

1f.pdf

110101, 110101, 2000-2020 The People's Republic of China

1.1 (a) M_1, M_2 的起始状态是各自的 q_1

(b) M_1 的接受状态: $\{q_2\}$, M_2 的接受状态: $\{q_1, q_4\}$

(c) $M_1: q_1 \xrightarrow{a} q_2 \xrightarrow{a} q_3 \xrightarrow{b} q_1 \xrightarrow{b} q_1$

$M_2: q_1 \xrightarrow{a} q_1 \xrightarrow{a} q_1 \xrightarrow{b} q_2 \xrightarrow{b} q_4$

(d) M_1 不接受, M_2 接受.

1.2 M_1 的形式化描述:

$Q = \{q_1, q_2, q_3\}$

$\Sigma = \{a, b\}$

start state = q_1

$F = \{q_2\}$

$\delta(q_1, a) = q_2$

$\delta(q_1, b) = q_1$

$\delta(q_2, a) = q_3$

$\delta(q_2, b) = q_3$

$\delta(q_3, a) = q_2$

$\delta(q_3, b) = q_1$

M_2 的形式化描述:

$Q = \{q_1, q_2, q_3, q_4\}$

$\Sigma = \{a, b\}$

start state = q_1

$F = \{q_1, q_4\}$

$\delta(q_1, a) = q_1$ $\delta(q_1, b) = q_2$

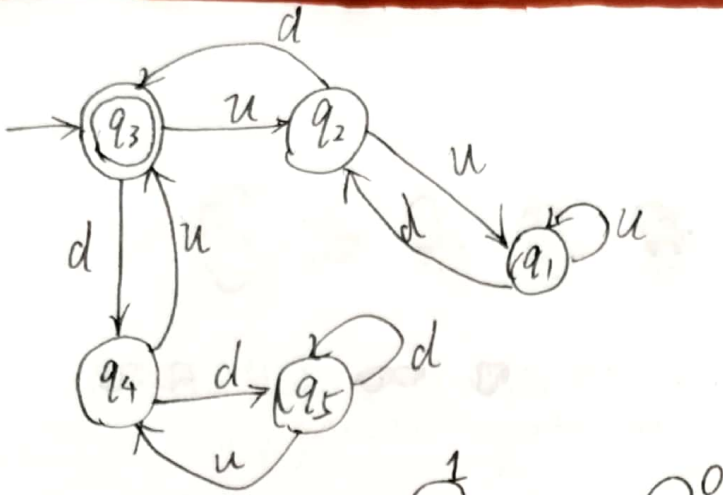
$\delta(q_2, a) = q_3$ $\delta(q_2, b) = q_4$

$\delta(q_3, a) = q_2$, $\delta(q_3, b) = q_1$

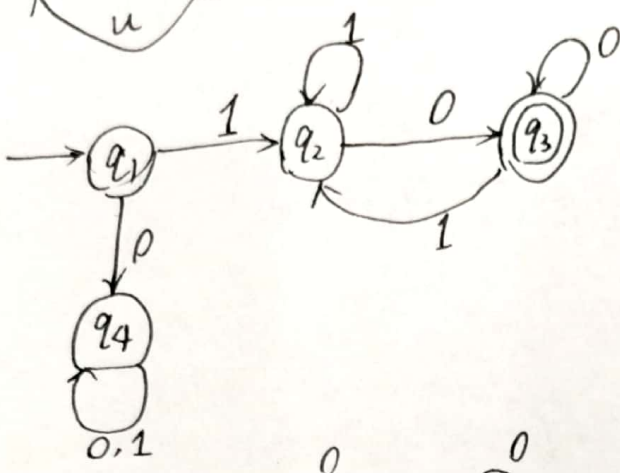
$\delta(q_4, a) = q_3$ - $\delta(q_4, b) = q_4$.



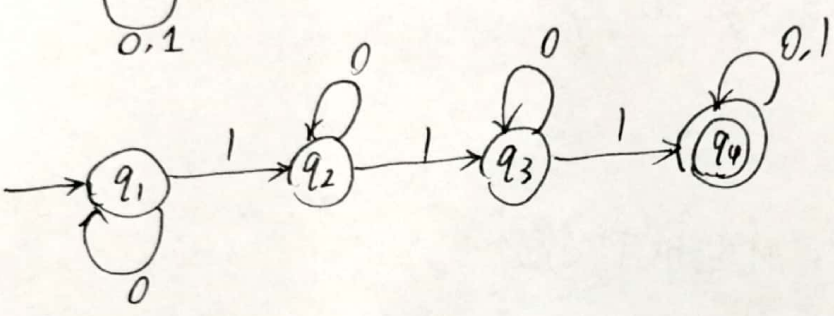
1.3



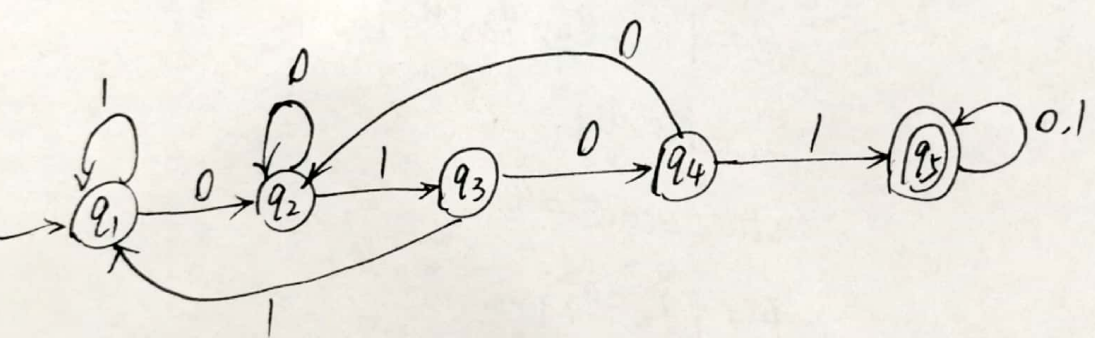
1.6 (a)



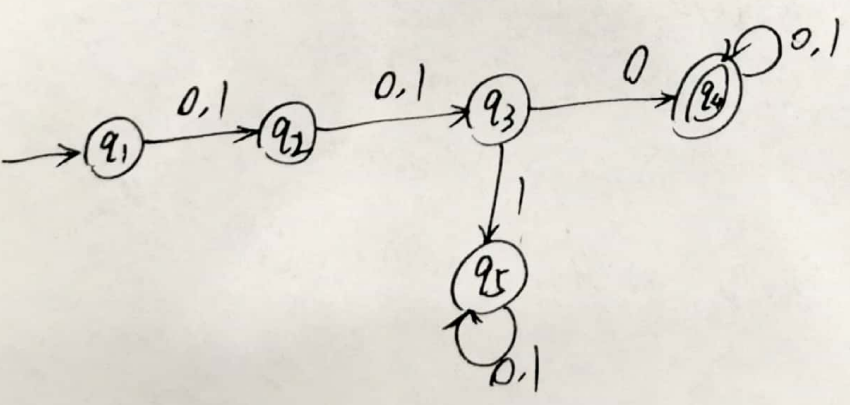
(b)



(c)



(d)



2f. pdf

a. Compute the complete truth table of the formula:

(1) $((p \rightarrow q) \rightarrow p) \rightarrow p$

| p | q | $p \rightarrow q$ | $(p \rightarrow q) \rightarrow p$ | $((p \rightarrow q) \rightarrow p) \rightarrow p$ |
|---|---|-------------------|-----------------------------------|---|
| T | T | T | T | T |
| T | F | F | T | T |
| F | T | T | F | T |
| F | F | T | F | T |

(2) $(p \wedge q) \rightarrow (p \vee q)$

| p | q | $p \wedge q$ | $p \vee q$ | $(p \wedge q) \rightarrow (p \vee q)$ |
|---|---|--------------|------------|---------------------------------------|
| T | T | T | T | T |
| T | F | F | T | T |
| F | T | F | T | T |
| F | F | F | F | T |

(3) $(p \rightarrow q) \vee (p \rightarrow \neg q)$

| $\neg q$ | p | q | $p \rightarrow q$ | $p \rightarrow \neg q$ | $(p \rightarrow q) \vee (p \rightarrow \neg q)$ |
|----------|---|---|-------------------|------------------------|---|
| F | T | T | T | F | T |
| T | T | F | F | T | T |
| F | F | T | T | T | T |
| T | F | F | T | T | T |

(4) $((p \vee q) \rightarrow r) \rightarrow ((p \rightarrow r) \vee (q \rightarrow r))$

| p | q | r | $p \vee q$ | $(p \vee q) \rightarrow r$ | $p \rightarrow r$ | $q \rightarrow r$ | $(p \rightarrow r) \vee (q \rightarrow r)$ | $((p \vee q) \rightarrow r) \rightarrow ((p \rightarrow r) \vee (q \rightarrow r))$ |
|---|---|---|------------|----------------------------|-------------------|-------------------|--|---|
| T | T | T | T | T | T | T | T | T |
| T | T | F | T | F | F | F | F | T |
| T | F | T | T | T | T | T | T | T |
| T | F | F | T | F | F | T | T | T |
| F | T | T | T | T | T | T | T | T |
| F | T | F | T | F | T | F | T | T |
| F | F | T | F | T | T | T | T | T |
| F | F | F | F | T | T | T | T | T |



1. (a) $\forall x (P(x) \rightarrow A(m, x))$

(b) $\exists x (P(x) \wedge A(x, m))$

(c) $A(m, m)$

(d) ~~$\exists x (S(x) \wedge \forall y$~~

Student x attended every lecture: $S(x) \wedge \forall y (L(y) \rightarrow B(x, y))$

~~$\neg \exists x (S(x) \wedge \forall y (L(y) \rightarrow B(x, y)))$~~

(e) Lecture y was attended by every student: $L(y) \wedge \forall x (P(x) \rightarrow B(x, y))$

~~$\neg \exists y (L(y) \wedge \forall x (P(x) \rightarrow B(x, y)))$~~

(f) Lecture y was not attended by any student: $\neg \exists x (P(x) \wedge B(x, y))$

$\forall y (\neg \exists x (P(x) \wedge B(x, y)))$

2. (a) 将 $P^M \stackrel{\text{def}}{=} \{(m, n) \mid m < n\}$ 代入 ϕ , ϕ 变成 $\forall x \exists y \exists z (x < y \wedge z < y \wedge (x < z \rightarrow z < x))$
可满足, 对于任意的 x , 取 $y = x+1$, $z = x$ 就可使 ϕ 得到满足
自然数

(b) 将 $P^M \stackrel{\text{def}}{=} \{(m, 2 \cdot m) \mid m \text{ natural number}\}$ 代入 ϕ .

$\phi = \forall x \exists y \exists z (y = 2x \wedge y = z \wedge (z = 2x \rightarrow x = 2z))$

可满足. 对于任意的自然数 x , 取 $y = 2x$, $z = x$ 就可使 ϕ 得到满足.

(c) 将 P^M 代入 ϕ , $\phi = \forall x \exists y \exists z (x < y+1 \wedge z < y+1 \wedge (x < z+1 \rightarrow z < x+1))$

可满足. 对于任意的自然数 x , 取 $y = x$, $z = x$ 就可使 ϕ 得到满足.

