Advanced Enterprise Computing -Lecturenotes SoSe2016

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1 Replication and State Management (25.04. - 09.05.)

1.1 Motivation and Background

1.1.1 Replication

Definition - Replication Process of maintaining multiple Copies of an Entity (Data / Process / File ...)

Advantages of Replication in General

- System Availability / Fault tolerance in case
 - A Server fails
 - B Data is corrputed.
- Performance / Scalability
 - A Workloads are spread across distributed Replicas
 - B Geodistribution for processing demands in client's proximity

Disadvantages of Replication in General

• Consistency vs. Performance

Kinds of Replication In general there are the following kinds of "physical" Replication. We do only consider (B).
!!!!!!!!!!!!ueberpruefen am Ende!!!!!!!!!!!

Replication Strategies PAGE 17

Synchronous vs. Asynchronous

Synchronous / eager

Asynchronous / lazy

Primary Copy vs. Update Everywhere

Primary Copy / master

Update Everywhere / group

Name explanation Protokolltitle implementation Atomacity either perform Transaction entirely or roll-back atomic commitment protocol 2PC

ACID

Atomiticity:

Consistency: does not mean Data-Consistency but that the transaction produces consistent changes.

Isolation: Transactions are isolated from one another

Durability: Once the transaction is ready (commits) it remains.

Both the Atomitcity and the Isolation are managed by the **Transaction** Manager It acquires locks on behalf of all transactions and tries to come up with a serializable execution, that is, make it look like the transactions were executed one after the other. If the transactions follow 2 Phase Locking, serializability is guaranteed. Thus, the scheduler only needs to enforce 2PL behavior.

What happens to ACID in case of Replication? Atomicity can be guaranteed using 2PC (but expensive) Problem: Serialization order must be the same at all replicas.

Synchronous ACID properties apply to all copy updates

- 1.2 Managing Replication
- 1.3 Implications of Replication
- 1.4 Paxos and CRDTs
- 2 Prototyping
- 3 Experiments
- 4 DevOps and Microservices
- 5 Reading Assignment

Tabelle 1: My caption

	Procedure 1: My caption	Advantages / Disadvantages
		- ACID (no
		Inconsistencies)
		- High response time
	1 propagate Data	(high execution time,
Synchronous	to everybody	response time)
		- Availability
	2 Wait until	(in case one Copy fails)
	everybody responded	- Response Time
		- Availability
	3 commit	- Data inconsistency
	1 Update local copy	(local read does not always
Asynchronous		return the latest value)
	2 commit	- No guarantee that the changes
		arrive at each copy
	3 Propagato Data	- Replication is not guaranteed
Drimary Cony	3 Propagate Data one Primary copy and	
Primary Copy	several read-only copies	
Undata Everywhere	Each site is able to	
Update Everywhere	initiate changes	

6 Lecturenotes

Lecture $05?_{start}$ @ 81 Für Donnerstag paper mitbringen und Paxos anschauen. 2016-05-09

Paxos (Represent as State-machine) - P. 77

Proposer

Phase 1 - Proposer choses Number larger than any value chosen before by Propposer. - Broadcast the integer prepare(n), e.g. prepare(50)

Acceptors a) Not respond at all b) recject Reject, in case a higher value has been accepted. 50; something b) prommise(n) in case 50; everything. Also Send everything that has already been accepted.

If prposer receives majority of prommise respons, -; proceed to Phase 2 ELSE -; Phase 1

Phase 2 - Check whether any ¡n, value¿ have been returned. - YES: take max n's value - accept (n, value)

Xtensions Paxos Multi-paxos Determine Leader once Stay in phase 2, attatch the leader identifier Leader is the one to accept values

Purpose: Optimize Speed (get rid of the first phase, Master-Slave setup)

Fast Paxos

Generalized Paxos - Assumption: The execution order does not matter.

CRDT Conflict free / Communitive replicated Datatypes

Some operations are commutative, others not.

State- Based vs. Operation based.

theoretically it is possible to converge them but ... practice

IDEA INTEGER - example: e.g. not store int values but operations (increment / decrement))

SET - example

State - based Set

7 Begriffe und Abkürzungen

Replication Strategy to maintain mutiple copies of an entity on multiple Servers.

Replica

CRDT conflict-free replicated data

Paxos

Commit In case a Transaction commits, it is ready.

Concurrency control protocol guarantees isolation of Transactions

2PL Two phase locking (one concurrency control protocol)

Snapshot Isolation other cuncurrency control protocol implementation atomic committeent protocol guarantees atomaticity

2PC Two phase Commit

Transaction Manager Middleware Component; Manages Atomacity and Isolation of Transactions

ACID Atomacity + Consistency + Isolation + Durability