Efficient Memory Disaggregation with Infiniswap

Juncheng Gu, Youngmoon Lee, Yiwen Zhang, Mosharaf Chowdhury, Kang G. Shin



Agenda

Motivation and related work

Design and system overview

Implementation and evaluation

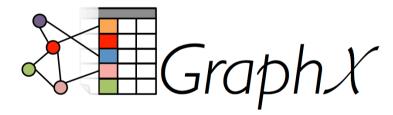
Future work and conclusion

Memory-intensive applications



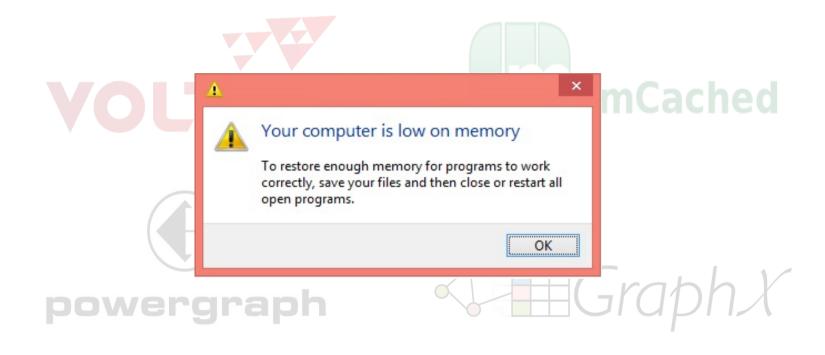






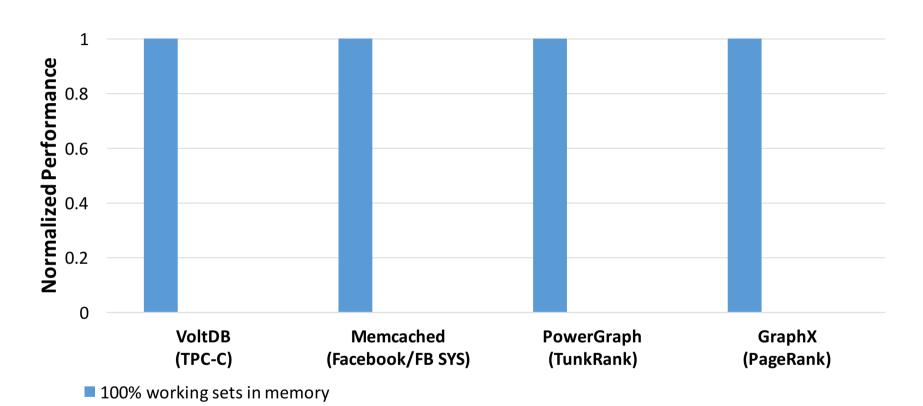
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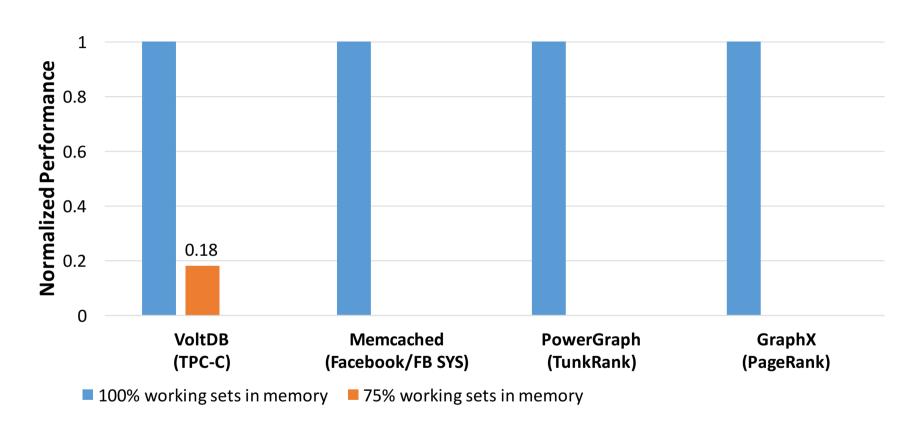
Memory-intensive applications



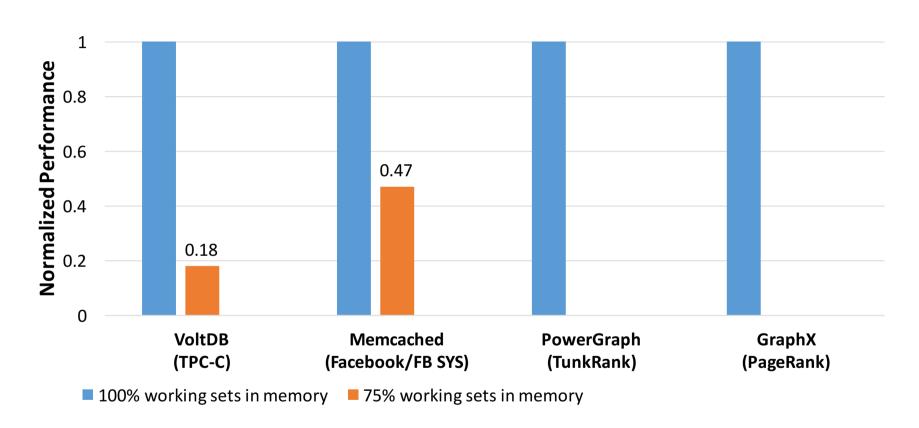
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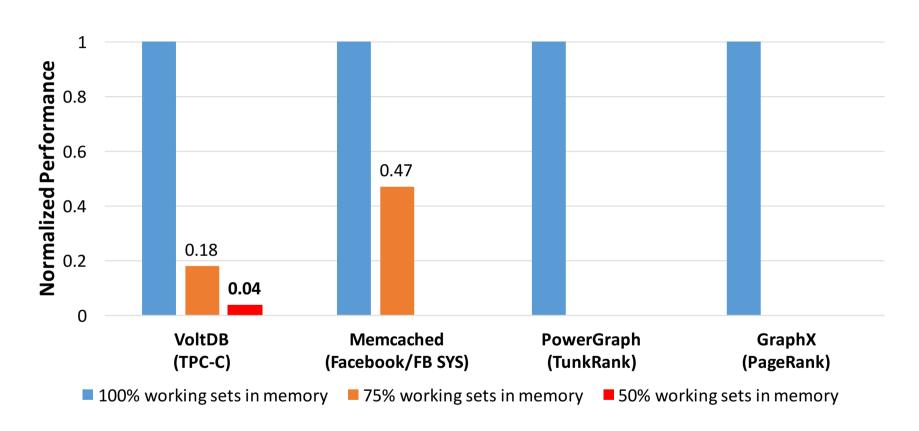
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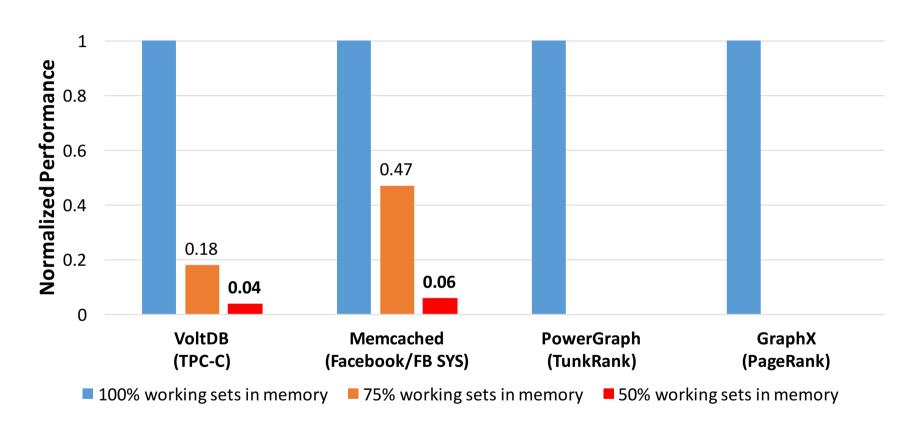


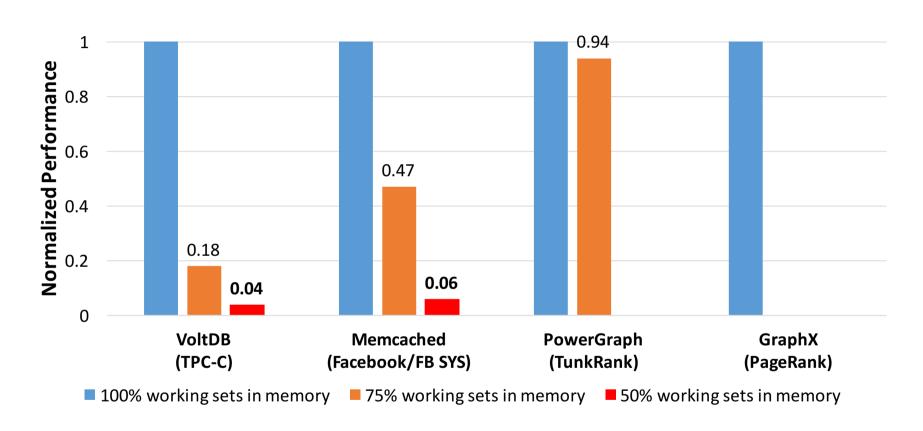


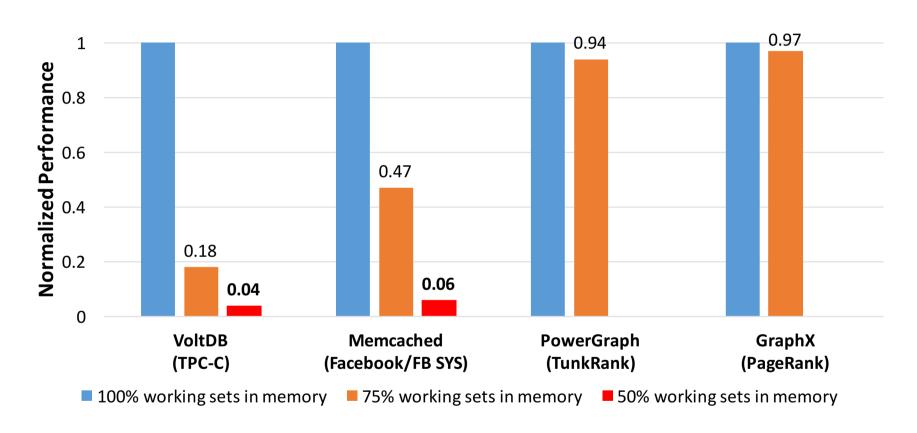
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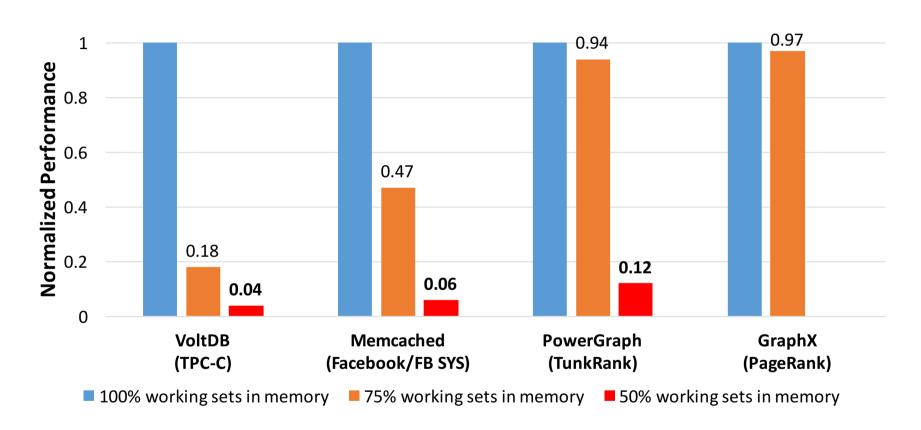


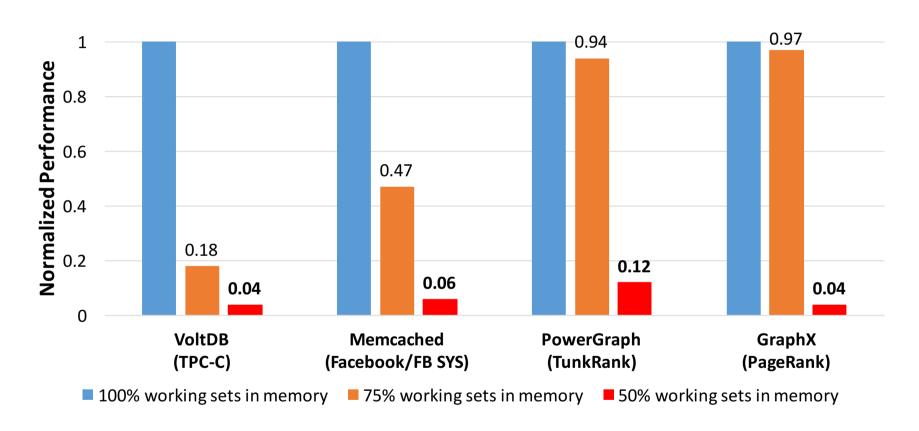


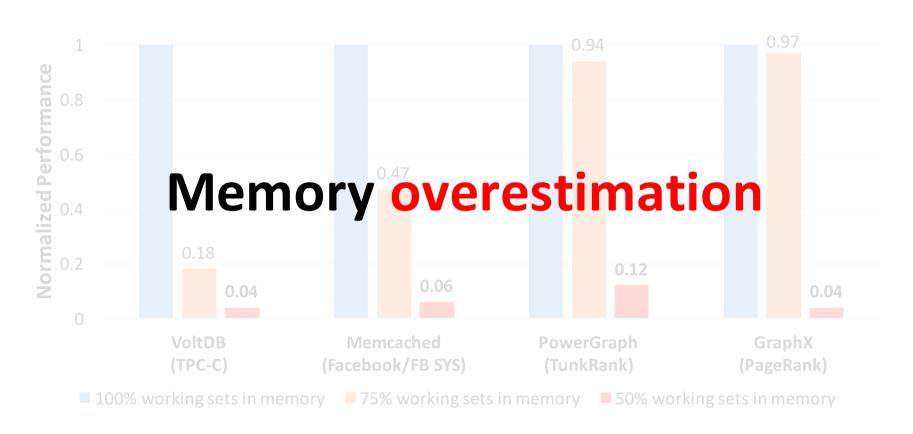




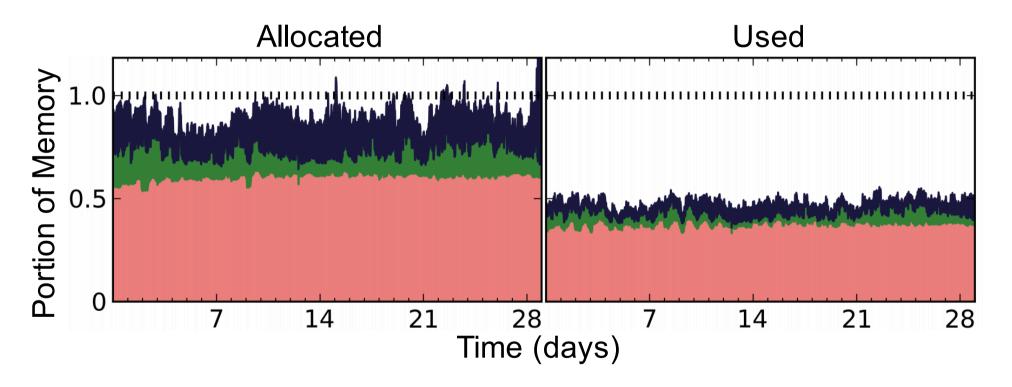




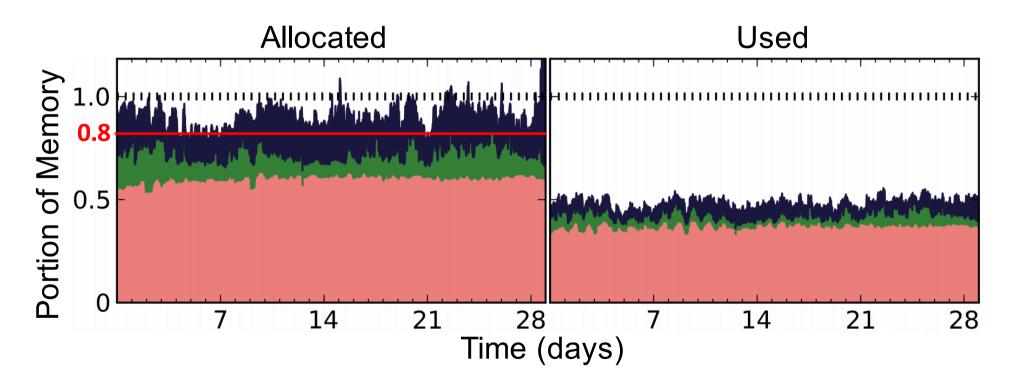




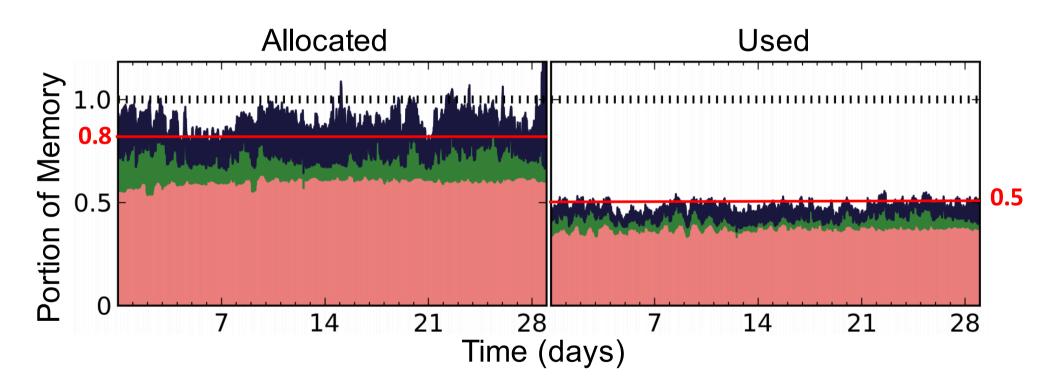
• Google Cluster Analysis_[1]



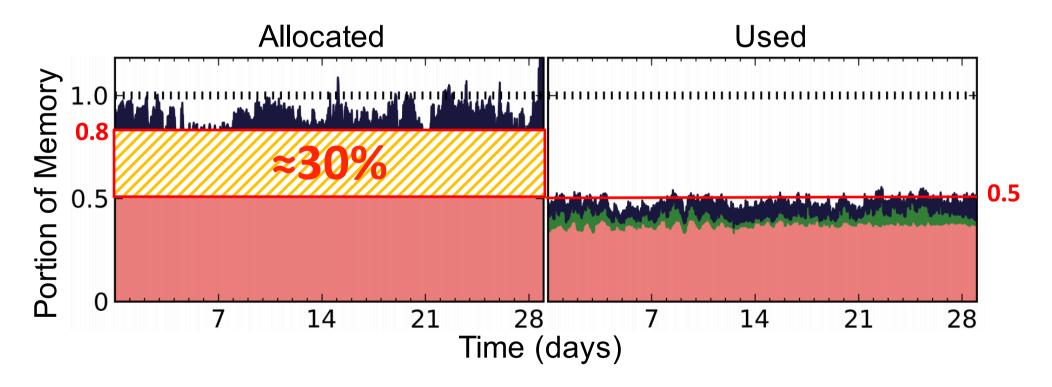
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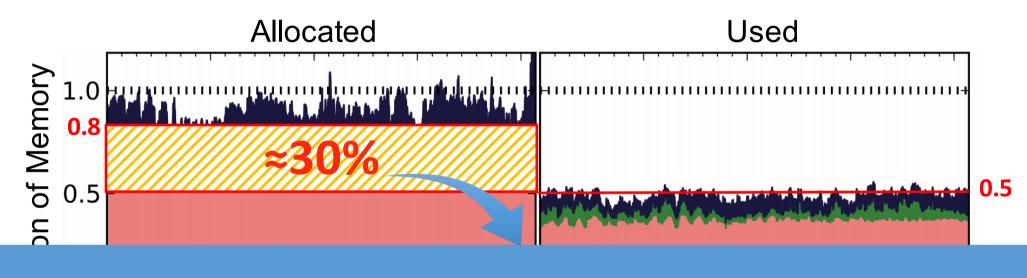
• Google Cluster Analysis



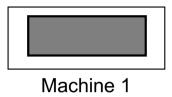
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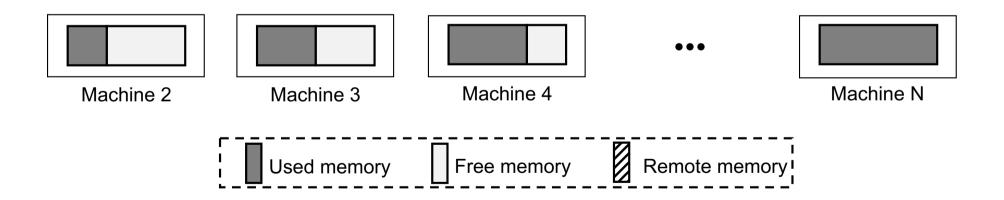


• Google Cluster Analysis₍₁₎

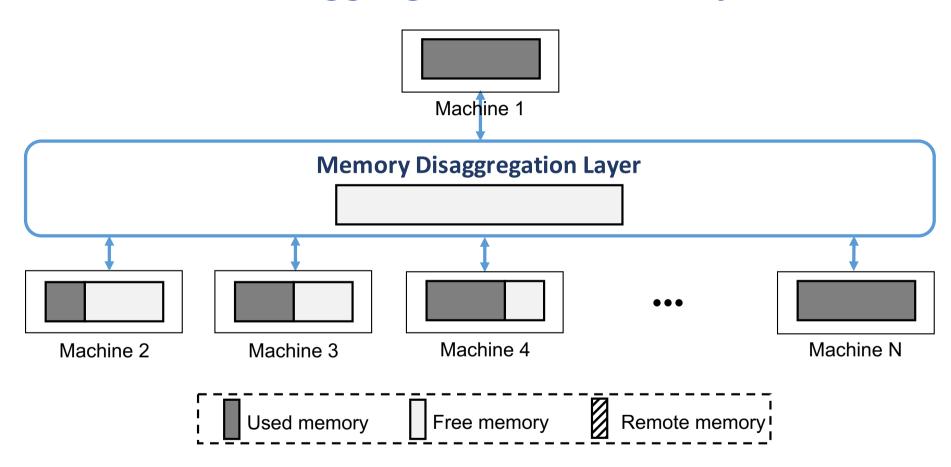


Can we utilize this memory?

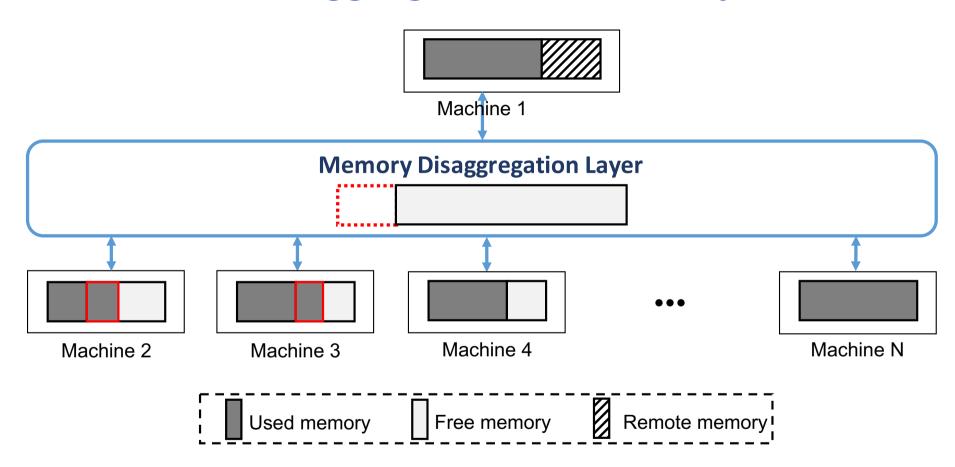




Disaggregate free memory



Disaggregate free memory



What are the challenges?

- Minimize deployment overhead
 - No hardware design
 - No application modification
- Tolerate failures
 - e.g. network disconnection, machine crash

Manage remote memory at scale

Recent work on memory disaggregation

	No HW design	No app modification	Fault- tolerance	Scalability
Memory Blade[ISCA'09]	×			
HPBD[CLUSTER'05] / NBDX[1]			×	X
RDMA key-value service (e.g. HERD[SIGCOMM'14], FaRM[NSDI'14])		×		
Intel Rack Scale Architecture (RSA)[2]	×			
Infiniswap				

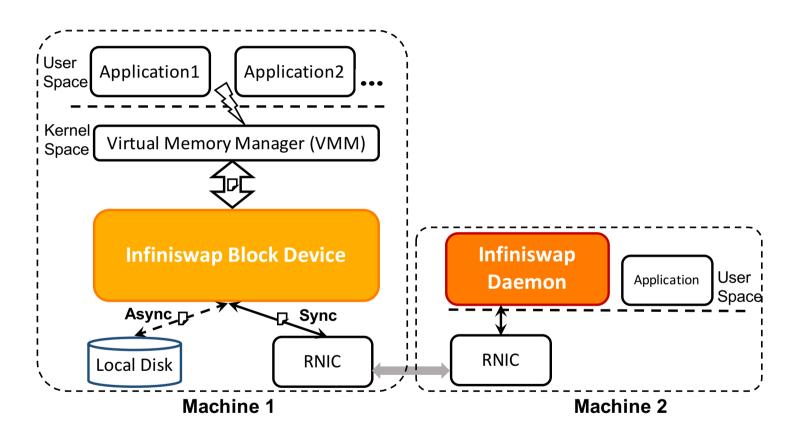
Agenda

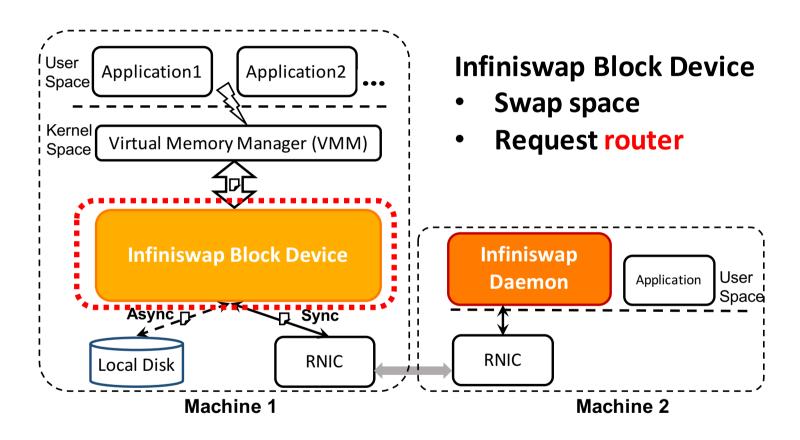
Motivation and related work

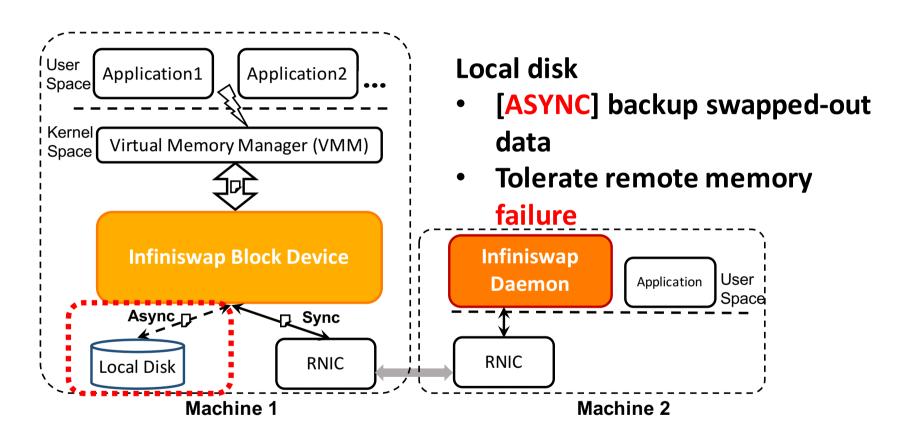
Design and system overview

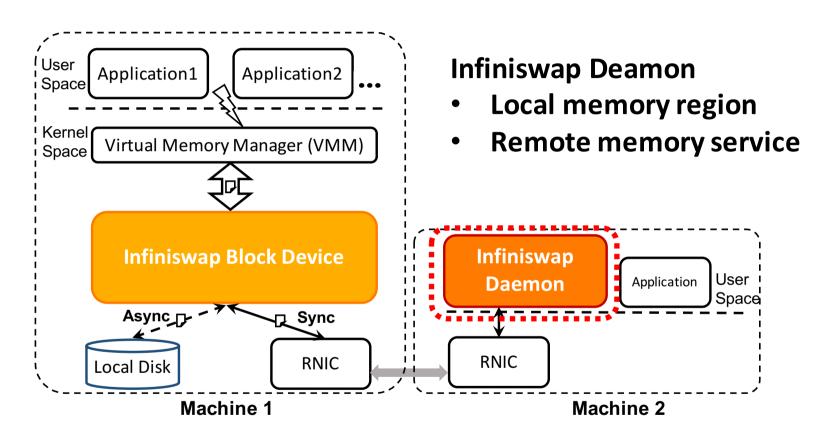
Implementation and evaluation

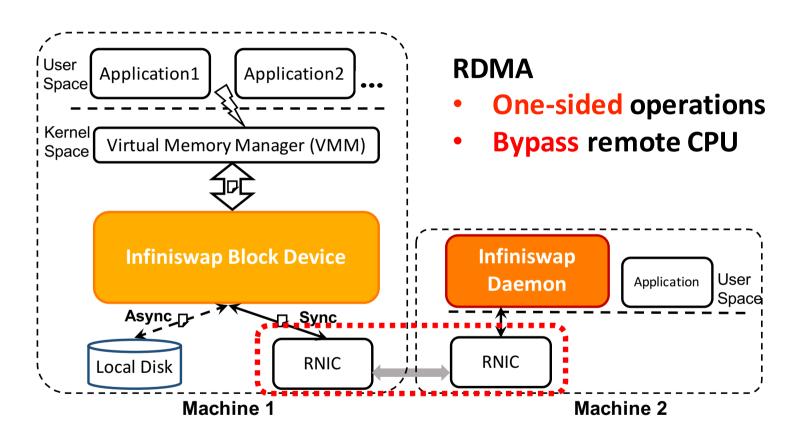
Future work and conclusion







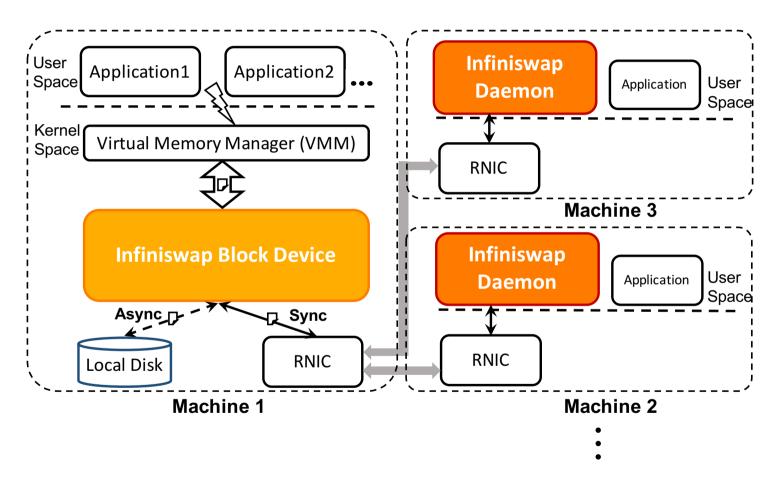




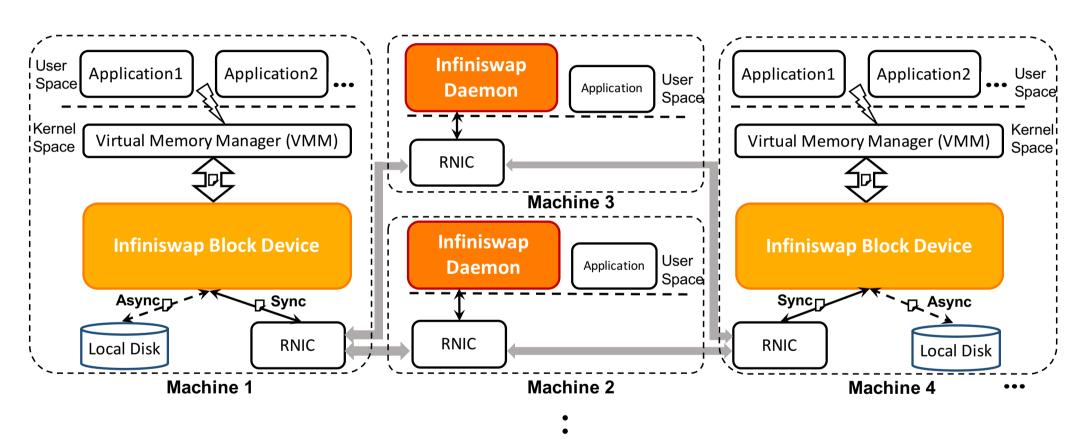
How to meet the design objectives?

Objectives	Ideas	
No hardware design	Remote paging	
No application modification		
Fault-tolerance	Local backup disk	

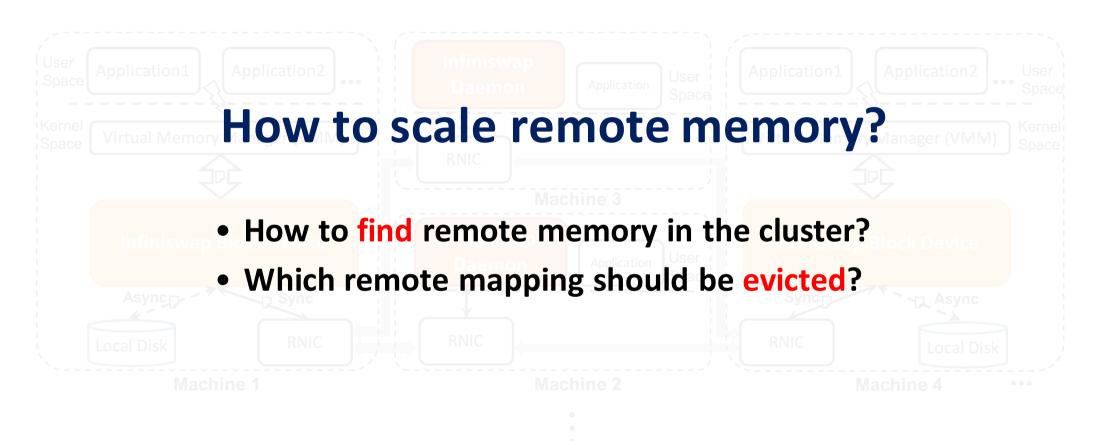
One-to-many



Many-to-many



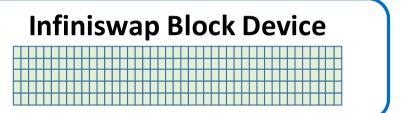
Many-to-many

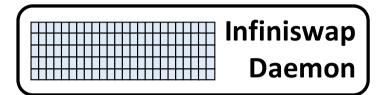


How to meet the design objectives?

Objectives	Ideas	
Scalability	Decentralized remote memory management	

Management unit: memory page?







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Management unit: memory page?





Local Page	Remote Page
p100	<s1, p1=""></s1,>

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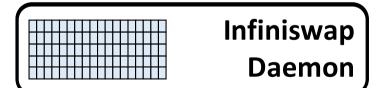
1GB = 256K entries

1GB = 256K RTTs

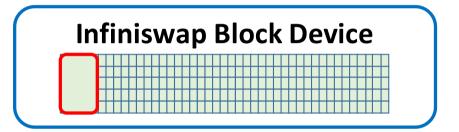


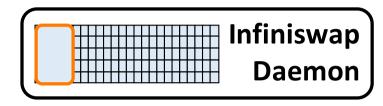


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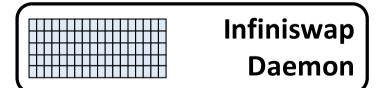
Management unit: memory slab!



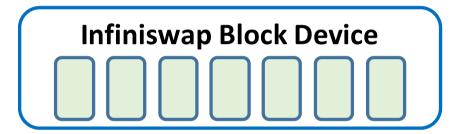


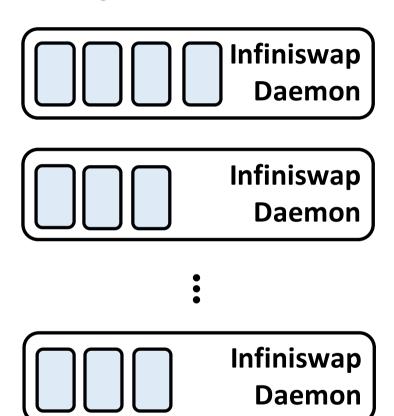


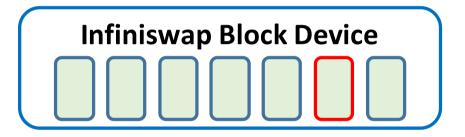
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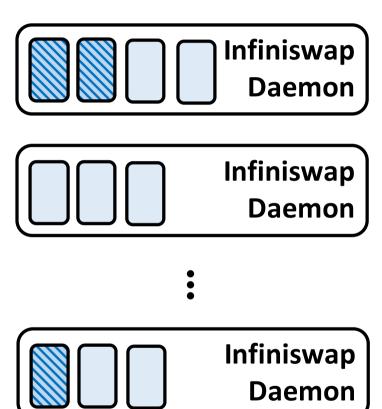


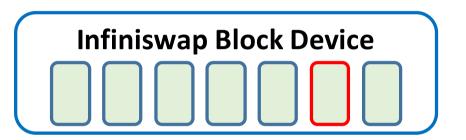
Management unit: memory slab!

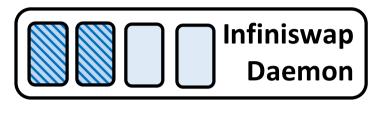








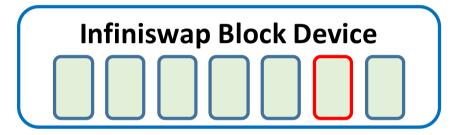




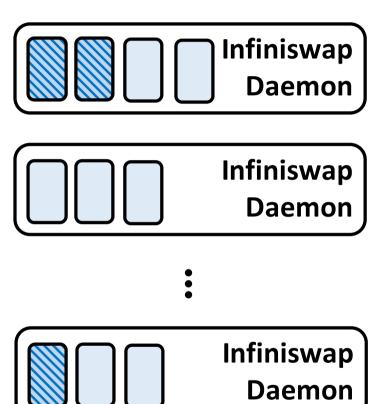


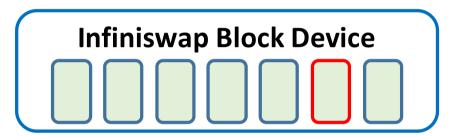
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Goal: balance memory utilization

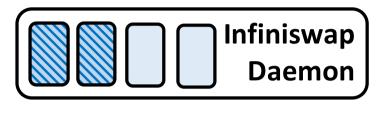


> Central controller



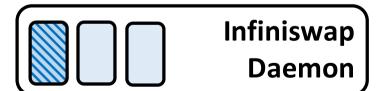


- > Central controller
- > Decentralized approach

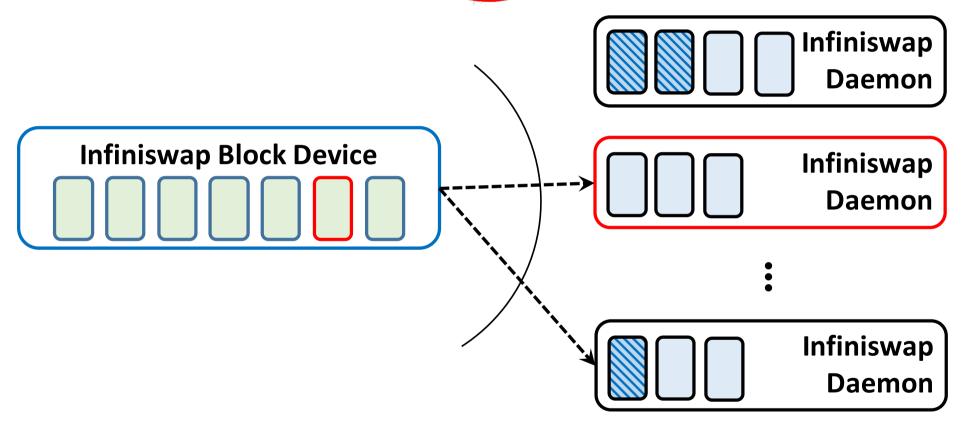




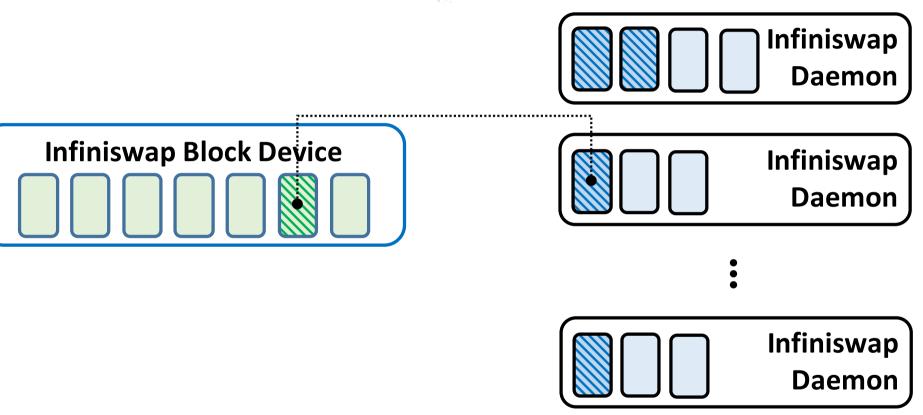
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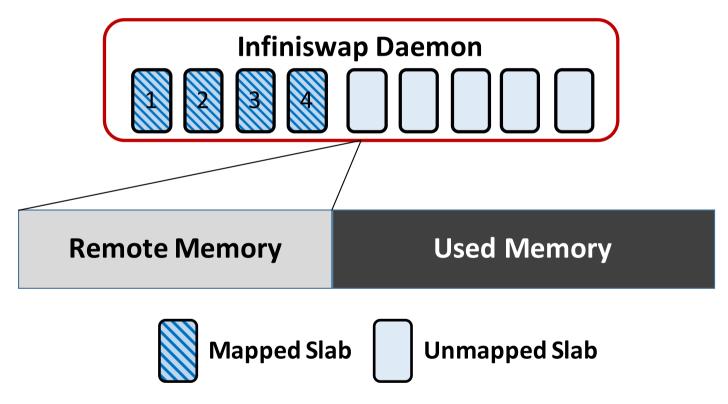






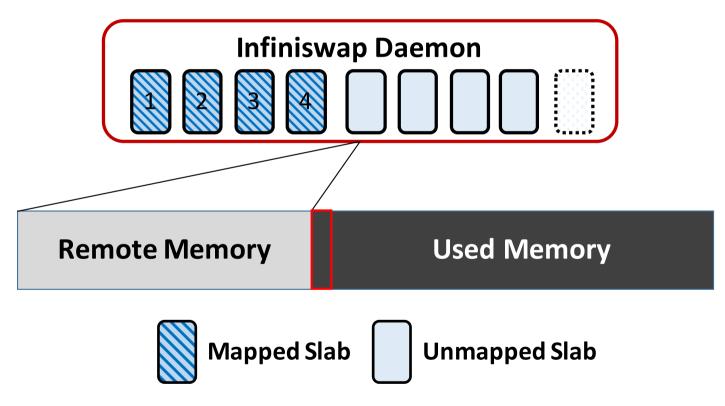


Slab eviction



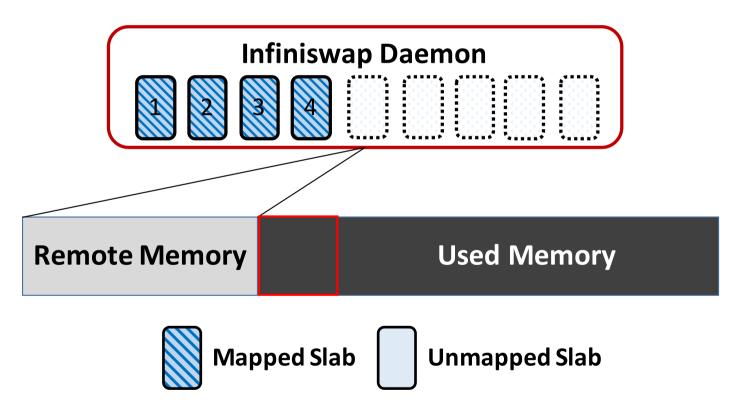
3/30/17 45

Slab eviction



3/30/17 46

Slab eviction



3/30/17 47

Which slab should be evicted?



Daemon: Does not know the swap activities

Which slab should be evicted?



Daemon: Too expensive to query all the slabs

Power of multiple choices_[1]



Select E least-active slabs from E+E' random slabs

Power of multiple choices_[1]



Select E least-active slabs from E+E' random slabs

Power of multiple choices_[1]



Select E least-active slabs from E+E' random slabs

Agenda

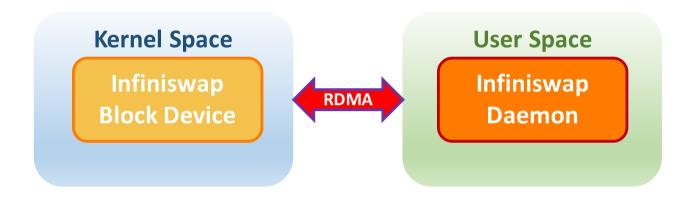
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Implementation



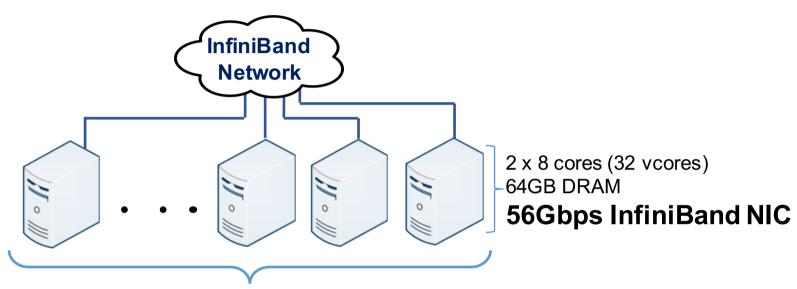
- Connection Management
 - One RDMA connection per active block device daemon pair
- Control Plane
 - SEND, RECV
- Data Plane
 - One-sided RDMA READ, WRITE

What are we expecting from Infiniswap?

- Application performance
- Cluster memory utilization
- Network usage
- Eviction overhead
- Fault-tolerance overhead
- Performance as a block device

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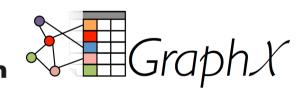
Evaluation



32-node cluster

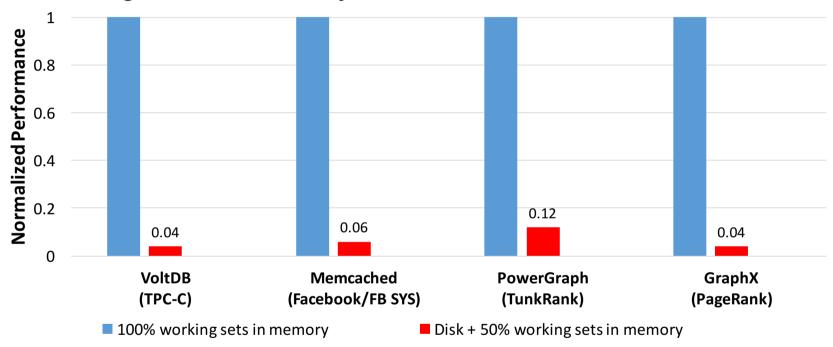






Application performance

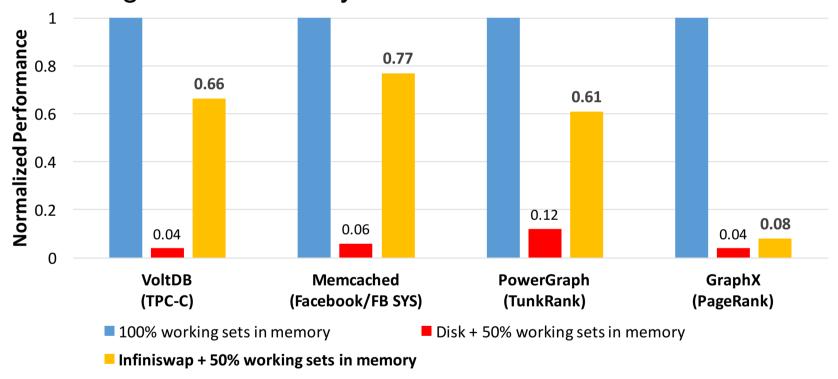
• 50% working sets in memory



Application performance is improved by 2-16x

Application performance

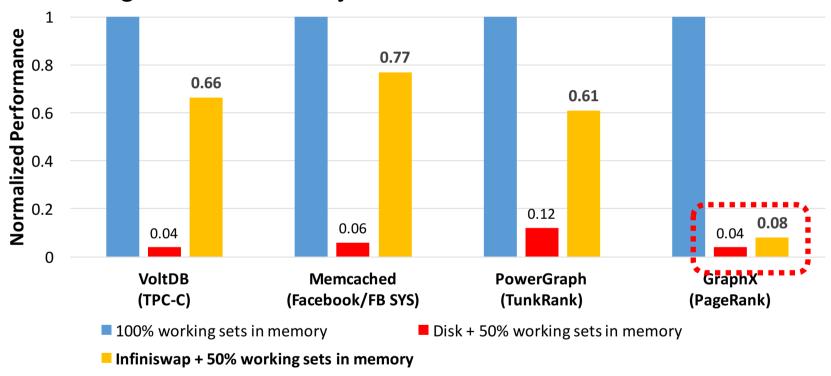
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Application performance

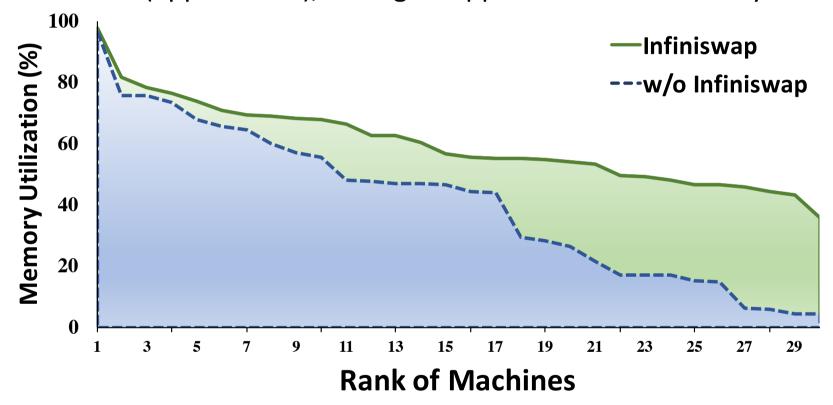
• 50% working sets in memory



Application performance is improved by 2-16x

Cluster memory utilization

• 90 containers (applications), mixing all applications and memory constraints.



Cluster memory utilization is improved from 40.8% to 60% (1.47x)

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Limitations and future work

- Trade-off in fault-tolerance
 - Local disk is the bottleneck
 - Multiple remote replicas
 - Fault-tolerance vs. space-efficiency
- Performance isolation among applications
 - W/o limitation on each application's usage
 - W/o mapping between remote memory and applications

Conclusion

- Infiniswap: remote paging over RDMA
 - Application performance
 - Cluster memory utilization
- Efficient, practical memory disaggregation
 - No hardware design
 - No application modification
 - Fault-tolerance
 - Scalability

Source code is coming soon!

https://github.com/Infiniswap/infiniswap.git

Thank You!