

Date: May 20, 24

Day: \_\_\_\_\_

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Subject : Analysis & Algorithm

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## Assignment # 02

### Quick Sort

Array = [13, 22, 23, 48, 70]

Sol: 

$p(\text{low})$				$q(\text{high})$
13	22	23	48	70
index: 0	1	2	3	4

Pivot element  $P = 13$

$r = 13$  at 0 index

$q = \text{Partition}(A, 0, 4)$

$\because P = \text{pivot element}$

$\because q = \text{high index}$

$\because r = \text{low index}$

$P = 13$  at 0 index

$P = A[r] \rightarrow P = A[0]$

$i = q + 1$

$i = 4 + 1 \Rightarrow 5$

$j = 4$  to 1

$j = 4, i = 5$

if  $13 > 22$

$i = 4$

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Similarly,

$$\begin{array}{ll} j=3 & , \quad i=3 \\ j=2 & , \quad i=2 \\ j=1 & , \quad i=1 \end{array}$$

Swap  $A[i-1]$  with  $A[r]$   
 $A[0]$  with  $A[0] \rightarrow$  no change

Now Quick Sort ( $A, 0, -1$ )QS ( $A, 1, 4$ ) $P=22 \rightarrow r$  is at index 1 $i=5$  $j=4$  to 2 $j=4$  ,  $i=5$ if  $23 > 22$  $i=4$ 

Similarly,

 $j=3$  ,  $i=3$  $j=2$  ,  $i=2$ Swap  $A[2-1]$  with  $A[r]$  $A[1]$  with  $A[1] \rightarrow$  no change

Now,

Quick Sort ( $A, 1, 0$ )QS ( $A, 2, 4$ ) $i=5$  ,  $P=23$  $j=4$  ,  $i=5$ if  $48 > 23$  $i=4$ 

Similarly,

 $j=3$  ,  $i=3$ swap  $A[3-1]$  with  $A[r]$  $A[2]$  with  $A[2] \rightarrow$  no change

Quick Sort (A, 2, 1)

QS (A, 3, 4)

P = 48

i = 3

j = 4 to 4

j = 4

if 70 > 48

i = 4

swap A[i-1] with A[p]

A[3] with A[3] → no change

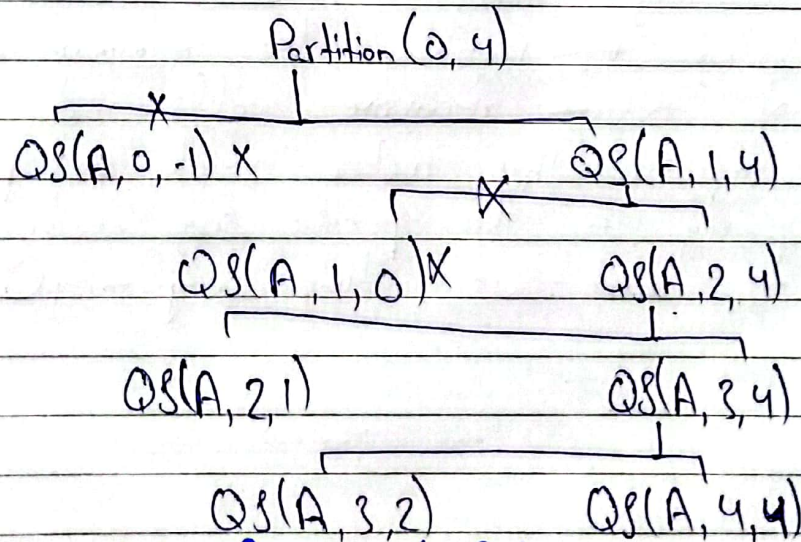
Now

QS (A, 3, 2)

QS (A, 4, 4)

So, final Sorted Array = 13 22 23 48 70

Tree:



Pseudo Code for Quick Sort

(Left to Right)

QS(A, r, q)

if (r < q)

Partition (A, r, q)

QS(A, r, P-1)

QS(A, P+1, q)



Partition ( $A, r, q$ )

Pivot ( $P$ ) =  $A[r]$

$i = q + 1$

for  $j = q$  down to  $r + 1$

if  $A[j] \geq \text{Pivot}$

$i = i + 1$

swap  $A[i]$  with  $A[j]$

~~swap~~ swap  $A[i-1]$  with  $A[r]$

return  $i - 1$

### Sorted Array Worst Case:

The ~~worst~~ worst case scenario occurs because each partitioning step only removes one element and the remaining array is still essentially sorted ~~then~~ requiring the maximum number of ~~recursion~~ recursive calls and comparisons. This makes quick sort perform poorly in this specific case leading to its worst-case time complexity  $O(n^2)$ .

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