

20th CSEC – Past Year Paper Solution 2015-2016 Sem 1
MH 1812 – Discrete Mathematics

1)

$$\begin{aligned}
 \text{a) } p \leftrightarrow q &\equiv (p \rightarrow q) \wedge (q \rightarrow p) && \text{(by definition)} \\
 &\equiv (\neg p \vee q) \wedge (\neg q \vee p) && \text{(by definition)} \\
 &\equiv [(\neg p \vee q) \wedge \neg q] \vee [(\neg p \vee q) \wedge p] && \text{(DeMorgan's)} \\
 &\equiv [(\neg p \wedge \neg q) \vee (q \wedge \neg q)] \vee [(\neg p \wedge p) \vee (p \wedge q)] && \text{(DeMorgan's)} \\
 &\equiv (\neg p \wedge \neg q) \vee (p \wedge q) \\
 &\equiv p \wedge q \vee \neg p \wedge \neg q && \text{(Operator precedence)}
 \end{aligned}$$

$$\text{b) } \neg(\exists x \in X, \forall y \in Y, P(x, y)) \equiv (\forall x \in X, \exists y \in Y, P(x, y))$$

X can only be 2 or 3:

① When $x = 2, y = 7$:

$$x \equiv y \pmod{5} = 2$$

That is, $P(2, 7)$ is true.

② When $x = 3, y = 8$:

$$x \equiv y \pmod{5} = 3$$

That is, $P(3, 8)$ is true.

So $\forall x \in X, \exists y \in Y, P(x, y)$ is true, which means $\neg(\exists x \in X, \forall y \in Y, P(x, y))$ is true.

c) A valid argument satisfies: If the premises are true, then the conclusion is true.

$$\begin{aligned}
 p \wedge q &&& \text{[premise 1]} \\
 \therefore p \text{ is true, and } q \text{ is true} &&& \text{[i]} \\
 p \rightarrow \neg r \text{ and } p \text{ is true} &&& \text{[premise 2 and i]} \\
 \therefore \neg r \text{ is true} &&& \text{[ii]} \\
 q \rightarrow \neg s \text{ and } q \text{ is true} &&& \text{[premise 3 and i]} \\
 \therefore \neg s \text{ is true} &&& \text{[iii]}
 \end{aligned}$$

So from (ii) and (iii), $\neg r \wedge \neg s$ is true.

According to the definition, this argument is valid.

2)

a)

i) Let $y \in f(S \cup T)$ arbitrary.

Then there exists $x \in (S \cup T)$ such that $f(x) = y$.

Since $x \in (S \cup T)$, $x \in S$ or $x \in T$, then $f(x) \in f(S)$ or $f(x) \in f(T)$.

That is, $f(x) \in f(S) \cup f(T)$

$$\text{Thus, } f(S \cup T) \subseteq f(S) \cup f(T) \quad [1]$$

Let $y \in f(S) \cup f(T)$ arbitrary.

Then there exists $x \in S$ or $x \in T$ such that $f(x) = y$.

Since $x \in S$ or $x \in T$, $x \in (S \cup T)$, then $f(x) \in f(S \cup T)$

$$\text{Thus, } f(S) \cup f(T) \subseteq f(S \cup T) \quad [2]$$

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From [1] and [2], $f(S \cup T) = f(S) \cup f(T)$.

ii) Disprove:

Suppose set $A = \{x1, x2, x3\}$, set $Y = \{y1, y2\}$, $f(x1) = f(x3) = y1, f(x2) = y2$.

If $S = \{x1, x2\}, T = \{x2, x3\}$, then $S \cap T = x2, f(x2) = y2$.

However, $f(S) = \{y1, y2\}, f(T) = \{y1, y2\}, f(S) \cap f(T) = \{y1, y2\}$, which is not equal to $y2$.

b) Yes, we can use membership table to solve this question. We should notice that premise is $A \oplus B = B \oplus C$, and we should get the conclusion that $A = C$.

| $A \oplus B$ | $B \oplus C$ | B | A | C |
|--------------|--------------|-----|-----|-----|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 0 |

And we can get the conclusion that $A = C$.

3) We can use matrix representation to solve this question.

Since the relation should be reflexive, which means every element of A is related to itself, the matrix's diagonal entries should all be true while the others can be true or false.

$$\begin{bmatrix} T & TorF & TorF & TorF \\ TorF & T & TorF & TorF \\ TorF & TorF & \ddots & TorF \\ TorF & TorF & TorF & T \end{bmatrix}$$

So we can construct 2^{n*n-n} distinct matrixes. That is, there are 2^{n*n-n} distinct reflexive relations on a set with n elements.

4)

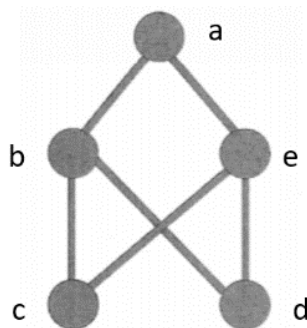


Figure 1: Graph

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- a) Yes, this graph is bipartite. Bipartite graph is a graph whose vertices can be partitioned into 2 (disjoint) subsets V and W such that each edge only connects a $v \in V$ and a $w \in W$.
Suppose $V = \{a, c, d\}$, $W = \{e, b\}$, then each edge only connects a node in V and a node in W .
Thus, this graph is bipartite.
- b) An Euler path (Eulerian trail) is a walk on the edges of a graph which uses each edge in the original graph exactly once. (The beginning and end of the walk may or not be the same vertex).
If we follow this order: $e \rightarrow c \rightarrow b \rightarrow d \rightarrow e \rightarrow a \rightarrow b$, we can walk on the edges of the graph which use each edge exactly once.

--End of Answers--

Solver: Deng Yue