

# Leet-Code Review

junyan.xu\*

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## 1 Important Algorithm

Before the solution, here are some summary of common graph algorithms

### 1.1 A Star Algorithm

The key idea is to always get minimum

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\*junyanxu5513@gmail.com

$$k(n) = g(n) + f(n) \quad (1)$$

as searching start

If  $f(n) \leq \tilde{f}(n)$ , then the  $f(n)$  is admissible, and the searching result is guaranteed to be smallest. Dijkstra is a special case for A\* algorithm where  $f(n) = 0$

Use a tree set as a priority queue for getting next.

## 2 Chapter 2

## 3 Chapter 3

## 4 Chapter 4

## 5 Chapter 5

### 5.1 Delete A Node In BST

Recursive is the best, remember to divide and conquer. Keep deleting what you copied.  
**time complexity  $O(\log(n))$ , mem  $O(\log(n))$**

```
class Solution {
public:
    TreeNode* deleteNode(TreeNode* root, int key) {
        if(!root) return NULL;
        if(root->val < key){
            root->right = deleteNode(root->right, key);
        }
        else if(root->val > key){
            root->left = deleteNode(root->left, key);
        }
        else{
            if(!root->left || !root->right){
                root = root->left ? root->left: root->right;
            }
            else{
                auto temp = root->right;
                while(temp->left)
                    temp = temp->left;
                root->val = temp->val;
                // key of the recursion
                root->right = deleteNode(root->right, temp->val);
            }
        }
    }
}
```

```

        return root;
    }
};

```

## 5.2 Minimum Number of Arrows to Burst Balloons

Greedy Algo, very similar to merge intervals. **time complexity**  $O(N \log(N))$ , **space is**  $O(1)$

```

class Solution {
public:
    int findMinArrowShots(vector<pair<int, int>>& points) {
        sort(points.begin(), points.end());
        int count = 0;
        int right = 0;
        for(int i=0; i<points.size(); i++){
            if(i==0){
                right = points[i].second;
                count++;
                continue;
            }
            if(points[i].first <= right)
                right = min(right, points[i].second);
            else{
                count++;
                right = points[i].second;
            }
        }
        return count;
    }
};

```

## 5.3 Minimum Moves to Equal Array Elements

Find minimum first, **time complexity**  $O(N)$ , **space is**  $O(1)$ . A set of good function in **algorithm** to be used

1. nth\_element(a.begin(), a.begin()+n, a.end())
2. \*min\_element(a.begin(), a.end())
3. \*max\_element(a.begin(), a.end())

```

class Solution {
public:
    int minMoves(vector<int>& nums) {
        int m = *min_element(nums.begin(), nums.end());
        long long res = 0;
        for(auto x: nums)
            res += abs((long long)x - m);
        return res;
    }
};

```

## 5.4 4 Sum Two

Use unordered\_map. **time complexity**  $O(N^2)$ , **space is**  $O(N^2)$

```

class Solution {
public:
    int fourSumCount(
        vector<int>& A,
        vector<int>& B,
        vector<int>& C,
        vector<int>& D
    ) {
        unordered_map<int, int> a, b;
        for(auto x: A) for(auto y: B) a[x+y]++;
        for(auto x: C) for(auto y: D) b[x+y]++;
        int res = 0;
        for(auto x: a){
            if(b.count(-x.first))
                res += x.second*b[-x.first];
        }
        return res;
    }
};

```

## 5.5 Assign Cookie

Sort and double pointer. **time complexity**  $O(N\log(N))$ , **space is**  $O(1)$

```

class Solution {
public:
    int findContentChildren(vector<int>& g, vector<int>& s) {
        sort(g.begin(), g.end());
        sort(s.begin(), s.end());
        int i=0, j=0, count=0;
        while(i < g.size() && j < s.size()){

```

```

        if(g[i] <= s[j]){
            i++;
            j++;
            count++;
        }
        else{
            j++;
        }
    }
    return count;
}
};

```

## 5.6 132 Pattern

Using reverse stack, keep track of the second largest. **time complexity  $O(N)$ , space is  $O(N)$**

```

class Solution {
public:
    bool find132pattern(vector<int>& nums) {
        stack<int> s;
        int s2 = INT_MIN;
        for(int i=nums.size()-1; i>=0; i--){
            if(nums[i] < s2)
                return true;
            while(!s.empty() && s.top() < nums[i]){
                s2 = s.top();
                s.pop();
            }
            s.push(nums[i]);
        }
        return false;
    }
};

```

## 5.7