

SDSC3001 Tutorial 3

Matrix Multiplication

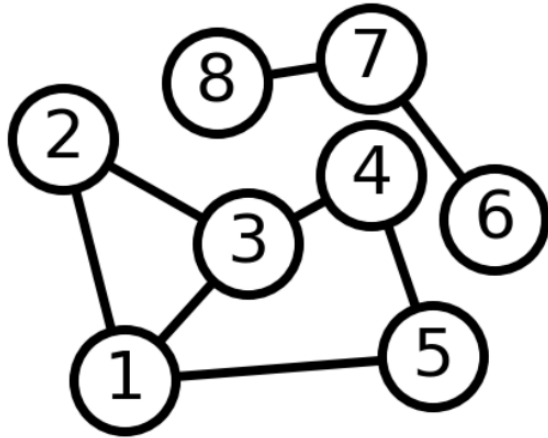
2023.09.25

Outline

- Review of Graph Data
- Matrix Multiplication and Representation of Sparse Matrix
- Power Iteration of PageRank

Review of Graph Data

Undirected Graph



8 × 8 matrix

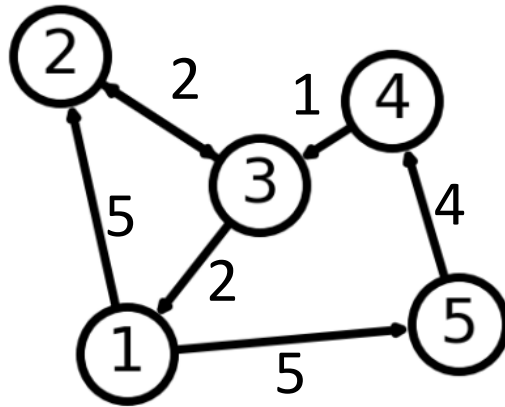
$$A_{8,8} = \begin{bmatrix} 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

$a_{ij} = 1$, if node i, j are connected

$a_{ij} = 0$, otherwise

Review of Graph Data

Directed Graph



No edge from node 2 to node 1

An edge with weight of 5 from node 1 to node 2

$$A_{5,5} = \begin{bmatrix} 0 & 5 & 0 & 0 & 5 \\ 0 & 0 & 2 & 0 & 0 \\ 2 & 2 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 4 & 0 \end{bmatrix}$$

Time complexity

Matrix-vector multiplication

- A matrix of shape (n, n) multiplied by a vector of shape $(n, 1)$

$$y = Ax = \begin{bmatrix} 0 & 2 & 1 \\ 4 & 0 & 3 \\ 1 & 4 & 0 \end{bmatrix} \cdot \begin{bmatrix} 2 \\ 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \times 2 + 2 \times 1 + 1 \times 0 \\ 4 \times 2 + 0 \times 1 + 3 \times 0 \\ 1 \times 2 + 4 \times 1 + 0 \times 0 \end{bmatrix} = \begin{bmatrix} 2 \\ 8 \\ 6 \end{bmatrix}$$

- $n \times n = O(n^2)$ multiplications

Time complexity

Matrix-matrix multiplication

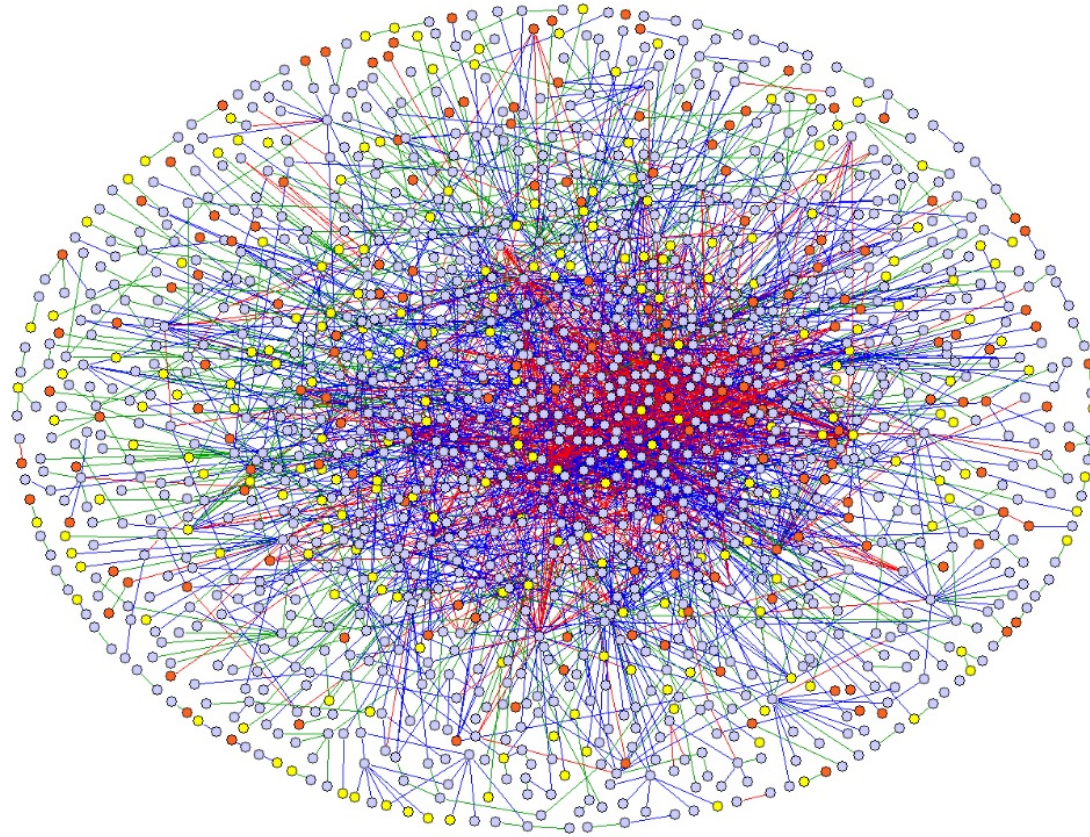
- Two matrices of shape (n, n)

$$y = AB = \begin{bmatrix} 0 & 2 & 1 \\ 4 & 0 & 3 \\ 1 & 4 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0 & 2 & 4 \\ 3 & 0 & 0 \\ 1 & 5 & 0 \end{bmatrix} = A \cdot \begin{bmatrix} 0 \\ 3 \\ 1 \end{bmatrix} + A \cdot \begin{bmatrix} 2 \\ 0 \\ 5 \end{bmatrix} + A \cdot \begin{bmatrix} 4 \\ 0 \\ 0 \end{bmatrix}$$

\downarrow
 $O(n^2)$

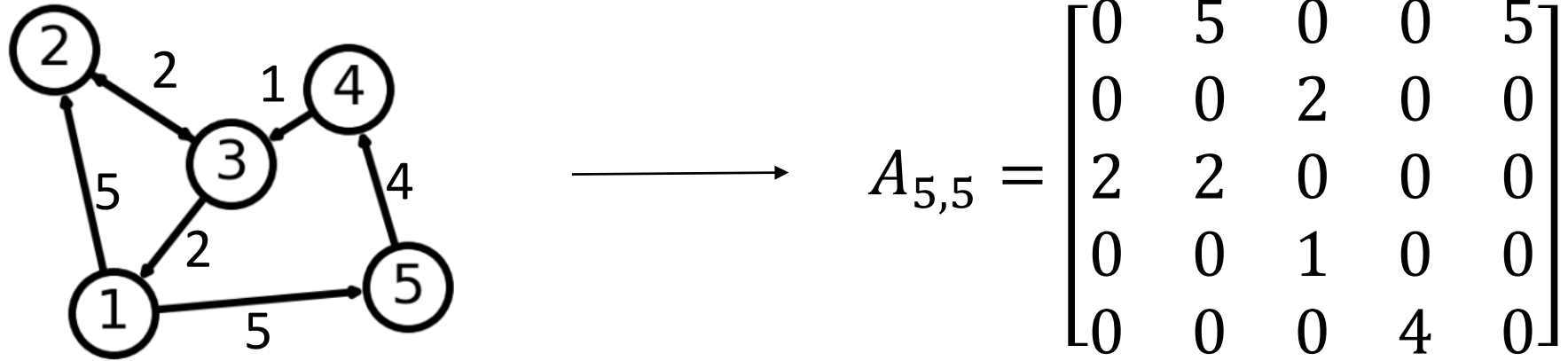
- $n^2 \times n = O(n^3)$ multiplications

Real-world Large Graph



n is large!

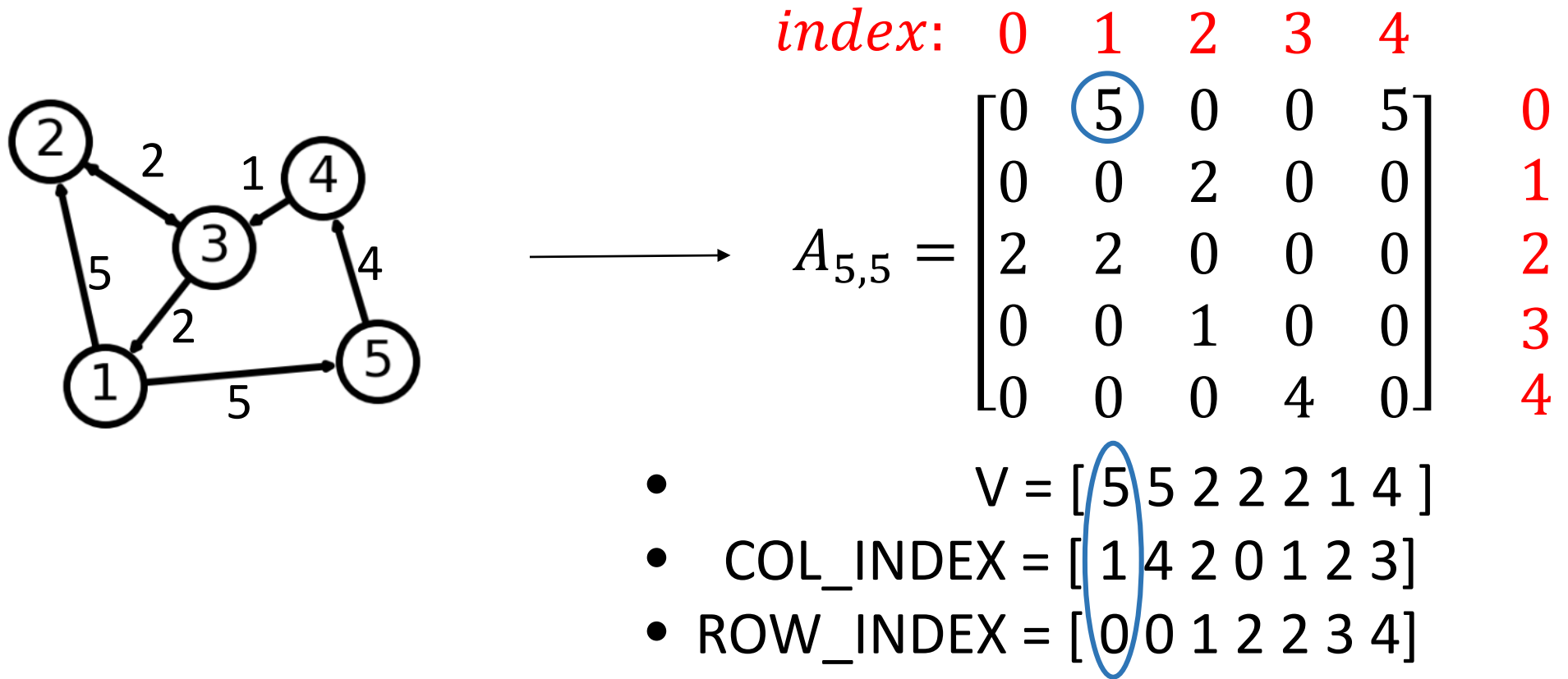
Sparse Matrix



- Space complexity is $O(n^2)$
- Only a few number of elements are not zero

Coordinate list (COO)

Let m be the number of non-zero values.



The space complexity is $O(3m)$.

Compressed sparse row (CSR)

Let m be the number of non-zero values and n be the number of rows.

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix}$$

The space complexity is $O(2m + n + 1)$.

- $V = [10\ 20\ 30\ 40\ 50\ 60\ 70\ 80]$
- $COL_INDEX = [0\ 1\ 1\ 3\ 2\ 3\ 4\ 5]$
- $ROW_INDEX = [0\ 2\ 4\ 7\ 8]$



Explanation 1

Element with (index, value) = (i, k) represents that there are k non-zero elements in the first i rows.

Compressed sparse row (CSR)

Let m be the number of non-zero values and n be the number of rows.

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix} \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix}$$

- $V = [10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80]$
- $COL_INDEX = [0 \ 1 \ 1 \ 3 \ 2 \ 3 \ 4 \ 5]$
- $ROW_INDEX = [0 \ 2 \ 4 \ 7 \ 8]$

Value $k=4$ and the index $i=2$

⇒ in the first 2 rows there are 4 non-zero elements (10, 20, 30, 40).

Compressed sparse row (CSR)

Let m be the number of non-zero values and n be the number of rows.

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix}$$

The space complexity is $O(2m + n + 1)$.

- $V = [10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80]$
- $COL_INDEX = [0 \ 1 \ 1 \ 3 \ 2 \ 3 \ 4 \ 5]$
- $ROW_INDEX = [0 \ 2 \ 4 \ 7 \ 8]$



Explanation 2

Another explanation: Element (except the last element) with (index, value) = (i, k) represents that the first non-zero element in row i is $V[k]$.

Compressed sparse row (CSR)

Let m be the number of non-zero values and n be the number of rows.

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix} \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix}$$

- $V = [10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80]$
- $COL_INDEX = [0 \ 1 \ 1 \ 3 \ 2 \ 3 \ 4 \ 5]$
- $ROW_INDEX = [0 \ 2 \ 4 \ 7 \ 8]$

Value $k=4$ and the *index* $i=2 \Rightarrow V[4]=50$.

The first non-zero element in row **2** is 50.

CSR Two Problems

- 1. Given a sparse matrix, how to define (V, COL_INDEX, ROW_INDEX)?
- 2. Given (V, COL_INDEX, ROW_INDEX), how to rebuild the sparse matrix/ find the location of a certain value?

1. Given a sparse matrix, how to define (V, COL_INDEX, ROW_INDEX)?

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix} \quad \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix}$$

Explanation 1

Element with (index, value) = (i, k) represents that there are k non-zero elements in the first i rows.

Explanation 2

- $V = [10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80]$
- $COL_INDEX = [0 \ 1 \ 1 \ 3 \ 2 \ 3 \ 4 \ 5]$
- $ROW_INDEX = [\quad]$

Element (except the last element) with (index, value) = (i, k) represents that the first non-zero element in row i is $V[k]$.

Q_1: how many elements are there in "ROW_INDEX"?

$i:0 \sim 4$, the length of "ROW_INDEX" is 5

$ROW_INDEX = [\quad]$

1. Given a sparse matrix, how to define (V, COL_INDEX, ROW_INDEX)?

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix} \quad \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix}$$

Explanation 1

Element with (index, value) = (i, k) represents that there are k non-zero elements in the first i rows.

k is obtained by counting the number

• $V = [10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80]$

$i=0$

In the first 0 rows (DO NOT EXIST), there are 0 non-zero elements

ALWAYS

ROW_INDEX = [0]

1. Given a sparse matrix, how to define (V, COL_INDEX, ROW_INDEX)?

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix}$$

Explanation 1

0 Element with (index, value) = (i, k) represents
1 that there are k non-zero elements in the first
2 i rows.

3 k is obtained by counting the number

• $V = [10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80]$

In the first 1 rows, there are 2 non-zero elements

$i=1$

ROW_INDEX = [0 2 7]

$i=3$

In the first 3 rows, there are 7 non-zero elements

1. Given a sparse matrix, how to define (V, COL_INDEX, ROW_INDEX)?

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix} \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix}$$

Explanation 2

Element (except the last element) with (index, value) = (i, k) represents that the first non-zero element in row i is $V[k]$.

k is obtained by finding index in V

• $V = [10, 20, 30, 40, 50, 60, 70, 80]$
Idx_V: 0 1 2 3 4 5 6 7

$i=0$

first non-zero element in row 0 is "10"



Its idx in V is "0"



$K=0$

ROW_INDEX = [0]

1. Given a sparse matrix, how to define (V, COL_INDEX, ROW_INDEX)?

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix} \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix}$$

Explanation 2

Element (except the last element) with (index, value) = (i, k) represents that the first non-zero element in row i is $V[k]$.

k is obtained by finding index in V

• $V = [10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80]$
 Idx_V: 0 1 2 3 4 5 6 7

$i=2$

first non-zero element in row 2 is "50"

↓
 $V(4)=50$; "50" idx in V is "4"

↓
 $K=4$

ROW_INDEX = [0 2 4]

Compressed sparse row (CSR)

Let m be the number of non-zero values and n be the number of rows.

index: 0 1 2 3 4 5

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix} \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix}$$

- $V = [\textcolor{red}{10} \textcolor{black}{20} \textcolor{orange}{30} \textcolor{black}{40} \textcolor{green}{50} \textcolor{black}{60} \textcolor{black}{70} \textcolor{blue}{80}]$
- $\text{COL_INDEX} = [0 \textcolor{black}{1} \textcolor{black}{1} \textcolor{black}{3} \textcolor{black}{2} \textcolor{black}{3} \textcolor{black}{4} \textcolor{black}{5}]$
- $\text{ROW_INDEX} = [\textcolor{red}{0} \textcolor{orange}{2} \textcolor{green}{4} \textcolor{blue}{7} \textcolor{black}{8}]$

Point to the first non-zero element in each row

2. Given $(V, \text{COL_INDEX}, \text{ROW_INDEX})$, how to rebuild the sparse matrix/ find the location of a certain value?

- $V = [10\ 20\ 30\ 40\ 50\ 60\ 70\ 80]$
- $\text{COL_INDEX} = [0\ 1\ 1\ 3\ 2\ 3\ 4\ 5]$
- $\text{ROW_INDEX} = [0\ 2\ 4\ 7\ 8]$

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix}$$

Explanation 1

Element with $(\text{index}, \text{value}) = (i, k)$ represents that there are k non-zero elements in the **first i rows**.

Explanation 2

Element (except the last element) with $(\text{index}, \text{value}) = (i, k)$ represents that the first non-zero element in **row i** is $V[k]$.

Q_1: Size of Matrix A?

$\text{Length}(\text{ROW_INDEX}) - 1$

*

$\max(\text{COL_INDEX}) + 1$

2. Given $(V, \text{COL_INDEX}, \text{ROW_INDEX})$, how to rebuild the sparse matrix/ find the location of a certain value?

Explanation 1

- $V = [10\ 20\ 30\ 40\ 50\ 60\ 70\ 80]$
- $\text{COL_INDEX} = [0\ 1\ 1\ 3\ 2\ 3\ 4\ 5]$
- $\text{ROW_INDEX} = [0\ 2\ 4\ 7\ 8]$

Element with $(\text{index}, \text{value}) = (i, k)$ represents that there are k non-zero elements in the first i rows.

$i=0$

$\text{ROW_INDEX}(0)=0$ always

$i=1$ $\text{ROW_INDEX}(1)=2$

In the first 1 rows, there are 2 non-zero elements (10&20)

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

2. Given $(V, \text{COL_INDEX}, \text{ROW_INDEX})$, how to rebuild the sparse matrix/ find the location of a certain value?

Explanation 1

- $V = [10\ 20\ 30\ 40\ 50\ 60\ 70\ 80]$
- $\text{COL_INDEX} = [0\ 1\ 1\ 3\ 2\ 3\ 4\ 5]$
- $\text{ROW_INDEX} = [0\ 2\ 4\ 7\ 8]$

Element with $(\text{index}, \text{value}) = (i, k)$ represents that there are k non-zero elements in the **first i rows**.

$i=2$ $\text{ROW_INDEX}(2)=4$

In the first **2** rows, there are **4** non-zero elements (10 20 30 40)

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

2. Given $(V, \text{COL_INDEX}, \text{ROW_INDEX})$, how to rebuild the sparse matrix/ find the location of a certain value?

Explanation 1

- $V = [10\ 20\ 30\ \textcolor{red}{40}\ 50\ 60\ 70\ 80]$
- $\text{COL_INDEX} = [0\ 1\ 1\ \textcolor{red}{3}\ 2\ 3\ 4\ 5]$
- $\text{ROW_INDEX} = [0\ \textcolor{red}{2}\ \textcolor{red}{4}\ 7\ 8]$

Element with $(\text{index}, \text{value}) = (i, k)$ represents that there are k non-zero elements in the **first i rows**.

How to get the coordinates of the element “40” (Column=?, Row=?) ?

1. Find idx in V : $k = V.\text{index}(\textcolor{brown}{40})=3$

2. Find the col: $\text{COL_INDEX}.\text{index}(k)=3$

3. Find the row: $\text{ROW_INDEX}[1]=2 \leq k < \text{ROW_INDEX}[2]=4 \Rightarrow \text{Row} = 1$

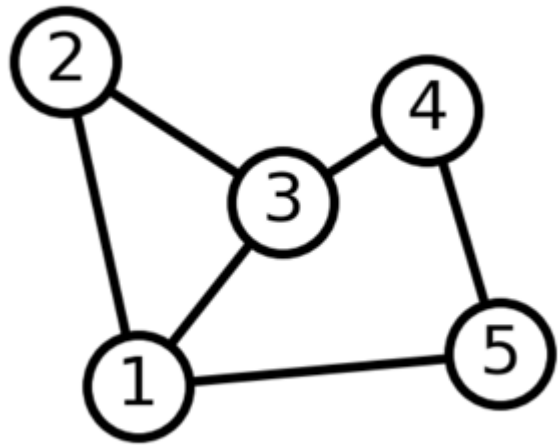
So the coordinate of element $\textcolor{brown}{40}$ is (1, 3)

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & \textcolor{red}{40} & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix}$$

Coding: PageRank

PageRank (PR) :

- an algorithm used by Google Search to rank web pages in their search engine results
- used to represent the likelihood that a person randomly clicking on links will arrive at any particular page.



Adjacency Matrix

$$A_{5,5} = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

Degree Matrix

$$D = \begin{bmatrix} 3 & 0 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 3 & 0 & 0 \\ 0 & 0 & 0 & 2 & 0 \\ 0 & 0 & 0 & 0 & 2 \end{bmatrix}$$

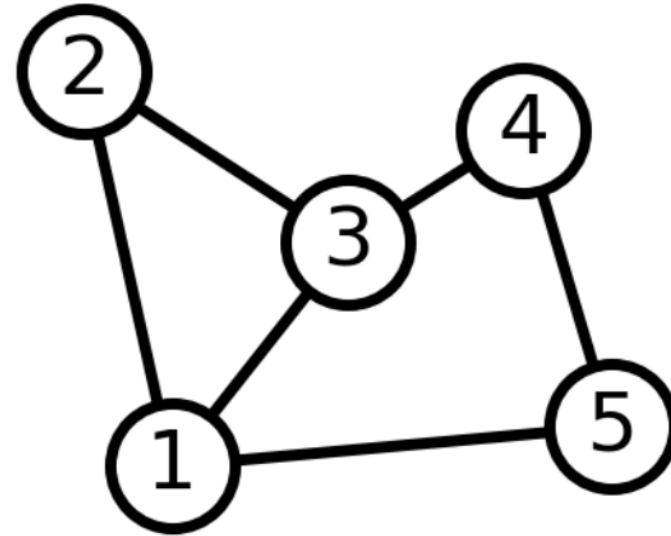
* $\deg(v_i)$ of a vertex counts the number of times an edge terminates at that vertex.

Coding: PageRank

- $\Pr(1 \rightarrow 2) = \Pr(1 \rightarrow 3) = \Pr(1 \rightarrow 5) = 1/3$
- $\Pr(2 \rightarrow 1) = \Pr(2 \rightarrow 3) = 1/2$

$$P = D^{-1}A$$
$$P = \begin{bmatrix} 0 & 1/2 & 1/3 & 0 & 1/2 \\ 1/3 & 0 & 1/3 & 0 & 0 \\ 1/3 & 1/2 & 0 & 1/2 & 0 \\ 0 & 0 & 1/3 & 0 & 1/2 \\ 1/3 & 0 & 0 & 1/2 & 0 \end{bmatrix}$$

transition probability



PageRank

► Power iteration method

starting with an arbitrary vector

► $\pi(0) = (\frac{1}{n}, \frac{1}{n}, \dots, \frac{1}{n})^\top$ operator

the operator is applied in succession

► $\pi(t) \leftarrow \left((1 - \alpha) \mathbf{P}^\top + \frac{\alpha}{n} \mathbf{E} \right) \pi(t - 1)$

► Repeat iterations until $\|\pi(t) - \pi(t - 1)\| \leq \epsilon$

$$(1 - \alpha) P^T \pi(t - 1) + \frac{\alpha}{n} E \pi(t - 1)$$

Sparse **Matrix-vector multiplication**

$\pi()$: probability distribution
P: transition probability matrix
E: matrix with all ones

Using CSR for Sparse Matrix-vector multiplication

$$A = \begin{bmatrix} 10 & 20 & 0 & 0 & 0 & 0 \\ 0 & 30 & 0 & 40 & 0 & 0 \\ 0 & 0 & 50 & 60 & 70 & 0 \\ 0 & 0 & 0 & 0 & 0 & 80 \end{bmatrix}$$
$$x = [1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 6]^T$$

```
1  # CSR format matrix A
2  v = [10, 20, 30, 40, 50, 60, 70, 80]
3  COL_INDEX = [0, 1, 1, 3, 2, 3, 4, 5]
4  ROW_INDEX = [0, 2, 4, 7, 8]
5
6  # vector
7  x = [1, 2, 3, 4, 5, 6]
8
9  # y = Ax
10 y = [0, 0, 0, 0]
11 m = ROW_INDEX[-1]  # the number of non-zero values
12
13
14 '''
15 In each row, figure out two things:
16 - 1. what are the non-zero elements?
17 = 2. their corresponding col idx.
18 '''
19
20 for i in range(len(ROW_INDEX) - 1):
21     row_index_start = ROW_INDEX[i]
22     row_index_end = ROW_INDEX[i+1]
23
24     for j in range(row_index_start, row_index_end):
25         # Length of the range list is the number of valid value in each row.
26         # v[j]: non-zero element
27         # COL_INDEX[j]: corresponding col idx
28         |
29         col = COL_INDEX[j]
30         y[i] += v[j] * x[col]
31
32 print(y)
```

[50, 220, 740, 480]

► Power iteration method

- $\pi(0) = (\frac{1}{n}, \frac{1}{n}, \dots, \frac{1}{n})^\top$
- $\pi(t) \leftarrow \left((1 - \alpha) \mathbf{P}^\top + \frac{\alpha}{n} \mathbf{E} \right) \pi(t - 1)$
- Repeat iterations until $\|\pi(t) - \pi(t - 1)\| \leq \epsilon$

$$(1 - \alpha)P^T \pi(t - 1) + \frac{\alpha}{n} E \pi(t - 1)$$

Sparse **Matrix-vector multiplication**

```

16 num_nodes = len(ROW_INDEX) - 1
17 pi = np.ones(num_nodes) / num_nodes  $\pi(0)$ 
18 print(pi)
19
20 alpha = 0.01
21 second_term = alpha * np.ones(num_nodes) / num_nodes
22 while True:
23     new_pi = np.zeros(num_nodes)
24     # (1-alpha)P * pi + alpha 1/n Matrix-vector multiplication
25     for i in range(len(ROW_INDEX) - 1):
26         row_index_start = ROW_INDEX[i]
27         row_index_end = ROW_INDEX[i + 1]
28
29         for j in range(row_index_start, row_index_end):
30             col = COL_INDEX[j]
31             value = v[j]
32             new_pi[i] += value * pi[col]  $P^T \pi(t - 1)$ 
33     new_pi = new_pi * (1-alpha) + second_term * pi
34
35     # check convergence by L1-distance(a, b) = sum_i |a_i - b_i|
36     if np.sum(np.abs(new_pi - pi)) < 1e-4:
37         break
38     else:
39         pi = new_pi termination condition
40
41 print(pi)
42
43 f = open("pi.pickle", 'wb')
44 pickle.dump(pi, f)
45

```

```

[0.5 0.5 0.03225806 ... 0.14285714 0.14285714 0.14285714]
[ 6547 15408 17793 ... 216687 256538 260898]
[ 0 2 33 ... 2312483 2312490 2312497]
[3.54731947e-06 3.54731947e-06 3.54731947e-06 ... 3.54731947e-06

```

Thank you!