

School of Informatics



Quantum Cyber Security Coursework

B232245
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Question 1

(a)

- $|0\rangle \rightarrow \Phi_q(|0\rangle\langle 0|) = (1-q)|0\rangle\langle 0| + q \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -i \end{bmatrix} = (1-q)|0\rangle\langle 0| + q \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} = |0\rangle\langle 0|$
- $|1\rangle \rightarrow \Phi_q(|1\rangle\langle 1|) = (1-q)|1\rangle\langle 1| + q \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -i \end{bmatrix} = (1-q)|1\rangle\langle 1| + q \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} = |1\rangle\langle 1|$
- $|+\rangle \rightarrow \Phi_q(|+\rangle\langle +|) = (1-q)|+\rangle\langle +| + \frac{q}{2} \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -i \end{bmatrix} = (1-q)|+\rangle\langle +| + \frac{q}{2} \begin{bmatrix} 1 & 1 \\ i & i \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -i \end{bmatrix} = (1-q)|+\rangle\langle +| + \frac{q}{2} \begin{bmatrix} 1 & -i \\ i & -1 \end{bmatrix} = |+\rangle\langle +| (1-q) + q|+_y\rangle\langle +_y|,$
since $|+_y\rangle\langle +_y| = \begin{bmatrix} 1 & -i \\ i & -1 \end{bmatrix}$
- $|-\rangle \rightarrow \Phi_q(|-\rangle\langle -|) = (1-q)|-\rangle\langle -| + \frac{q}{2} \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -i \end{bmatrix} = (1-q)|-\rangle\langle -| + \frac{q}{2} \begin{bmatrix} 1 & -1 \\ -i & i \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -i \end{bmatrix} = (1-q)|-\rangle\langle -| + \frac{q}{2} \begin{bmatrix} 1 & i \\ -i & -1 \end{bmatrix} = |-\rangle\langle -| (1-q) + q|-_y\rangle\langle -_y|,$
since $|-_y\rangle\langle -_y| = \begin{bmatrix} 1 & i \\ -i & -1 \end{bmatrix}$

(b)

- Probability that Alice sent the state $|0\rangle$ and that Bob measured $|1\rangle$:

$$\text{Tr}(\Phi_q(|0\rangle\langle 0|)|1\rangle\langle 1|) = \langle 1|\Phi_q(|0\rangle\langle 0|)|1\rangle = \langle 1|0\rangle\langle 0|1\rangle = 0$$

- Probability that Alice sent the state $|1\rangle$ and that Bob measured $|0\rangle$:

$$\text{Tr}(\Phi_q(|1\rangle\langle 1|)|0\rangle\langle 0|) = \langle 0|\Phi_q(|1\rangle\langle 1|)|0\rangle = \langle 0|1\rangle\langle 1|0\rangle = 0$$

- Probability that Alice sent the state $|+\rangle$ and that Bob measured $|-\rangle$:

$$\text{Tr}(\Phi_q(|+\rangle\langle +|)|-\rangle\langle -|) = \dots = \frac{q}{2}$$

- Probability that Alice sent the state $|-\rangle$ and that Bob measured $|+\rangle$:

$$\text{Tr}(\Phi_q(|-\rangle\langle -|)|+\rangle\langle +|) = \dots = \frac{q}{2}$$

$$\implies e_b = 0 \text{ and } e_p = \frac{\frac{q}{2} + \frac{q}{2}}{2} = \frac{q}{2}$$

(c)

$$Q = 1, \xi = 1, \Delta(n, \epsilon) = 0$$

$$\implies R_{BB84} = \frac{1}{2}(1 - h(e_b) - h(e_p)) = \frac{1}{2}(1 + \frac{q}{2} \log \frac{q}{2} + (1 - \frac{q}{2}) \log (1 - \frac{q}{2})),$$

since

$$e_b = 0 \implies h(e_b) = 0$$

$$e_p = \frac{1}{2} \implies h(e_p) = -\frac{q}{2} \log \frac{q}{2} - (1 - \frac{q}{2}) \log (1 - \frac{q}{2})$$

Secret key is possible to distil when $\frac{1}{2}(1 + \frac{q}{2} \log \frac{q}{2} + (1 - \frac{q}{2}) \log (1 - \frac{q}{2})) > 0$, or when $q = 0$. This holds $\iff q \in (0, 1) \cap (1, 2)$. Since q is a probability, it holds that it is possible to distil the secret key for any $q \in [0, 1)$.

Question 2

(a)

- $h(\frac{1}{8}) = -\frac{1}{8} \log \frac{1}{8} - \frac{7}{8} \log \frac{7}{8} = -\frac{1}{8}(-3) - \frac{7}{8} \log \frac{7}{8} = \frac{3}{8} - \approx 0.544$
- $h(\frac{1}{16}) = -\frac{1}{16} \log \frac{1}{16} - \frac{15}{16} \log \frac{15}{16} = -\frac{1}{16}(-4) - \frac{15}{16} \log \frac{15}{16} = \frac{1}{4} - \frac{15}{16} \log \frac{15}{16} \approx 0.337$

(b)

We first calculate the density matrix for ρ :

$$\rho = \frac{1}{2} |0\rangle \langle 0| + \frac{1}{2} |+\rangle \langle +| = \frac{1}{2} \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} + \frac{1}{2} \begin{bmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{bmatrix} = \begin{bmatrix} \frac{3}{4} & \frac{1}{4} \\ \frac{1}{4} & \frac{1}{4} \end{bmatrix}$$

To find the eigenvalues, we first calculate the determinant of $\begin{bmatrix} \frac{3}{4} - \lambda & \frac{1}{4} \\ \frac{1}{4} & \frac{1}{4} - \lambda \end{bmatrix}$:

$$(\frac{3}{4} - \lambda)(\frac{1}{4} - \lambda) - \frac{1}{16} = \frac{3}{16} - \lambda + \lambda^2 - \frac{1}{16} = \lambda^2 - \lambda + \frac{1}{8}$$

Then we see that $\lambda^2 - \lambda + \frac{1}{8} = 0 \iff \lambda = \frac{2 \pm \sqrt{2}}{4}$

Finally, we calculate the von Neumann entropy as follows:

$$H(\rho) = \frac{2 - \sqrt{2}}{4} \log(\frac{2 - \sqrt{2}}{4}) - \frac{2 + \sqrt{2}}{4} \log(\frac{2 + \sqrt{2}}{4}) \approx 0.122 + 0.059 = 0.181$$

(c)

$$\begin{aligned} \Phi(\rho) &= E_0 \rho E_0^\dagger + E_1 \rho E_1^\dagger = \\ &= \sqrt{p} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \frac{1}{2} \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} \sqrt{p} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} + \sqrt{1-p} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \frac{1}{2} \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} \sqrt{1-p} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} = \\ &= \frac{1}{4} \left(\begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} + \begin{bmatrix} 0 & 0 \\ 0 & 2 \end{bmatrix} \right) = \frac{1}{4} \begin{bmatrix} 1 & -1 \\ -1 & 3 \end{bmatrix} \end{aligned}$$

(d)

$$\begin{aligned}
F(\rho, \sigma) &= \text{Tr}(\sqrt{\rho^{\frac{1}{2}} \sigma \rho^{\frac{1}{2}}}) = \text{Tr}(\sqrt{\begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{bmatrix} \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}}) = \text{Tr}(\sqrt{\begin{bmatrix} 0 & 0 \\ 0 & \frac{1}{2} \end{bmatrix}}) = \frac{1}{\sqrt{2}} \\
D(\rho, \sigma) &= \frac{1}{2} \text{Tr}(|\rho - \sigma|) = \frac{1}{2} \text{Tr}(|\begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} - \begin{bmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{bmatrix}|) = \frac{1}{2} \text{Tr}(|\frac{1}{2} \begin{bmatrix} -1 & -1 \\ -1 & 1 \end{bmatrix}|) = \\
&= \frac{1}{2} \text{Tr}(\sqrt{\frac{1}{4} \begin{bmatrix} -1 & -1 \\ -1 & 1 \end{bmatrix} \begin{bmatrix} -1 & -1 \\ -1 & 1 \end{bmatrix}}) = \frac{1}{2} \text{Tr}(\sqrt{\begin{bmatrix} \frac{1}{2} & 0 \\ 0 & \frac{1}{2} \end{bmatrix}}) = \frac{1}{2} \text{Tr}(\begin{bmatrix} \sqrt{\frac{1}{2}} & 0 \\ 0 & \sqrt{\frac{1}{2}} \end{bmatrix}) = \frac{1}{\sqrt{2}}
\end{aligned}$$

The maximum probability that Charlie can correctly identify the state is $\frac{1}{2}(1 + D(\rho, \sigma)) = \frac{1}{2}(1 + \frac{1}{\sqrt{2}}) \approx 0.854$.

Question 3

(a)

The following is an example of the Mermin-Peres magic square game. It is known that the best classical strategy can ensure the win probability of $\frac{8}{9}$. In order to satisfy the rules, the players can choose the following sets of values:

- Alice can either fill the row with the numbers 1,-1,-1 or 1,1,1.
- Bob can either fill the column with the numbers: 1,1,-1 or -1,-1,-1

In order to maximize the probability of winning, they must pick a set of numbers where most of the match is. There are two equivalent solutions:

- Alice picks 1,-1,-1 and Bob picks -1,-1,-1
- Alice picks 1,1,1 and Bob picks 1,1,-1

In both cases, there are five equal values and one different. Let us assume that Alice and Bob decide: Alice picks 1,-1,-1 and Bob picks -1,-1,-1. Since they always share one cell, the probability of winning is the same for any of the nine scenarios. Let us observe the scenario where Alice gets the first row, and Bob gets the first column. As we can see, the only scenario where they lose the game is:

1,-1	-1	-1
1		
1		

There are nine possible outcomes (since we need to distinguish between all the values), and the probability of them losing is $\frac{1}{9}$. Since that is the case for any row and column Alice and Bob get, the final probability of them winning the game is $\frac{8}{9}$.

(b)

As it was shown in many papers, if the players are allowed to share an entangled quantum state, it is possible for them to win the magic square game with certainty. In order to see why this strategy is valid (the products of row and column values are correct), let us consider the following:

- The three matrices in any one row, or in any one column, of this grid all commute with each other. That is, we can multiply any two or three of them together in any order, and the result will not be affected by the order.
- If we multiply all three matrices in any row, we always get the identity matrix.
- If we multiply all three matrices in any column, we always get the opposite of the identity matrix.
- For each of the nine matrices, all of their eigenvectors have eigenvalues of 1 or -1 .

As a result, each cell within the specified context is filled with either a 1 or a 0, where the product of the row cells is equivalent to 1 and the product of the column cells is equivalent to -1 .

Since all the matrices commute with one another, it is possible to identify four orthogonal vectors in the 4-dimensional space that serve as simultaneous eigenvectors of all three matrices at once. For instance, considering the third row, the standard basis vectors:

$$\{(1, 0, 0, 0), (0, 1, 0, 0), (0, 0, 1, 0), (0, 0, 0, 1)\}$$

can serve as eigenvectors for all three matrices, and the four corresponding eigenvalues can be obtained from the diagonals of each of the three matrices consecutively:

$$\{1, -1, 1, -1\}$$

$$\{1, -1, -1, 1\}$$

$$\{1, 1, -1, -1\}$$

Therefore, we can construct a single measurement on the quantum system that tells us all three quantities measured by the three matrices in the row or column.

In order to understand how the answers of Alice and Bob are correlated, they must share maximally entangled states. A maximally entangled state possesses a specific type of symmetry such that if Ψ is maximally entangled and U is any unitary operation, then $U \otimes U^\dagger |\Psi\rangle = |\Psi\rangle$. Therefore, if Bob measures in the same basis as Alice, their outcomes will be fully correlated. Therefore even if Alice and Bob use different measurement choices, their outcomes at the intersections will agree.

Assuming that A represents any observable applicable to Alice's system, and B corresponds to the corresponding observable for Bob's system, the correlation value of $\langle \Psi | AB | \Psi \rangle = 1$ ensures that the players win the game every time.

(c)

The protocol can work as follows:

1. Any trusted or untrusted party distributes to Alice and Bob n copies of the state Ψ
2. Alice measures chooses a random row from the grid and measures in all of the three observables $x^{(i)} = (x_1^{(i)}, x_2^{(i)}, x_3^{(i)})$. She obtains triplets of results $a^{(i)} = (a_1^{(i)}, a_2^{(i)}, a_3^{(i)})$. It holds that $a_j^{(i)} \in \{-1, 1\}$ for any j . She stores pairs of $(a^{(i)}, x^{(i)})_n$
3. Bob measures chooses a random row from the grid and measures in all of the three observables $y^{(i)} = (y_1^{(i)}, y_2^{(i)}, y_3^{(i)})$. He obtains triplets of results $b^{(i)} = (b_1^{(i)}, b_2^{(i)}, b_3^{(i)})$. It holds that $b_j^{(i)} \in \{-1, 1\}$ for any j . He stores pairs of $(b^{(i)}, y^{(i)})_n$
4. Alice and Bob announce the bases $x^{(i)}, y^{(i)}$ and they keep positions where they used the same basis
5. If there was no eavesdropping then $a^{(i)} = b^{(i)} \forall i$ of the raw key

Intuition: If non-locality exists, then it is impossible for Eve to have a perfect correlation with Alice's string. This is due to the monogamy of entanglement, which is valid for maximum violation and applies to any violation because the presence of local hidden variables would be implied by perfect correlation.