does SQL fit into a modern software engineering workflow, with Git-based version control, CI, refactoring, and AI-assisted coding?

These are all questions that drove the creation of Morel. Morel is a new functional programming language that embeds relational algebra, so it is as powerful as SQL. Morel's compiler, like that of any good SQL planner, generates scalable distributed programs, including federated SQL. But unlike SQL, Morel is Turing-complete, which means that you can solve the whole problem without leaving Morel.

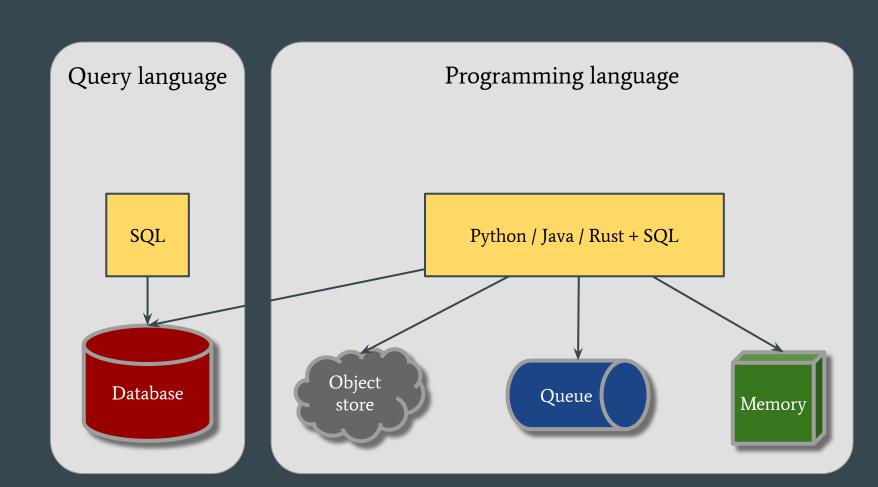
This session will discuss the challenges and opportunities of query languages, especially for streaming and data engineering tasks, and provide a gentle introduction to the Morel language. It is presented by Morel's creator, Julian Hyde, who created Apache Calcite and also pioneered streaming SQL.

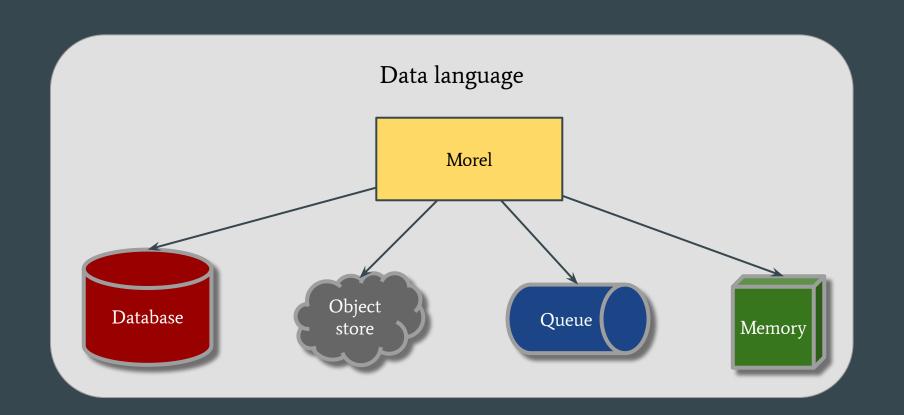
**Location** Breakout Room 285 **Level** Intermediate

Audience Data Engineer/Scientist, Data Streaming Engineers

**Track** Storage and Analytics







#### About me











# Agenda

- 1. Data language
- 2. High level language
- 3. SQL
- 4. A comfortable programming language
- 5. Better query

# 1. Data language

#### 1. Data

```
"Hello, world!"
> val it = "Hello, world!" : string
[1, 2, 3];
> val it = [1,2,3] : int list
(3.14, true);
> val it = (3.14, true) : real * bool
{empno = 100, ename = "SCOTT", job = "MANAGER"};
> val it : {empno:int,ename:string,job:string}
val depts = [
  {deptno = 10, dname = "SALES", emps = []},
  {deptno = 20, dname = "MARKETING", emps = [
    {empno = 100, ename = "SCOTT", job = "MANAGER"},
    {empno = 110, ename = "OATES", job = "CLERK"}
> val depts
  : {deptno:int,dname:string,
      emps:{empno:int, ename:string, job:string} list} list
```

- 1. Data
- 2. Functions

```
fn x => x * x;
> val it = fn : int -> int

map (fn x => x * x) [1, 2, 3, 4];
> val it = [1,4,9,16] : int list

fun factorial 1 = 1
    | factorial n = n * (factorial (n - 1));
> val factorial = fn : int -> int
```

- 1. Data
- 2. Functions
- 3. Types

```
type employee =
  {empno:int, ename:string, is_mgr:bool,
   mgrno:int option};
> type employee
datatype color = BLUE | GREEN | RED;
> datatype color = BLUE | GREEN | RED
datatype 'a option = NONE | SOME of 'a;
> datatype option
SOME "abc";
> val it = SOME "abc" : string option
NONE:
> val it = NONE : 'a option
```

- 1. Data
- 2. Functions
- 3. Types
- 4. Constraints

```
type nat = int check (fn v => v >= 0);
> type nat
type empno = nat;
> type empno
type hr = {
  emps: employee bag check (fn emps =>
    not (exists e in emps
      group e.empno compute count
      where count > 1).
  depts: department bag check (fn depts =>
    not (exists d in depts
      group d.deptno compute count
      where count > 1)
} check (fn hr =>
    not (exists e in hr.emps yield e.deptno
      except distinct
      (from d in hr.depts yield d.deptno)));
> type hr
```

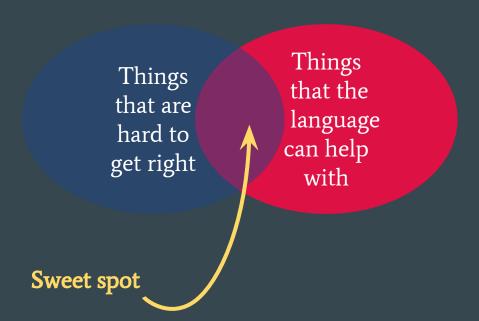
- 1. Data
- 2. Functions
- 3. Types
- 4. Constraints
- 5. Expressions

```
substring ("abcde", 1, 2);
> val it = "bc": string
tl [1, 2, 3];
> [2,3] : list
from i in [1, 2, 3, 4, 5]
 where i \mod 2 = 1
yield i * i;
> [1,9,25]: int list
fun categorize (x, y) =
  case (x mod 2, y mod 2) of
      (0, 0) => "both even"
    | (1, 1) => "both odd"
    | (_, _) => "odd and even";
> val categorize = fn : int * int -> string
```

# 2. High level language

# The "sweet spot" of language design

- 1. Organization of data
- 2. Choice of algorithms
- 3. Parallelism/federation
- 4. Mutation
- 5. Evolving schema
- 6. Abstraction



#### high-level programming language n.

A programming language that requires you to specify only the details that matter.

#### morel n.

- 1. An edible mushroom with a honeycomb-like cap.
- 2. (*comp. sci.*) A high-level data programming language that combines Standard ML with relational algebra.

Part 1: Java ArrayList

```
record LogEntry(String timestamp, String message) {}
List<LogEntry> logsList = List.of(
    new LogEntry("2025-10-25T08:15:23", "User login: alice"),
    new LogEntry("2025-10-25T08:16:45", "API call: /users"),
    new LogEntry("2025-10-25T09:23:11", "Error: timeout"),
    new LogEntry("2025-10-25T10:05:33", "User login: bob"),
    new LogEntry("2025-10-25T14:22:01", "API call: /orders")
    // ... millions of log entries
List<String> logsInRange(String startTime,
   String endTime, List<LogEntry> logs) {
 var result = new ArrayList<String>();
  for (var entry : logs) {
    if (entry.timestamp.compareTo(startTime) >= 0
        && entry.timestamp.compareTo(endTime) <= 0) {
      result.add(entry.message);
  return result:
logsInRange("2025-10-2509:00:00", "2025-10-2511:00:00", logsList);
```

Part 2: Java SortedMap

```
SortedMap<String, String> logsMap =
    new TreeMap<>(
        Map.ofEntries(
            Map.entry("2025-10-25T08:15:23", "User login: alice"),
            Map.entry("2025-10-25T08:16:45", "API call: /users"),
            Map.entry("2025-10-25T09:23:11", "Error: timeout"),
            Map.entry("2025-10-25T10:05:33", "User login: bob"),
            Map.entry("2025-10-25T14:22:01", "API call: /orders")
           // ...
        ));
List<String> logsInRange(String startTime,
    String endTime, SortedMap<String, String> logs) {
  var result = new ArrayList<String>();
  var subMap = logs.subMap(startTime, true, endTime, true);
  for (var message : subMap.values()) {
    result.add(message);
  return result;
logsInRange("2025-10-2509:00:00", "2025-10-2511:00:00", logsMap);
```

Part 3: Rust BTreeMap

```
let mut logs_map = BTreeMap::new();
logs_map.insert("2025-10-25T08:15:23".to_string()
    "User login: alice".to_string()
logs_map.insert("2025-14-75148") 10-45
                                      .to_string(
    "API call: /users".to_string()
logs_map.insert("2025-14-75140-73:
                                      .to_string(
    "Error: timeout".to_string()
logs_map.insert("2025-10-25110:05:33".to_string()
    "User login: bob".to_string());
logs_map.insert("2025-12-2
                                      .to string()
    "API call: /orders".to_string()
fn logs_in_range(start_time: & tr, end_time: & tr,
    logs: & TreeMap<String, String>) -> Vec<String> {
    let mut result = Vec::new();
    for (_timestamp, message)
      in logs.range(start_time_to_string()..=end_time.to_string()
        result.push(message.clone());
  result
logs_in_range("2025-10-2509:00:00", "2025-10-2511:00:00", logsList);
```

Part 4: Morel list

```
val logsList = [
  ("2025-10-25T08:15:23", "User login: alice"),
  ("2025-10-25T08:16:45", "API call: /users"),
  ("2025-10-25T09:23:11", "Error: timeout"),
  ("2025-10-25T10:05:33", "User login: bob"),
  ("2025-10-25T14:22:01", "API call: /orders")
  (* ... millions of log entries *)
fun logsInRange (startTime, endTime, logs) =
  from (timestamp, message) in logs
    where timestamp >= startTime
      andalso timestamp <= endTime</pre>
    yield message;
logsInRange ("2025-10-2509:00:00", "2025-10-2511:00:00", logsList);
```

# The "sweet spot" of language design

- 1. Organization of data
- 2. Choice of algorithms
- 3. Parallelism/federation
- 4. Mutation
- 5. Evolving schema
- 6. Abstraction

Things that are hard to get right

Things that the language can help with 3. SQL



# Algebra – a collection of values, operators, and laws

#### On numbers

$$a + (b + c) \rightarrow (a + b) + c$$

$$(a * b) + (a * c) \rightarrow a * (b + c)$$

#### On relations

$$P \cup (Q \cup R) \rightarrow (P \cup Q) \cup R$$

$$(P \bowtie Q) \cup (P \bowtie R) \rightarrow P \bowtie (Q \cup R)$$

# Relational algebra in a functional programming language

#### Relational algebra

```
∪ union
\ minus
∩ intersect
σ filter
Π project
⋈ join
```

#### Relational operators as functions

#### SQL

Built-in operators

Relation is the only bulk type

Algebraic rewrite

Not Turing complete

Distributed engine

# SQL + functional

Immutable data

Type inference

Inlining

#### Morel

User-defined operators & types

Algebraic rewrite

Turing complete

Choose your own engine

#### **Functional**

User-defined operators

User-defined types

Functions as values

Local optimizations

Turing complete

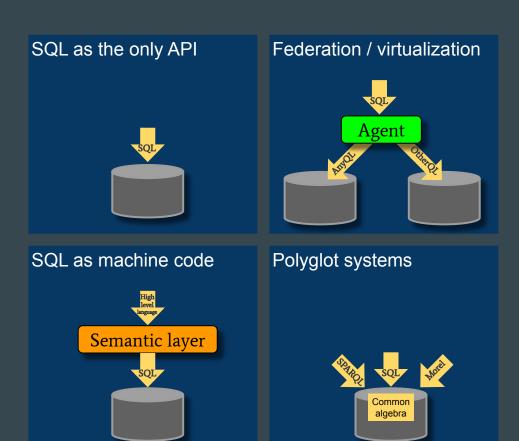
Local execution

# Future of SQL

## **Future of SQL**

#### I believe:

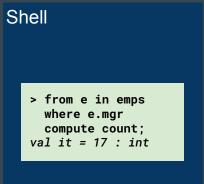
- SQL continues to grow/sprawl
- ISO standard declines in importance
- Databases speak more than one language
- Databases are not the only systems that speak SQL

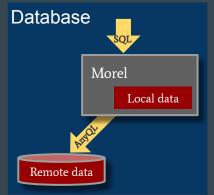


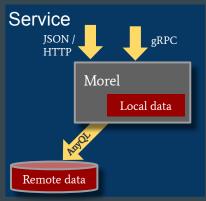
# Data systems

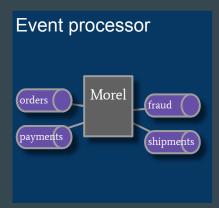
## Data systems

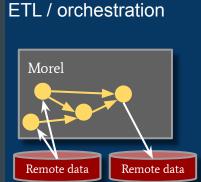


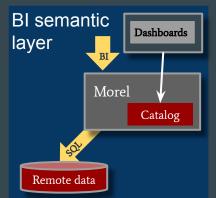


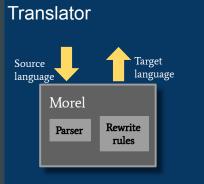












4. A comfortable programming language for your data



# Many nice things about programming languages

If it compiles, it probably works

Refactoring

Autocompletion

Git

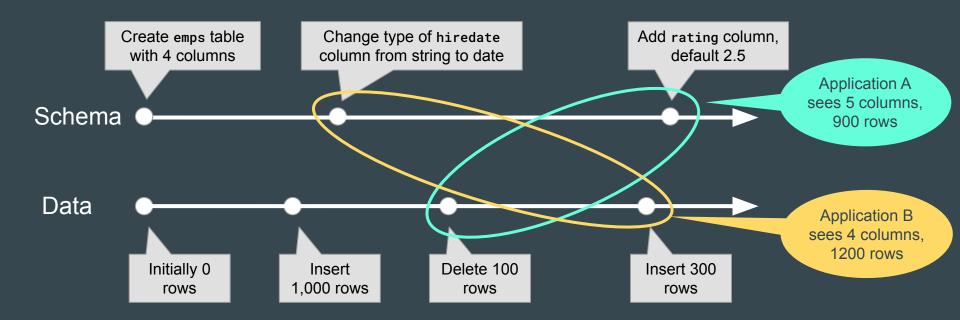
Documentation in the code

Unit tests in the same language

Modules

Abstraction

# Types and data evolve independently



# Database as a module (shims for schema evolution)

```
(* Initial database value and type (schema). *)
val scott1 = db
  : {emps: {name: string, empno: int, deptno: int,
            hiredate: string} bag,
    depts: {deptno: int, name: string} bag};
(* Shim that makes a v1 database look like v2. *)
fun scott2on1shim scott1 =
  \{emps =
   fn () => from e in scott1.emps
      yield {e with hiredate = Date.fromString(e.hiredate)},
   depts = fn () => scott1.depts};
(* Shim that makes v3 database look like v1. *)
fun scott1on3shim scott3 =
  \{emps =
   fn () => from e in scott3.emps
      yield {e with hiredate = Date.toString(e.hiredate)
               removing rating},
   depts = fn () => scott3.depts};
(* An application writes its queries & views against version 2;
   shims make it work on any actual version. *)
val scott = scott2;
fun recentHires () =
 from e in scott.emps
    where e.hiredate > Date.subtract(Date.now(), 100);
```

# 5. Better query



```
SQL
SELECT item, COUNT(*) AS c,
   SUM(sales) AS total
FROM ProduceSales
WHERE item != 'bananas'
   AND category IN ('fruit', 'nut')
GROUP BY item
ORDER BY item DESC;
item   c total
===== ==== =====
apples   2   9
```

# SQL

# GoogleSQL pipe syntax

```
FROM ProduceSales
|> WHERE item != 'bananas'
    AND category IN ('fruit', 'nut')
|> AGGREGATE COUNT(*) AS c, SUM(sales) AS total
    GROUP BY item
|> ORDER BY item DESC;
```

```
SQL
SELECT item, COUNT(*) AS c,
   SUM(sales) AS total
FROM ProduceSales
WHERE item != 'bananas'
   AND category IN ('fruit', 'nut')
GROUP BY item
ORDER BY item DESC;
item c total
===== ==== ==== apples 2 9
```

# SQL

# PRQL

```
# PRQL
from produceSales
filter item != "bananas"
filter category in ["fruit", "nut"]
group item (
   aggregate {
     c = count this,
     total = sum sales
   }
)
sort -item
```

# SQL



```
(* Morel *)
from p in produceSales
  where p.item <> "bananas"
     andalso p.category elem ["fruit", "nut"]
  group p.item compute {c = count,
     total = sum over p.sales}
  order DESC item
```

### Morel query syntax

```
exp →
    from scan [, scan... ] [ step... ]
     [ terminalStep ]
  exists scan [, scan...] [ step...]
   | forall scan [, scan...] [ step...]
     require condition
   | aggFn over exp<sup>1</sup>
   \mathsf{l} elements^1
   | current<sup>2</sup>
   ∣ ordinal<sup>2</sup>
    DESC exp
  (* many other kinds of expression *)
scan →
    pat in collection [ on condition ]
  | pat = exp [ on condition ]
  l var
<sup>1</sup> Aggregate expression and elements only in compute
<sup>2</sup> current and ordinal only in query
```

```
step →
    join scan [, scan...]
   where condition
   yield exp
   yieldmany collection
   distinct
   group exp [ compute exp ]
   order exp
   unorder
   skip number
   take number
   through pat in exp
   union [ distinct ] collection [, collection... ]
    intersect [ distinct ] collection [, collection... ]
   except [ distinct ] collection [, collection... ]
terminalStep →
    into exp
   compute exp
```

## **How Morel improves on SQL's SELECT**

- Ordered collections
- Collections of non-records
- Queries that read like pipelines
- Function values (lambdas)
- Easily refactor query into functions
- Write your own aggregate functions
- Any expression can be correlated
- Use a query anywhere that an expression is allowed
- Define variables, intermediate results or functions anywhere
- Polymorphic types
- Sum types
- 2-valued boolean logic

The billion-dollar mistake

```
SELECT e.ename, e.sal
FROM emps AS e
WHERE e.deptno = 10
AND e.sal >
    (SELECT MAX(e2.sal)
    FROM emps AS e2
    WHERE e2.deptno = 20
    AND e2.job = 'PROGRAMMER')
```

### The billion-dollar mistake

### Why do we need NULL?

- Empty strings
- Missing references
- Missing values
- N/A
- Unknown boolean
- Aggregation over the empty set

### We don't need NULL!

- Empty strings
- Option type
- Option type
- Option type
- Unnecessary
- Aggregate functions that return a default value, or an option

```
SELECT e.ename, e.sal
FROM emps AS e
WHERE e.deptno = 10
AND e.sal >
    (SELECT MAX(e2.sal)
    FROM emps AS e2
    WHERE e2.deptno = 20
    AND e2.job = 'PROGRAMMER')
```

```
from e in scott.emp
  where e.deptno = 10
    andalso
    (forall e2 in scott.emp
      where e2.deptno = 20
        andalso e2.job = "PROGRAMMER"
      require e.sal > e2.sal)
    yield e.ename
```

**Modifying data** 

### SQL DML

```
-- Delete employees who earn more than 1,000.
DELETE FROM hr.emps
WHERE sal > 1000;
-- Add one employee.
INSERT INTO hr.emps (empno, deptno, ename, job, sal)
VALUES (100, 20, 'HYDE', 'ANALYST', 1150);
-- Double the salary of all managers.
UPDATE hr.emps
SET sal = sal * 2
WHERE job = 'MANAGER';
-- Commit.
COMMIT;
```

## Morel DML (first attempt)

```
(* Delete employees who earn more than 1,000. *)
delete e in hr.emps
 where e.sal > 1000;
(* Add one employee. *)
insert hr.emps
  [\{empno = 100, deptno = 20, ename = "HYDE", \}
    job = "ANALYST", sal = 1150}];
(* Double the salary of all managers. *)
update e in hr.emps
 where e.job = "MANAGER"
 assign (e, {e with sal = e.sal * 2});
(* Commit. *)
commit;
```

### Morel DML

```
(* Delete employees who earn more than 1,000. *)
val emps2 =
 from e in hr.emps
   where not (e.sal > 1000);
(* Add one employee. *)
val emps3 = emps2 union
  [{empno = 100, deptno = 20, ename = "HYDE", job = "ANALYST",
   sal = 1150};
(* Double the salary of all managers. *)
val emps4 =
  from e in emps3
   yield if e.job = "MANAGER"
      then {e with sal = e.sal * 2}
      else e;
(* Commit. *)
commit {hr with emps = emps4};
```

### Incremental computation

```
(* New and removed employees. *)
val empsAdded = emps4 except hr.emps;
val empsRemoved = hr.emps except emps4;
(* Compute the updated summary table. *)
val summary2 =
 from s in hr.summary
    union
      (from e in empsAdded
        yield {e.deptno, c = 1, sum_sal = e.sal})
    union
      (from e in empsRemoved
        yield {e.deptno, c = \sim 1, sum_sal = \sim e.sal})
    group s.deptno compute c = sum of c, sum_sal = sum of sum_sal
    where c <> 0:
(* Commit. *)
commit {hr with summary = summary2};
```

## Optimizing data-intensive programs

Query optimization (e.g. change join order, push down filters, parallelize)

Program optimization (e.g. remove unused variables, unwind loops, convert tail-recursion to loop)

Efficiently compute deltas

Recommend materialized views

Multi-query optimization [CALCITE-6188]

Add temporary tables for shared intermediate results

Remove unused temporary tables

