

Continuation passing style (CDS)

- (ontimation - passing is a technique where functions don't return a value, but instead call a given function (callback) to produce the next result.

process: a -> b

vs. process': a -> (cb: (q->b)) => b

or type (ont(r,a) = (a->r) -> r

process : a -> Cont < b, q >

-we can think of continuations as suspended computations: they need a further function to complete.

- Another way of bothing at the situation is that a value is meaningless until we apply (use) it:

let Rive f = } 5 (let Rive = (1>) 5)

do live (printfin "/. 2")

let add x y = x + y
let square x = x * x

let pythagares x y = add (square x) (square y)

let add' x y = fun f -> f (add x y)

let squae' x = fun f -> f (square x)

let pythagaras' x y = fun f ->

Square' x (fun x' ->

Square' y (fun y' ->

add' x' y' f.).

- Example: thrice

let thrice (f: a-a) -> (x:a)-> a $= f(f(f_X))$

= XX + D+ D +

= 1 >> 1>> 1

let thrice' (1: a -> (a->r)->r)-> (x: a) -> (a->r)->+)=
or let thrice' (1: a-> (ontco))-> (x:a) -> (ontco) =

fun k -> fl x

(Pun fx > f'fx

(fun ffx -> f'ffx k)

- We see a pottern! Let's abstract!

let chain (s: (a>r)>r)>

(f:(a→(b→r)→r)→ (b→r)→r=

or let chain (s: Cont(as)) -> (f: a -> (out(b)) -> (out(b))

= fun k -> s'(fun x -> f x k)

= Cont. bind!

let Unrice 1 x = (1)x >= 3 >= 3 >= 3

- (out <r, a) is also a lunctor in a:

(aul, map: (f: a > b) > (x: (out <r, a>) > (out <r, b)

(out, map $f \times = fvn \quad g \rightarrow x \quad (fvn \quad a \rightarrow g \quad (f \quad a))$ (b > r)

- (out, return $x = fvn \quad f \rightarrow f \quad x$ (a-r) $\rightarrow r = (out \langle r, a \rangle)$

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Now we can create a computation expression:

type (ont Builder () =

member Bind (a,b) = a >>= b

member Return x = Cout. neturn x

member Roturntrom x = lun k -> x k

let cont = ContBuilder () == 5 = +01

let thrice f x =

(out {

let! x' = f x

let! $\chi'' = f x'$

3 return! f x"



- Continuations are of great practical and theoretical importance
 - The continuation monad is the mother of all monads: every monad can be generated from it! (see Dan Piponi (sig & pe))
- There is one type of continuation which is superbly more common than others:

type Async (a) = (a -> (1) -> ()

- Since (out <(), a) = Async<'a) doesn't return
 anything useful (it penforms IO or mutation)
 it can be scheduled to run on a separate
 thread, and return immeliately
- Agync is sundemental for multi-thrading, concurrency, parallelism, and reactive programming.
- The async & I workflow makes it easy and natural to work with asynchronous computations

This works super for pure functions

INF-3910 Lecture 11



Managing shared state using Mailbox Processo

- Pure sunctions ga a long way, but something we need communication between threads.

It we are careless, this can become a disaster

elegantly (bring a serial region!)

let private mutable state = 0

let agent = MailboxProcessor. Start (fun inbox ->

let rec loop () = async {
let! msg = inbox. Receive ()

State < Msg return! loop()

3

let volate state x = ayent. Post x