

## chapter 4 : computing eigenvalues

### 4.1 introduction

problem : Given  $A$ , find  $\lambda$  and  $x \neq 0$  such that  $Ax = \lambda x$ .

$\lambda$  : e-value (e.g. frequency, growth rate, energy level)

$x$  : e-vector (e.g. normal mode, principal component, bound state)

thm : Assume  $A$  is real and symmetric. Then the e-values  $\lambda_i$  are real and the e-vectors  $q_i$  form an orthonormal basis, i.e.  $q_i^T q_j = 0$  for  $i \neq j$ ,  $\|q_i\|_2 = 1$ , and any  $x$  can be written as a linear combination of the  $q_i$ . (pf : omit)

ex :  $A = \begin{pmatrix} 2 & -1 \\ -1 & 2 \end{pmatrix}$

$$f_A(\lambda) = \det \begin{pmatrix} 2-\lambda & -1 \\ -1 & 2-\lambda \end{pmatrix} = (2-\lambda)^2 - 1 = \lambda^2 - 4\lambda + 3 = (\lambda - 3)(\lambda - 1)$$

$$\lambda_1 = 3 : Ax = 3x \Rightarrow \begin{pmatrix} 2 & -1 \\ -1 & 2 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = 3 \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}$$

choose  $x_1 = 1$ , then  $2x_1 - x_2 = 3x_1 \Rightarrow x_2 = -1, -x_1 + 2x_2 = 3x_2$  ok

$$q_1 = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ -1 \end{pmatrix} \Rightarrow \|q_1\|_2 = 1$$

$$\lambda_2 = 1 : Ax = x \Rightarrow \begin{pmatrix} 2 & -1 \\ -1 & 2 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}$$

choose  $x_1 = 1$ , then  $2x_1 - x_2 = x_1 \Rightarrow x_2 = 1, -x_1 + 2x_2 = x_2$  ok

$$q_2 = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \Rightarrow \|q_2\|_2 = 1, q_1^T q_2 = 0 \quad \text{ok}$$

### obvious method for computing e-values

step 1. form  $f_A(\lambda) = \det(A - \lambda I)$

step 2. solve  $f_A(\lambda) = 0$  by the methods of chapter 2

ex :  $A = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ ,  $\tilde{A} = \begin{pmatrix} 1+\epsilon & 0 \\ 0 & 1-\epsilon \end{pmatrix}$  : perturbed matrix

$$f_A(\lambda) = (1-\lambda)^2 = \lambda^2 - 2\lambda + 1 = 0 \Rightarrow \lambda = 1$$

$$f_{\tilde{A}}(\lambda) = (1+\epsilon-\lambda)(1-\epsilon-\lambda) = \lambda^2 - 2\lambda + 1 - \epsilon^2 = 0 \Rightarrow \lambda = 1 \pm \epsilon$$

1. A change in the elements of  $A$  of size  $\epsilon$  leads to a change in the e-values of size  $\epsilon$ .

2. A change in the coefficients of  $f_A(\lambda)$  of size  $\epsilon^2$  leads to a change in the roots of size  $\epsilon$ .

3. Hence the roots of  $f_A(\lambda)$  depend sensitively on the coefficients, and this implies that the obvious method for computing e-values is unstable.

ex (Wilkinson)

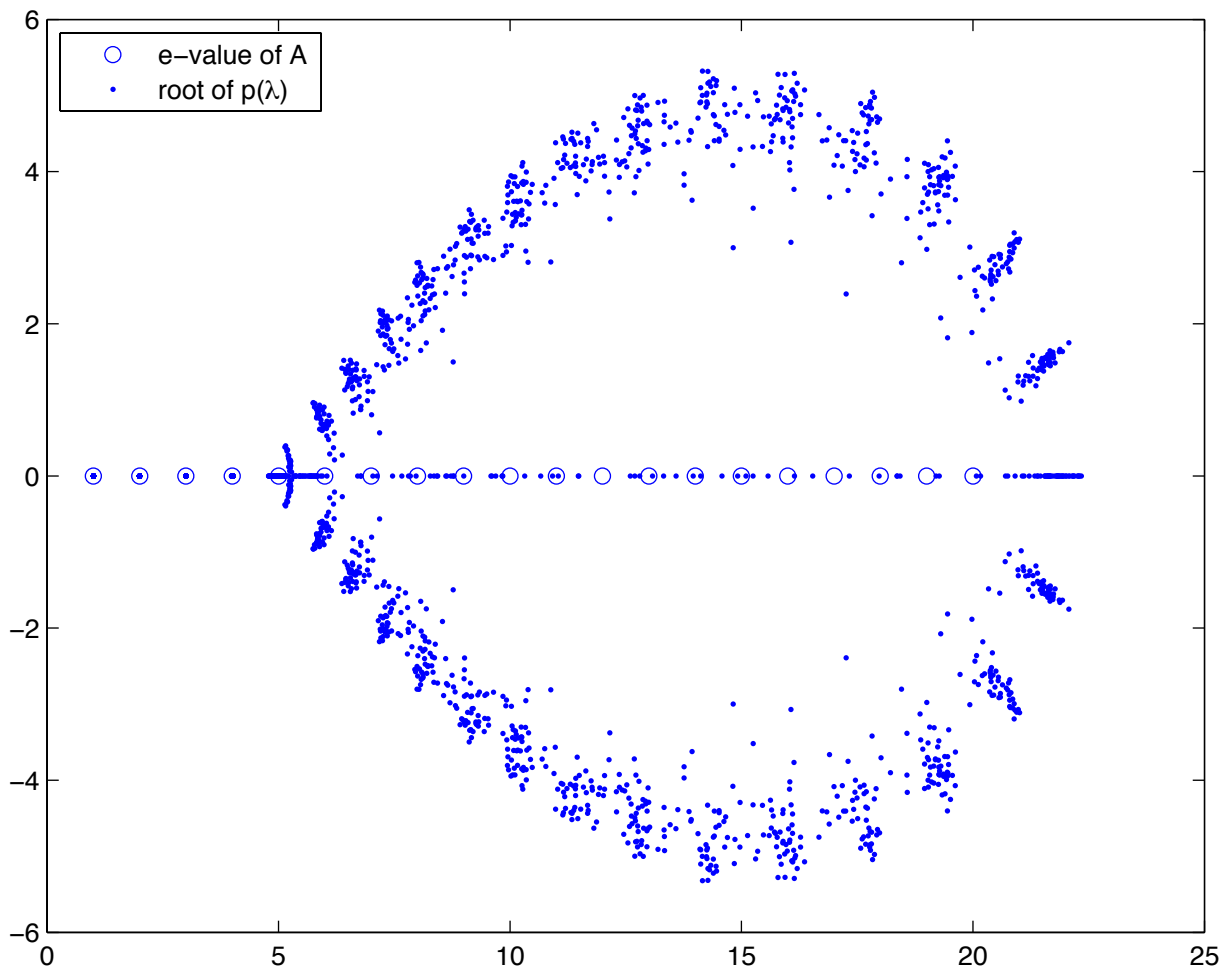
$$A = \text{diag}(1, 2, \dots, 20)$$

$$f_A(\lambda) = (1 - \lambda)(2 - \lambda) \cdots (20 - \lambda) = \sum_{k=0}^{20} a_k \lambda^k$$

$$\text{set } \tilde{a}_k = a_k(1 + 10^{-10}\epsilon_k), \epsilon_k \in (0, 1) : \text{random}, p(\lambda) = \sum_{k=0}^{20} \tilde{a}_k \lambda^k, \text{ roots} = ?$$

Matlab

```
plot(zeros(1,20),'o'); hold on;
for i=1:100
    r = roots(poly(1:20).*(ones(1,21)+1e-10*randn(1,21))));
    plot(r, '.'); axis([0,25,-6,6]);
end
```



1. This example shows that the roots of the characteristic polynomial are very sensitive to perturbations in the coefficients, and hence solving  $f_A(\lambda) = 0$  numerically is not a practical method for computing e-values.
2. question : What method does Matlab use to compute the roots of  $p(\lambda)$ ?

def : Given any  $x \neq 0$ , define  $R_A(x) = \frac{x^T A x}{x^T x}$  : Rayleigh quotient.

note

$$1. \text{ For } x = q_i, R_A(q_i) = \frac{q_i^T A q_i}{q_i^T q_i} = \frac{q_i^T \lambda_i q_i}{q_i^T q_i} = \lambda_i.$$

2. For  $x \approx q_i$ ,  $R_A(x)$  is an approximation to  $\lambda_i$  and we can derive an error estimate by Taylor expansion. First recall some notation.

$$f(x_1, x_2) = f(a_1, a_2) + \frac{\partial f}{\partial x_1}(a_1, a_2)(x_1 - a_1) + \frac{\partial f}{\partial x_2}(a_1, a_2)(x_2 - a_2) + \dots$$

$$f(x) = f(a) + \nabla f(a) \cdot (x - a) + O(\|x - a\|^2) \quad , \quad \nabla f = \left( \frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2} \right)$$

$$R_A(x) = R_A(q_i) + \nabla R_A(q_i) \cdot (x - q_i) + O(\|x - q_i\|^2)$$

$$\nabla R_A(x) = \nabla \left( \frac{x^T A x}{x^T x} \right) = \frac{x^T x \cdot \nabla(x^T A x) - x^T A x \cdot \nabla(x^T x)}{(x^T x)^2}$$

$$\nabla(x^T x) = \nabla(x_1^2 + x_2^2) = (2x_1, 2x_2) = 2x^T$$

$$x^T A x = (x_1, x_2) \begin{pmatrix} a_{11} & a_{12} \\ a_{12} & a_{22} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = a_{11}x_1^2 + 2a_{12}x_1x_2 + a_{22}x_2^2$$

$$\nabla(x^T A x) = (2a_{11}x_1 + 2a_{12}x_2, 2a_{12}x_1 + 2a_{22}x_2) = 2(Ax)^T$$

$$\nabla R_A(x) = \frac{x^T x \cdot 2(Ax)^T - x^T A x \cdot 2x^T}{(x^T x)^2} = \frac{2}{x^T x}((Ax)^T - R_A(x)x^T)$$

$$\nabla R_A(q_i) = \frac{2}{q_i^T q_i}((Aq_i)^T - R_A(q_i)q_i^T) = 2((\lambda_i q_i)^T - \lambda_i q_i^T) = 0$$

$$\Rightarrow R_A(x) = \lambda_i + O(\|x - q_i\|^2) : \text{quadratic approximation}$$

## 4.2 power method

idea :  $v, Av, A^2v, \dots$

algorithm

1.  $v^{(0)}$  : given ,  $\|v^{(0)}\|_2 = 1$  ,  $\lambda^{(0)} = (v^{(0)})^T A v^{(0)}$
2. for  $k = 1, 2, \dots$
3.  $w = Av^{(k-1)}$       % this can be done efficiently if  $A$  is sparse
4.  $v^{(k)} = w / \|w\|_2$       % this is done to avoid overflow/underflow
5.  $\lambda^{(k)} = (v^{(k)})^T A v^{(k)}$       %  $\lambda^{(k)} \rightarrow \lambda_1$  : largest e-value of  $A$  in absolute value

$$\underline{\text{ex}} : A = \begin{pmatrix} 2 & 1 & 1 \\ 1 & 3 & 1 \\ 1 & 1 & 4 \end{pmatrix} \Rightarrow \begin{cases} \lambda_1 = 5.214320 \\ \lambda_2 = 2.460811 \\ \lambda_3 = 1.324869 \end{cases}$$

$$v^{(0)} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \cdot \frac{1}{\sqrt{3}} \Rightarrow \lambda^{(0)} = (v^{(0)})^T A v^{(0)} = 5$$

power method

$k$	$\lambda^{(k)}$	$ \lambda^{(k)} - \lambda_1 $	$\frac{ \lambda^{(k)} - \lambda_1 }{ \lambda^{(k-1)} - \lambda_1 }$
0	5.000000	0.214320	-
1	5.181818	0.032502	0.151650
2	5.208193	0.006127	0.188513
$\downarrow$	$\downarrow$	$\downarrow$	$\downarrow$
$\infty$	$\lambda_1$	0	$(\lambda_2/\lambda_1)^2$

thm : Assume that  $|\lambda_1| > |\lambda_2| > \dots > |\lambda_n|$  and  $q_1^T v^{(0)} \neq 0$ .

Then  $\|v^{(k)} - (\pm q_1)\| = O\left(\left|\frac{\lambda_2}{\lambda_1}\right|^k\right)$ ,  $|\lambda^{(k)} - \lambda_1| = O\left(\left|\frac{\lambda_2}{\lambda_1}\right|^{2k}\right)$ .

pf :  $v^{(0)} = \alpha_1 q_1 + \alpha_2 q_2 + \dots + \alpha_n q_n$ , where  $\alpha_i = q_i^T v^{(0)}$

$$\begin{aligned} v^{(k)} &= \beta_k A^k v^{(0)} = \beta_k (\alpha_1 A^k q_1 + \alpha_2 A^k q_2 + \dots + \alpha_n A^k q_n) \\ &= \beta_k (\alpha_1 \lambda_1^k q_1 + \alpha_2 \lambda_2^k q_2 + \dots + \alpha_n \lambda_n^k q_n) \\ &= \beta_k \lambda_1^k \left( \alpha_1 q_1 + \alpha_2 \left(\frac{\lambda_2}{\lambda_1}\right)^k q_2 + \dots + \alpha_n \left(\frac{\lambda_n}{\lambda_1}\right)^k q_n \right) \end{aligned}$$

$\Rightarrow v^{(k)} \rightarrow \pm q_1$  as  $k \rightarrow \infty$ ,  $\pm$  depends on  $\text{sign}(\lambda_1)$  ok

note

1. The error is reduced by a constant factor in each step.
2. If  $\alpha_1 = q_1^T v^{(0)} = 0$ , then  $v^{(k)} \rightarrow \pm q_2$ ,  $\lambda^{(k)} \rightarrow \lambda_2$ .
3. If  $A = \frac{1}{h^2} \text{trid}(-1, 2, -1)$ , the matrix form of  $-D_+ D_-$ , then  $w = Av$ , can be coded as a loop.

for  $i = 1 : n$ ;  $w_i = (-v_{i-1} + 2v_i - v_{i+1})/h^2$ ; end % assuming  $v_0 = v_{n+1} = 0$

This is more efficient than forming  $A$  and computing  $w = Av$  by direct matrix-vector multiplication.

question : How can the other eigenvalues be obtained?

### 4.3 inverse power method

idea : apply power method to  $A^{-1}$

1.  $Aq_i = \lambda_i q_i \Rightarrow A^{-1}q_i = \lambda_i^{-1}q_i$  : same e-vectors , reciprocal e-values

The largest e-value of  $A^{-1}$  is  $\lambda_n^{-1}$ .

2.  $w = A^{-1}v \Leftrightarrow Aw = v$

#### algorithm

1.  $v^{(0)}$  : given ,  $\|v^{(0)}\|_2 = 1$
2. for  $k = 1, 2, \dots$
3. solve  $Aw = v^{(k-1)}$       % e.g.  $LU$  factorization, etc.
4.  $v^{(k)} = w/\|w\|_2$
5.  $\lambda^{(k)} = (v^{(k)})^T A v^{(k)}$

ex :  $A = \begin{pmatrix} 2 & 1 & 1 \\ 1 & 3 & 1 \\ 1 & 1 & 4 \end{pmatrix} \Rightarrow \begin{cases} \lambda_1 = 5.214320 \\ \lambda_2 = 2.460811 \\ \lambda_3 = 1.324869 \end{cases} , \quad v^{(0)} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \cdot \frac{1}{\sqrt{3}}$

#### inverse power method

$k$	$\lambda^{(k)}$	$ \lambda^{(k)} - \lambda_3 $	$\frac{ \lambda^{(k)} - \lambda_3 }{ \lambda^{(k-1)} - \lambda_3 }$
0	5.000000	3.675131	-
1	3.816327	2.491457	0.677923
2	1.864903	0.540034	0.216754
$\downarrow$	$\downarrow$	$\downarrow$	$\downarrow$
$\infty$	$\lambda_3$	0	hw

#### summary

1. **power method** :  $\lambda^{(k)} \rightarrow \lambda_1$  , convergence factor =  $(\lambda_2/\lambda_1)^2$   
     **inverse power method** :  $\lambda^{(k)} \rightarrow \lambda_n$  , convergence factor = hw
2. How to compute the other  $\lambda_i$ ? Several methods are discussed in Math 571.
  - a) shifted inverse power method :  $(A - \mu I)^{-1}$
  - b)  $QR$  method  
 $A^{(0)} = Q^{(0)}R^{(0)}$  ,  $Q$  : orthogonal ( $Q^T Q = I$ ) ,  $R$  : upper triangular  
 $A^{(1)} = R^{(0)}Q^{(0)} = Q^{(1)}R^{(1)}$  : same e-values as  $A$   
 $\dots$   
 $A^{(k)} \rightarrow \text{diagonal}$
3. recall : What method does Matlab use to compute the roots of a polynomial?  
     answer : **type roots** , ...