

DATA STRUCTURES AND ALGORITHMS

Lecture: Array Sorting

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SORTING



- Sorting takes an unordered collection and makes its an ordered one.
- Let A be a list of n elements A_1 , A_2 , ... A_n in memory.
- Sorting A refers to the operation of rearranging the contents of A
 - Ascending order
 - Descending order





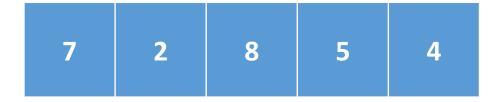


- In this lecture, we are going to look at three simple sorting techniques:
 - Bubble Sort
 - Selection Sort
 - Insertion Sort





- Traverse a collection of elements.
- Move an element from its position to the front (if smallest) or to the end (if largest) in case of ascending order and vice versa.
- "Bubble" the value to the front or end using the operations
 - Pair-wise comparison
 - Swapping



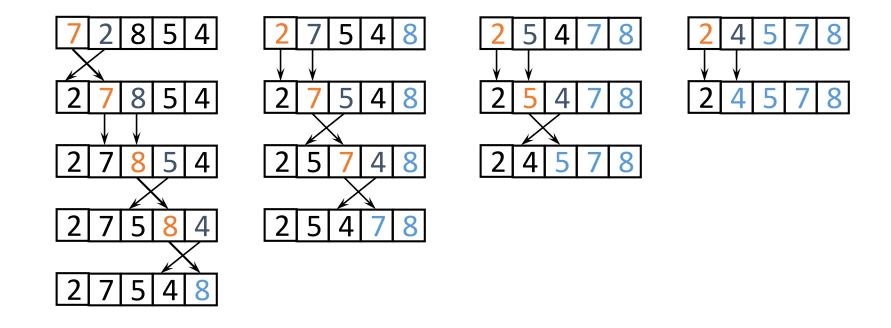




- Compare each element (except the last one) with its neighbor to the right.
 - If they are out of order, swap them
 - This puts the largest element at the very end
 - The last element is now in the correct and final place
- Compare each element (except the last two) with its neighbor to the right.
 - If they are out of order, swap them
 - This puts the second largest element next to last
 - The last two elements are now in their correct and final places
- Compare each element (except the last three) with its neighbor to the right
 - Continue as above until you have no unsorted elements on the left

BUBBLE SORT









```
void BubbleSort ( int arr[], int size )
int temp;
for ( int outer = size-1; outer > 0; outer-- )
    for ( int inner = 0; inner < outer; inner++ )</pre>
       if ( arr[inner] > arr[inner+1] )
           temp = arr[inner];
           arr[inner] = arr[inner+1];
           arr[inner+1] = temp;
```





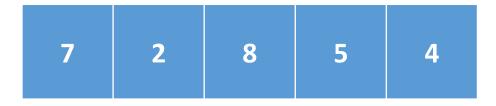
```
for ( int outer = size-1; outer > 0; outer-- ) {
for ( int inner = 0; inner < outer; inner++ ) {
   if ( arr[inner] > arr[inner+1] ) {
         //code for swapping is skipped
   } } }
```

- Let size = arr.length = size of the array
- The outer loop is executed n-1 times (call it n, that's close enough)
- Each time the outer loop is executed, the inner loop is executed
 - Inner loop executes n-1 times at first, linearly dropping to just once
 - On average, inner loop executes about n/2 times for each execution of the outer loop
 - In the inner loop, the comparison is always done (constant time), the swap might be done (also constant time)
- Result is n * n/2 * k, that is, $O(n^2/2 + k) = O(n^2)$





- Define the entire array as the unsorted portion of the array.
- While the unsorted portion of the array has more than one element:
 - Find its largest element
 - Swap with last element (assuming their values are different)
 - Reduce the size of the unsorted portion of the array by 1.







- Given an array of length n,
 - Search elements 0 through n-1 and select the smallest
 - Swap it with the element in location 0
 - Search elements 1 through n-1 and select the smallest
 - Swap it with the element in location 1
 - Search elements 2 through n-1 and select the smallest
 - Swap it with the element in location 2
 - Continue in this fashion until there's nothing left to search

SELECTION SORT



```
8 5 4
```

```
void SelectionSort ( int arr[], int size )
 int temp;
 int min;
 for ( int outer = 0; outer < size-1; outer++ )</pre>
     min = outer;
      for ( int inner = outer+1; inner < size; inner++ )</pre>
          if ( arr[inner] < arr[min] )</pre>
              min = inner;
      temp = arr[outer];
      arr[outer] = arr[min];
      arr[min] = temp;
```





- The Selection Sort might swap an array element with itself--this is harmless, and not worth checking for
- Analysis:
 - The outer loop executes n-1 times
 - The inner loop executes about n/2 times on average (from n to 2 times)
 - Work done in the inner loop is constant (swap two array elements)
 - Time required is roughly (n-1)*(n/2)
- You should recognize this as $O(n^2)$





- The list is divided into two parts: sorted and unsorted.
- In each pass, the following steps are performed:
 - First element of the unsorted part (i.e., sub-list) is picked up
 - Transferred to the sorted sub-list
 - Inserted at the appropriate place
- A list of n elements will take at most n-1 passes to sort the data

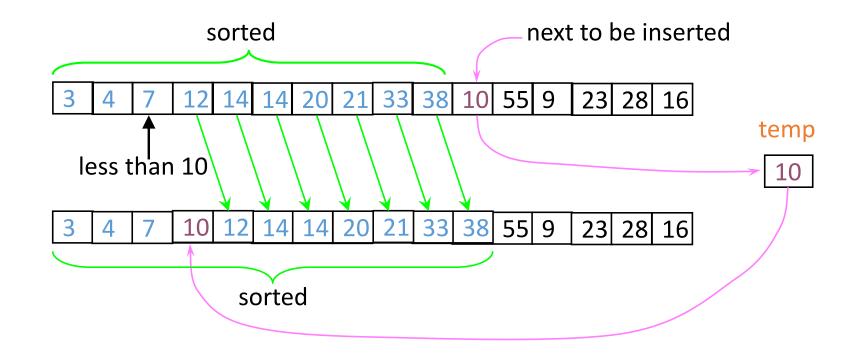




```
void InsertionSort ( int arr[], int size )
int temp;
for ( int outer = 1; outer < size; outer++ )</pre>
   temp = arr[outer];
   int inner = outer;
   while ( inner > 0 && arr[inner-1] >= temp )
       arr[inner] = arr[inner-1];
       inner--;
   arr[inner] = temp;
```











- We run once through the outer loop, inserting each of n elements;
 this is a factor of n
- On average, there are n/2 elements already sorted
 - The inner loop looks at (and moves) half of these
 - This gives a second factor of n/2
- Hence, the time required for an insertion sort of an array of $\frac{n}{2}$ elements is proportional to $\frac{n^2}{4}$
- Discarding constants, we find that insertion sort is $O(n^2)$

CONCLUSION

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- In this lecture we have studied:
 - Bubble Sort
 - Selection Sort
 - Insertion Sort

Question?