ROW TRANSPOSITION CIPHER IMPLEMENTATION

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Abstract

The cipher appeared thousands of years ago but was heavily exploited from the 2nd World War. In our group project, we have learn one of the cryptography method in the transposition ciphers filed called Row Transposition cipher. Our group implements the Row Transposition cipher's encryption and decryption algorithm by using Java programming language.

I. Introduction

The cipher appeared thousands of years ago but was heavily exploited from the 2nd World War in which information exchanged between enemies, even encrypted, was intercepted and deciphered. There are two widely explored types of encryptions (transposition and substitution) that, when worked independently, can be solved quickly. Still, protocols have been created and improved that fundamentally use these two techniques repetitively. By definition, transposition cipher changes the plaintext order and rearranges to get ciphertext. In this group project, we used "Row Transposition Cipher," where you write your plaintext in rows of fixed length (key size), and we write by columns in key order. We can use the procedure to enhance the complexity of a more complex cipher-text.

II. Encryption Implementation

To implement the row transposition encryption, we utilize the key as a sequence to switch the columns in a two-dimension matrix to form a row transposition matrix(Figure 3).

Take the key 'NYITV' as an example (Figure 1), the algorithm uses the 26 English letters to find the number sequence '14023'. Then, the algorithm arranges the columns by the order of this number sequence. The encryption algorithm writes letters of message out in rows over a specified number of columns which equals the key length '5' (Figure 2). Then, reorder columns in the matrix (Figure 3).

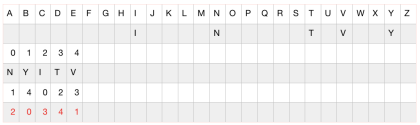

Figure 1: *Task1-Encryption Sequence Order*

The reorder sequence is '20341'(Figure1) for decryption in the columns of 2D Matrix, which transfers from the key order '14023'. For instance, the 0 column read 1st and it is in the RowMatrix column 2.

Regards for the current assignment, the empty space would be replaced by the capital letter 'X'. Append the rows to form the ciphertext.

N	Y	I	T	V		I	N	T	V	Y
0	1	2	3	4		2	0	3	4	1
c	r	y	p	t		y	c	p	t	r
o	l	o	g	y		o	o	g	y	l
i	s	t	h	e		t	i	h	e	s
p	r	a	c	t		a	p	c	t	r
i	c	e	a	n		e	i	a	n	c
d	s	t	u	d		t	d	u	d	s
y	o	f	t	e		f	y	t	e	o
c	h	n	i	q		n	c	i	q	h
u	e	s	f	o		s	u	f	o	e
r	s	e	c	u		e	r	c	u	s
r	e	c	o	m		c	r	o	m	e
m	u	n	i	c		n	m	i	c	u
a	t	i	o	n		i	a	o	n	t
i	n	t	h	e		t	i	h	e	n
p	r	e	s	e		e	p	s	e	r
n	c	e	o	f		e	n	o	f	c
t	h	i	r	d		i	t	r	d	h
p	a	r	t	i		r	p	t	i	a
e	s	c	a	l		c	e	a	l	s
l	e	d	a	d		d	l	a	d	e
v	e	r	s	a		r	v	s	a	e
r	i	e	s	x		e	r	s	x	i

Figure 2: *Task1-Row Transcription Matrix*

The following figure 3 displays the encrypted message.

yotaetfnsecniteeircdrecoipidycurrm
aipntpelvrpghcautifcoiohsortaasstye
tndeqoumcneefdildaXrlsrcsoheseut
nrchaseei

Figure 3: *Task1-Encrypted Message*

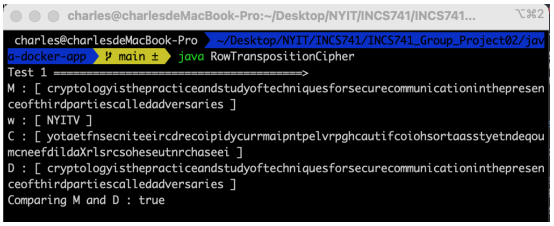

Figure 4: *Task1-Output*

Figure 5 shows the result of the encrypted plaintext. Further, we use the decrytion algorithm to double check the answer. The result match to the original text.

III. Decryption Implementation

The decryption algorithm put the letters in columns by the key order '14023'; however, as we fill the columns in the order '20341'. Take '0' as an example, The '0' is filled up 1st in the 2nd column. The algorithm first writes the encrypted message out in rows. Then, it reads off the message by recording columns.

I	N	T	V	Y		N	Y	I	T	V
2	0	3	4	1		0	1	2	3	4
e	t	r	o	h		t	h	e	r	o
r	w	a	n	t		w	t	r	a	n
o	s	s	i	p		s	p	o	s	i
o	t	n	c	i		t	i	o	n	c
h	i	e	r	p		i	p	h	e	r
a	i	f	a	s		i	s	a	f	a
l	i	y	s	r		i	r	l	y	s
p	i	l	e	m		i	m	p	l	e
s	e	y	t	a		e	a	s	y	t
m	o	p	l	i		o	i	m	p	l
e	e	n	t	m		e	m	e	n	t
p	c	h	e	i		c	i	p	h	e
t	r	i	s	i		r	i	t	i	s
r	a	a	n	t		a	t	r	a	n
o	s	s	i	p		s	p	o	s	i
o	t	n	c	i		t	i	o	n	c
h	i	e	r	p		i	p	h	e	r
a	t	t	f	h		t	h	a	t	f
l	o	o	w	l		o	l	l	o	w
s	s	i	m	a		s	a	s	i	m
e	p	r	u	l		p	l	e	r	u
f	l	o	r	e		l	e	f	o	r
x	m	i	n	i		m	i	x	i	n
p	g	t	h	u		g	u	p	t	h
h	e	a	r	c		e	c	h	a	r
t	a	e	r	c		a	c	t	e	r
n	s	t	h	i		s	i	n	t	h
l	e	a	i	p		e	p	l	a	i
e	n	x	t	t		n	t	e	x	t
f	t	o	r	o		t	o	f	o	r
h	m	e	c	t		m	t	h	e	c
h	i	e	r	p		i	p	h	e	r
x	t	t	x	e		t	e	x	t	x

Figure 5: *Task2-Decrypted Row Transposition Matrix*

therowtranspositioncipherisafairlysi
mpleeasytoimplementcipheritisatran
spositioncipherthatfollowsasimpleru
leformixingupthecharactersintheplai
ntexttoformtheciphertextx

Figure 6: *Task2-Decrypted Message*

The following is the pseudocode for the row transposition cipher algorithm :

The first part finds the decoding sequence from the key 'NYITV'. Then, write the encrypted message into a 2D RowTranspositionMatrix. Next, use the RowTranspositionMatrix to record the rearrangement of the RowTranspositionMatrix. In the last step, utilize the StringBuilder to build the encrypted message line by line through the RowMatrix.

Row Transposition Decryption Algorithm 1

input: 'w' : Key and 'C' Encrypted plain-text
output: Decrypted plain-text

```
1:  function RTCDECRYPTION(w, C)
2:      keylen ← w.length()
3:      keyArray ← key.toCharArray()
4:      messageArray ← C.toCharArray()
5:      keyPosition ← int[keylen]
6:
7:      Sort the keyArray        ▷ sort the keyArray and assign it to a string
8:      Strings ← String.valueOf(keyArray)
9:
10:     x ← 0
11:     for each char c in dArray do
12:         keyPosition[x] ← s.indexOf(c)
13:         Increament x by 1
14:     cols ← keylen
15:     rows ← 0
16:     if C's length mod cols equals 0 then    ▷ calculate rows
17:         rows ← C.length()/cols
18:     else                                    ▷ calculate columns
19:         rows ← C.length()/cols + 1
20:
21:     RowMatrix ← char[rows][cols]
22:     k ← 0
23:     for i to rows do
24:         for j to cols do
25:             if count k equals message C's length then
26:                 while k equals message's length and j
27:                 less than cols keep add 'X' to
28:                 RowMatrix[i][j]
29:                 break
30:                 assign RowMatrix[j][keyPosition[i]] from
31:                 messageArray[k]
32:                 Increament k by 1
33:
34:     StringBuilder str ← StringBuilder()
35:     for i to rows do
36:         for j to cols do
37:             if RowMatrix[i][j] unequal to 'X' then
38:                 str.append(RowMatrix[i][j])
39:     return str                              ▷ Decrypted message is str
```

The following figure 7 shows the result of the decrypted plaintext. Further, use the encryption algorithm to double-check the answer. The result matches the original encrypted text.

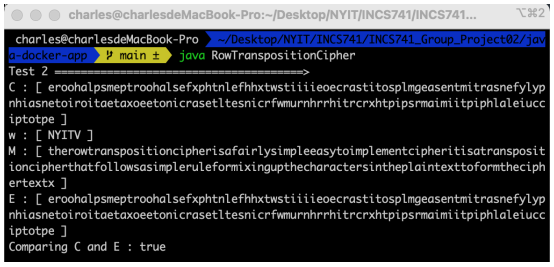

Figure 7: *Task2-Output*

IV. Conclusion

With the growing use of computers and the internet, and an increasing need to transmit information quickly and securely, encryption through existing protocols (AES, RSA, 3DES, etc.) became essential. In the project, we can see that using only one round of encryption (row transportation) and a minor key (5 letters), the information is already quite challenging to decipher, and with the use of the protocols mentioned above that repeatedly (using transportation, substitution, and other resources), it becomes harder to decipher the messages. We also demonstrate in the project that the information is decrypted, just doing the inverse of the encryption procedure that needs to be done by the person who will receive the message.

References

[1]N. Hamza, "Row transposition ciphers - ppt download," SlidePlayer. [Online]. Available: https://slideplayer.com/slide/13205094/. [Accessed: 28-Feb-2022].

[2]"Transposition cipher," Wikipedia, 23-Feb-2022. [Online]. Available: https://en.wikipedia.org/wiki/Transposition_cipher. [Accessed 28-Feb-2022].