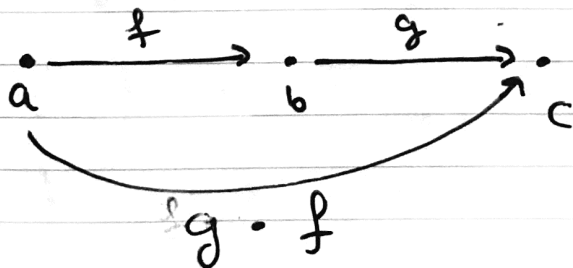
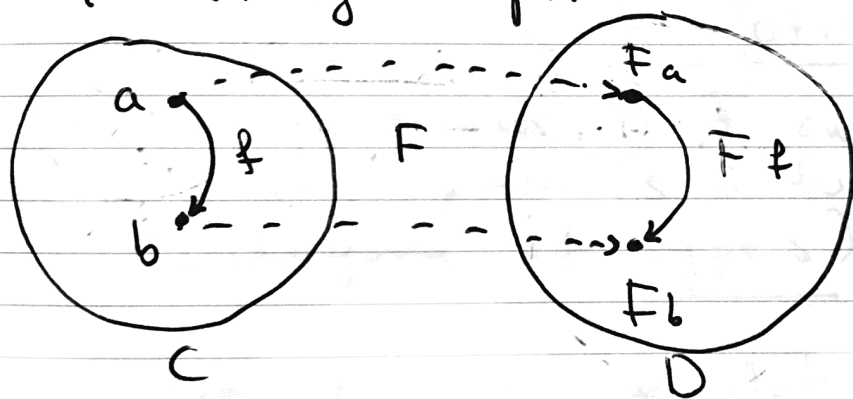


## Functors

- Functors are a simple, yet very powerful abstraction in functional programming
- We have already seen how map abstracts looping for `IEnumerable`s (list, array, seq)
- But functors are much more general than just looping! we can map over a lot of different structures!
- In category theory a category is defined in terms of objects, morphisms and composition of morphisms (think functions and elements):



- In category theory a functor is a function between two categories which maps every object in  $C$  to an object in  $D$ .
- Since morphisms (functions) are also objects, all functions are also mapped. This means that all connections between objects are preserved!
- Functors also preserve composition and identity morphisms:



note! this is very different from  $F(a \rightarrow b)$ !

$$F(a \rightarrow Fa)$$

$$F(f: a \rightarrow b) \rightarrow (Ff: Fa \rightarrow Fb)$$

$$F(g \cdot f) \rightarrow Fg \cdot Ff$$

$$F(id_a) \rightarrow id_{Fa}$$

Functor laws

← preserve identity

Note! Functors need not be injective!

They can collapse structure!

- ok, enough mathematical nonsense! What!?
- In programming, every time we have a generic type, we potentially have a functor!

type Option <'a> = | Some 'a | Nothing

↙  
this is a type function (constructor), taking a type 'a and returning a new type!

Option : 'a  $\rightarrow$  Option 'a

Compare this to F on the previous page. Option is a functor (in the category of types)!

- We can create an Option <int> from an int by using one of the data constructors.

let x = Some 42 : Option <int>

let y = Nothing : Option <int>

- But what about functions?

Some f : Option <int  $\rightarrow$  int>

is very different from

f' : Option <int>  $\rightarrow$  Option <int> !

- How do we transform  $f: a \rightarrow b$  to  $f': \text{Option}\langle a \rangle \rightarrow \text{Option}\langle b \rangle$ ?

- The answer is: we define a function usually called map (or fmap)!

let  $\text{map} (f: a \rightarrow b) (x: \text{Option}\langle a \rangle) : \text{Option}\langle b \rangle =$   
match  $x$  with

| Some  $v \rightarrow \text{Some} (f v)$

| Nothing  $\rightarrow \text{Nothing}$

- map is a function which has the following signature:

$\text{map} : (f: a \rightarrow b) \rightarrow \text{Option}\langle a \rangle \rightarrow \text{Option}\langle b \rangle$

which is exactly what we are looking for.

- But what about composition and identity?

$\text{fmap id} = \text{id} \quad (\text{id}: 'a \rightarrow 'a)$

$\text{fmap id Nothing} = \text{Nothing} = \text{id Nothing}$

$\text{fmap id (Some } x) = \text{Some (id } x)$

$= \text{Some } x$

$= \text{id (Some } x)$

composition:  $[fmap\ g \ll fmap\ f = fmap\ (g \ll f)]$

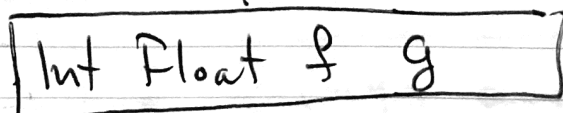
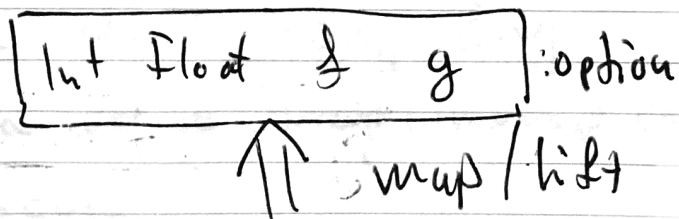
$$\begin{aligned} [fmap\ (g \ll f)\ Nothing &= Nothing \\ &= fmap\ g\ Nothing = fmap\ g\ (fmap\ f\ Nothing) \\ &= (fmap\ g \ll fmap\ f)\ Nothing \end{aligned}$$

$$\begin{aligned} [fmap\ (g \ll f)\ (Some\ x) &= Some\ (g\ (f\ x)) \\ &= fmap\ g\ (Some\ (f\ x)) \\ &= fmap\ g\ (fmap\ f\ (Some\ x)) \\ &= (fmap\ g \ll fmap\ f)\ (Some\ x) \end{aligned}$$

□

Note that for these relations to hold  $f$  and  $g$  must be pure, stateless functions!

- map/fmap should really be called lift instead, as it lifts functions into an "elevated world":



- lift can be generalized to lift N

- How is all this relevant to programmers?
- There is a very common pattern where we unwrap (pattern match) a type, apply a transformation and wrap it up again. This is handled much more elegantly using map!

Example:

```
[ let cust : Option<Customer> = get Customer 123  
  match cust with  
  | Some c → Some <| process Customer c  
  | None → None
```

vs.

```
[ Option.map process Customer (get Customer 123)
```

In particular, when we start composing transformations we get a lot of wrap/unwrap code. Functors to the rescue!

Keep your eyes open, these things are everywhere!

Examples of functors:

$\text{Option}\langle a \rangle$

-  $\text{List}\langle a \rangle$ :  $\text{map } f \ x =$

match  $x$  with

$|\ [] \rightarrow []$

$| h :: t \rightarrow f \ h :: \text{map } f \ t$

$\text{Identity}\langle a \rangle$ :  $\text{map } f = \text{function } | \text{Id } x \rightarrow \text{Id } (f \ x)$

$\text{Const}\langle a, b \rangle$

$\text{Const } o \ f \ a$

$\text{map } f =$

$| \text{Const } x \rightarrow \text{Const } x$

Functions! type  $\text{Fun}\langle a, r \rangle = \text{Fun } (a \rightarrow r)$

$\text{map } (f: r \rightarrow s) \ (x: a \rightarrow r) : a \rightarrow s$

let  $\text{map } f \ x =$

match  $x$  with

$| \text{Fun } g \rightarrow g \gg f \quad (= f \ll g)$

or simply let  $\text{map} = (\ll)$

$\text{map}$  is simply function composition!

- The function functor is also called the Reader functor