

Problem 1 Consider the language of natural numbers, whose abstract syntax is given by the following CFG:

$$\begin{array}{lll}
 e \in \text{expr} & ::= & \text{value} \quad \text{value} \\
 & | & \text{expr} + \text{expr} \quad \text{addition} \\
 & | & \text{expr} \times \text{expr} \quad \text{multiplication} \\
 \\
 v \in \text{value} & ::= & 0 \quad \text{zero} \\
 & | & S(\text{value}) \quad \text{successor}
 \end{array}$$

And the operational semantics is given by the following rules:

$$\begin{array}{lll}
 \frac{}{v \Rightarrow v} \text{VAL} & \frac{e_1 \Rightarrow 0 \quad e_2 \Rightarrow v_2}{e_1 + e_2 \Rightarrow v_2} \text{ADD0} & \frac{e_1 \Rightarrow S(v_1) \quad v_1 + e_2 \Rightarrow v_2}{e_1 + e_2 \Rightarrow S(v_2)} \text{ADDS} \\
 \\
 \frac{e_1 \Rightarrow 0 \quad e_2 \Rightarrow v_2}{e_1 \times e_2 \Rightarrow v_2} \text{MUL0} & \frac{e_1 \Rightarrow S(v_1) \quad v_1 \times e_2 \Rightarrow v_2 \quad v_2 + e_2 \Rightarrow v_3}{e_1 \times e_2 \Rightarrow v_3} \text{MULS}
 \end{array}$$

Problem 1-A Show that there exists some value v such that $S(0) + S(0) \Rightarrow v$.

Problem 1-B Show that there exists some value v such that $S(0) \times S(0) \Rightarrow v$.

Problem 2 Consider the CoinPython language. The abstract syntax is given by the following CFG:

$p \in \text{prog}$	$::=$	pass	No-op (do nothing)
		raise	Exception
		print (m)	Printing a message
		$p_1; p_2$	Sequencing
		if (\star) $\{p_1\}$ else $\{p_2\}$	Non-deterministic branching
		while (\star) $\{p\}$	Non-deterministic loop

Let the set of messages be the set $\{\odot, \ominus\}$. Let $s_1 + s_2$ denotes the concatenation of strings s_1 and s_2 . The operational semantics of CoinPython is defined by the following set of rules:

$$\begin{array}{c}
\frac{}{\text{pass} \Rightarrow ""} \text{PASS} \quad (\text{No rule for raise}) \quad \frac{}{\text{print}(m) \Rightarrow m} \text{PRINT} \quad \frac{p_1 \Rightarrow s_1 \quad p_2 \Rightarrow s_2 \quad (s_1 + s_2 = s_3)}{p_1; p_2 \Rightarrow s_3} \text{SEQ} \\
\\
\frac{p_1 \Rightarrow s_1}{\text{if } (\star) \{p_1\} \text{ else } \{p_2\} \Rightarrow s_1} \text{IFTRUE} \quad \frac{p_2 \Rightarrow s_2}{\text{if } (\star) \{p_1\} \text{ else } \{p_2\} \Rightarrow s_1} \text{IFFALSE} \\
\\
\frac{}{\text{while } (\star) \{p\} \Rightarrow ""} \text{WHILEFALSE} \quad \frac{p \Rightarrow s_1 \quad \text{while } (\star) \{p\} \Rightarrow s_2 \quad (s_1 + s_2 = s_3)}{\text{while } (\star) \{p\} \Rightarrow s_3} \text{WHILETRUE}
\end{array}$$

Let P be the program (in concrete syntax):

```

while (*):
  if (*):
    print(⊙)
  else (*):
    print(⊙)
  if (*):
    print(⊙)
  else (*):
    print(⊙)

```

Problem 2-A Draw the AST for P .

Problem 2-B Show that $P \Rightarrow \odot\odot\odot\odot$.

Problem 2-C How would you describe the set of all strings that *may* be printed by P ?

Problem 3 Let Q be the Python program (in concrete syntax):

```
i = 0
you_won_one_million_usd = None
while found is None:
    if collatz(i):
        print(☺)
        print(☺)
    else:
        you_won_one_million_usd = i
        print(☺)
    i += 1
```

Assume that `collatz(i)` returns `True` if and only if the Collatz sequence starting at `i` converges to 1. Recall that the Collatz sequence is obtained by repeatedly applying the following two rules: a) if the number is even, divide it by two, and b) if the number is odd, triple it and add one. The Collatz conjecture – which remains an open problem – states that any Collatz sequence eventually converges to 1. The first person to solve the conjecture wins 120 million JPY (approx 0.8 million USD).

Problem 3-A Abstract Q into a CoinPython program \hat{Q} (by keeping only those features that are present in the CoinPython language). Then draw the AST for \hat{Q} .

Problem 3-B Does $\hat{Q} \Rightarrow \text{☺☺}$ hold? If so, draw the derivation tree. Otherwise, draw the partial tree and indicate the exact place(s) where the proof(s) gets stuck. If $\hat{Q} \Rightarrow \text{☺☺}$ does/doesn't hold, what does it tell you about the printing behavior of Q ?

Problem 3-C Does $\hat{Q} \Rightarrow \text{☺☺☺☺}$ hold? If so, draw the derivation tree. Otherwise, draw the partial tree and indicate the exact place(s) where the proof(s) gets stuck. If $\hat{Q} \Rightarrow \text{☺☺}$ does/doesn't hold, what does it tell you about the printing behavior of Q ?

Problem 3-D How would you describe the set of all strings that *may* be printed by Q ?

Problem 4 (bonus) Can you write a CoinPython program that prints all and only those strings such that every ☺ is matched by exactly one ☺ to its right?