

Indian Institute of Technology Roorkee

CHN-323

Computer Applications in Chemical Engineering

Ashwini Kumar Sharma

Department of Chemical Engineering
Indian Institute of Technology Roorkee

Email: ashwini.fch@iitr.ac.in



Set of equations

- Determine the values $x_1, x_2, x_3, \dots, x_n$ that simultaneously satisfy a set of equations

$$\begin{aligned} f_1(x_1, x_2, \dots, x_n) &= 0 \\ f_2(x_1, x_2, \dots, x_n) &= 0 \\ &\vdots \\ f_n(x_1, x_2, \dots, x_n) &= 0 \end{aligned}$$

- Equations can be linear/nonlinear

Set of linear algebraic equations

$$a_{11} x_1 + a_{12} x_2 + \cdots + a_{1n} x_n = b_1$$

$$a_{21} x_1 + a_{22} x_2 + \cdots + a_{2n} x_n = b_2$$

\vdots

$$a_{n1} x_1 + a_{n2} x_2 + \cdots + a_{nn} x_n = b_n$$

- a 's are constant coefficients
- b 's are constants
- x 's are unknowns
- n is the number of equations

Importance of linear equations

- Linear equations are the basis for mathematical models of
 - economics,
 - weather prediction,
 - heat and mass transfer,
 - statistical analysis, and
 - a myriad of other applications.
- The methods for solving ordinary and partial differential equations depend on them

Set of linear equations and matrix algebra

$$a_{11} x_1 + a_{12} x_2 + \cdots + a_{1n} x_n = b_1$$

$$a_{21} x_1 + a_{22} x_2 + \cdots + a_{2n} x_n = b_2$$

\vdots

$$a_{n1} x_1 + a_{n2} x_2 + \cdots + a_{nn} x_n = b_n$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}$$

$$AX = B$$

$$\text{where } A = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix}, \quad X = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}, \quad B = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}$$

Matrix algebra

- Square matrix
- Symmetric matrix
- Diagonal matrix
- Identity matrix
- Upper triangular matrix
- Lower triangular matrix

Matrix algebra

- Matrix addition and subtraction
 - Commutative $[A] + [B] = [B] + [A]$
 - Associative $([A] + [B]) + [C] = [A] + ([B] + [C])$
- Multiplication of a matrix with a scalar
- Multiplication of two matrices
 - Associative $([A][B])[C] = [A]([B][C])$
 - Distributive $[A]([B] + [C]) = [A][B] + [A][C]$
 - Not commutative $[A][B] \neq [B][A]$

Determinant of a matrix

$$[A] = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix}$$

$$D = \begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix}$$

$$|A| = \begin{vmatrix} a & b \\ c & d \end{vmatrix} = (ad - bc)$$

$$\begin{aligned} |A| &= \begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix} = a \begin{vmatrix} e & f \\ h & i \end{vmatrix} - b \begin{vmatrix} d & f \\ g & i \end{vmatrix} + c \begin{vmatrix} d & e \\ g & h \end{vmatrix} \\ &= a(ei - fh) - b(di - fg) + c(dh - eg) \end{aligned}$$

Solving linear equations in MATLAB

➤ $AX = B$

```
>> a=[1 2; 3, 4];  
>> b=[3, 5];  
>> b=b';  
>> x=a\b
```

x =

-1
2

```
>> x=inv(A)*b
```

Unrecognized function or variable 'A'.

Did you mean:

```
>> x=inv(a)*b
```

x =

-1.0000
2.0000

Solving linear equation systems

- Three types of methods
 - Direct methods
 - Iterative methods
 - Decomposition methods

Solving linear equation systems

- A is a $n \times n$ matrix

$$[A][A]^{-1} = [A]^{-1}[A] = [I]$$

- Our system, $AX = B$
- Multiply both sides by A^{-1}

$$A^{-1}AX = A^{-1}B$$

$$X = A^{-1}B$$

$$A^{-1} = \frac{1}{|A|} \cdot \text{Adj } A$$

- $\text{Adj } A$ means the adjoint matrix of A
- The adjoint of a matrix is the transpose of the cofactor element matrix of the given matrix.

Cramer's rule

- Our system $AX = B$
- Solution is given by

$$x_j = \frac{\det(A_j)}{\det(A)} \quad ; \det(A) \neq 0 \quad ; j = 1, \dots, n$$

- A_j is obtained by replacing the j^{th} column of A by B

Cramer's rule

- Total number of determinants required = $n+1$
- Number of operations to calculate determinant of a $n \times n$ matrix = $(n-1)(n!)$
- Total operations in Cramer's rule = $(n+1)(n-1)(n!)$
 $\sim n^2 n!$

For a 100×100 matrix: $100^2 100! = 10^{162}$ calculations

Gauss elimination

- Consider the system

$$x_1 + x_2 + 2x_3 = 3$$

$$2x_1 + 3x_2 + x_3 = 2$$

$$3x_1 - x_2 - x_3 = 6$$

- Let us perform the following row operations

$$\text{Old Eq (2)} - 2 \text{ Eq (1)} \rightarrow \text{New Eq (2): } x_2 - 3x_3 = -4$$

$$\text{Old Eq (3)} - 3 \text{ Eq (1)} \rightarrow \text{New Eq (3): } -4x_2 - 7x_3 = -3$$

- The equivalent system of equations is

$$x_1 + x_2 + 2x_3 = 3$$

$$x_2 - 3x_3 = -4$$

$$-4x_2 - 7x_3 = -3$$

We have **eliminated** x_1 from Eq. (2) and (3).

Gauss elimination contd...

- Let us perform another elementary row operation

Old Eq (3) + 4 Eq (2) → New Eq (3): $-19x_3 = -19$

- The new equivalent system of equations is

$$x_1 + x_2 + 2x_3 = 3$$

$$0x_1 + x_2 - 3x_3 = -4$$

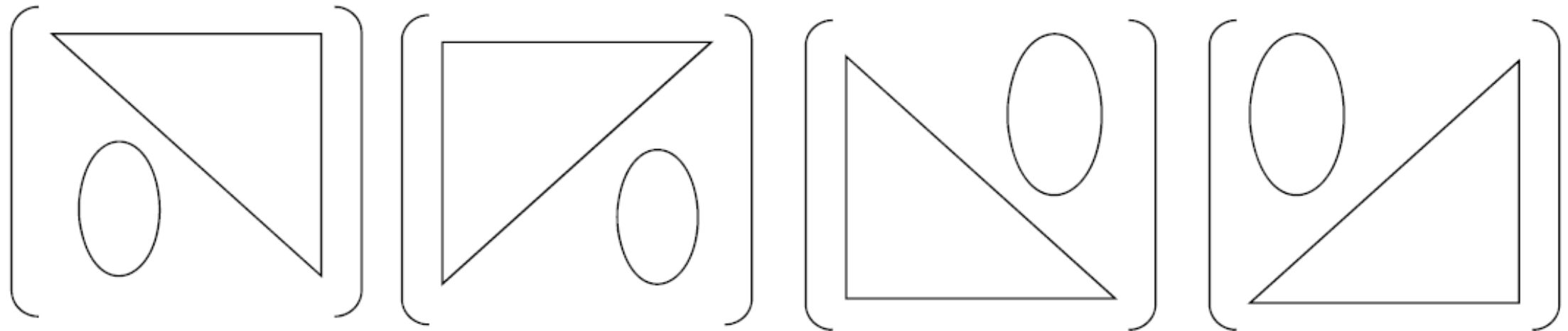
$$0x_1 + 0x_2 - 19x_3 = -19$$

We have **eliminated** x_2 from Eq. (3).

- Now, if we write in matrix format, the coefficient matrix is **upper triangular matrix**
 - That can be easily solved
 - Solve Eq. (3) to get $x_3 = 1$
 - Substitute $x_3 = 1$ in Eq. (2) to get $x_2 = -1$
 - Substitute $x_3 = 1$ and $x_2 = -1$ in Eq. (1) to get $x_1 = 2$

Gauss elimination contd...

- Idea behind: Convert "full" coefficient matrix into a lower or upper triangular matrix (Note: the constant matrix on RHS will also be modified)



Upper
Triangular Structure

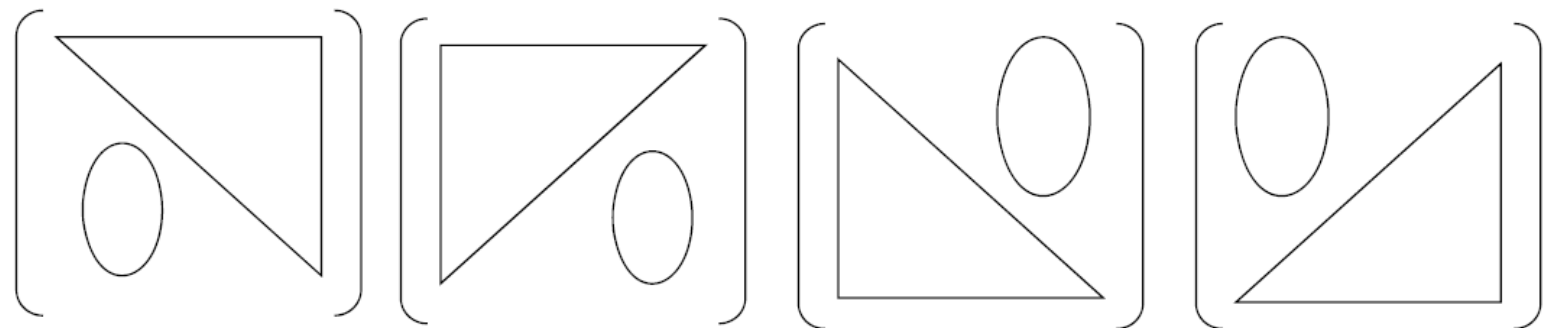
Reverse Upper
Triangular Structure

Lower
Triangular Structure

Reverse Lower
Triangular Structure

Gauss elimination contd...

- If coefficient matrix is
 - Upper triangular, the system can be solved in sequence x_n, x_{n-1}, \dots, x_1 using equations in order $n, n-1, \dots, 1$
 - Reverse upper triangular, the system can be solved in sequence x_1, x_2, \dots, x_n using equations in order $n, n-1, \dots, 1$
 - Lower triangular, the system can be solved in sequence x_1, x_2, \dots, x_n using equations in order $1, 2, \dots, n$
 - Reverse lower triangular, the system can be solved in sequence x_n, x_{n-1}, \dots, x_1 using equations in order $1, 2, \dots, n$



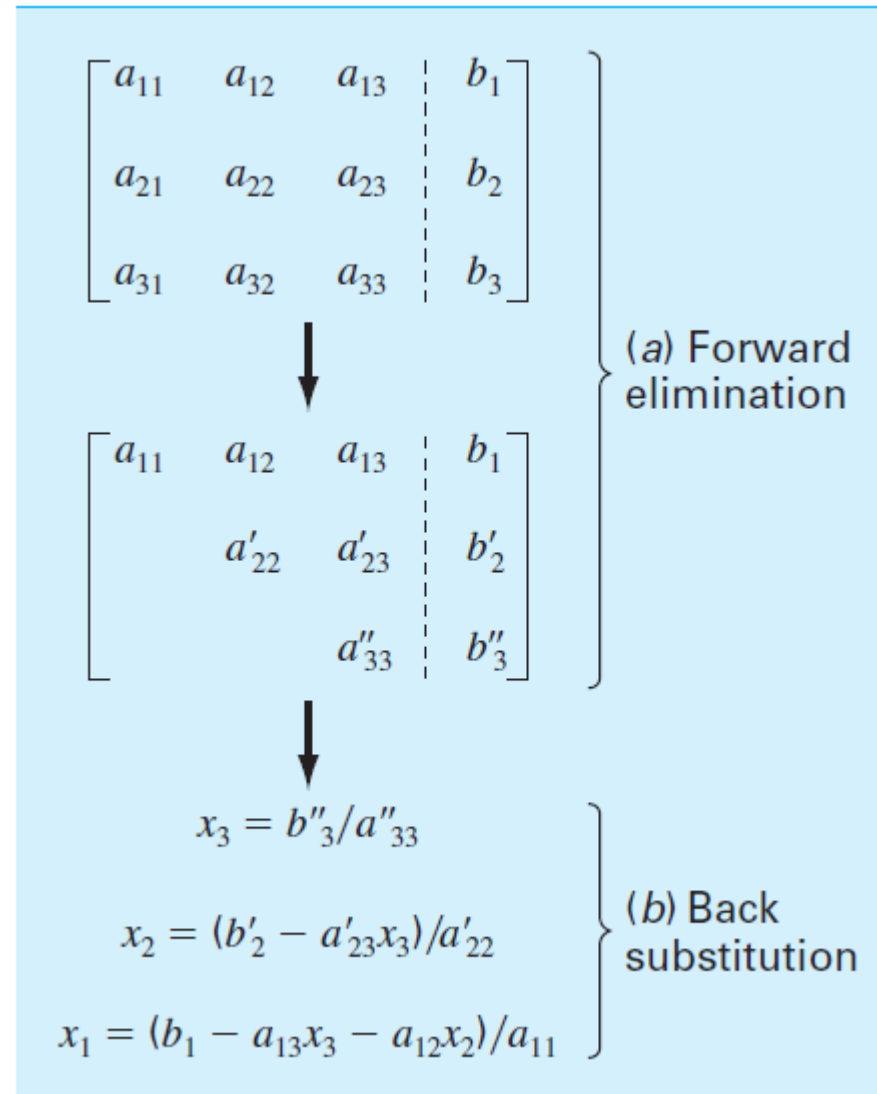
Upper
Triangular Structure

Reverse Upper
Triangular Structure

Lower
Triangular Structure

Reverse Lower
Triangular Structure

Two stages of Gauss elimination



Forward elimination: generalization

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1$$

$$a_{21}x_1 + a_{22}x_2 + a_{23}x_3 = b_2$$

$$a_{31}x_1 + a_{32}x_2 + a_{33}x_3 = b_3$$

- First, we eliminate x_1 from Eq. (2) and (3). To do so, we perform the following row operations

$$\text{New Eq (2)} = \text{Old Eq (2)} - (a_{21}^* \text{ Old Eq (1)} / a_{11})$$

$$\text{New Eq (3)} = \text{Old Eq (3)} - (a_{31}^* \text{ Old Eq (1)} / a_{11})$$

- a_{11} is the **pivot element** in the elimination of x_1 from Eq (2)&(3)

Forward elimination: generalization

- The new equivalent system is

$$\begin{aligned}a_{11}x_1 + a_{12}x_2 + a_{13}x_3 &= b_1 \\ \left(a_{22} - \left(\frac{a_{21}a_{12}}{a_{11}}\right)\right)x_2 + \left(a_{23} - \left(\frac{a_{21}a_{13}}{a_{11}}\right)\right)x_3 &= b_2 - \frac{a_{21}b_1}{a_{11}} \\ \left(a_{32} - \left(\frac{a_{31}a_{12}}{a_{11}}\right)\right)x_2 + \left(a_{33} - \left(\frac{a_{31}a_{13}}{a_{11}}\right)\right)x_3 &= b_3 - \frac{a_{31}b_1}{a_{11}}\end{aligned}$$

- In other words,

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1$$

$$a'_{22}x_2 + a'_{23}x_3 = b'_2$$

$$a'_{32}x_2 + a'_{33}x_3 = b'_3$$

Diagonal elements are pivot elements

- Let us perform another elementary row operation

$$\text{New Eq (3)} = \text{Old Eq (3)} - (a'_{32} * \text{Old Eq (2)} / a'_{22})$$

a'_{22} is the **pivot element** in the elimination of x_2 from Eq (3)

Forward elimination: generalization

- The new equivalent system is

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1$$

$$a'_{22}x_2 + a'_{23}x_3 = b'_2$$

$$\left(a'_{33} - \left(\frac{a'_{32}a'_{23}}{a'_{22}} \right) \right) x_3 = b'_3 - \frac{a_{32}b_2}{a_{22}}$$

- In other words,

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1$$

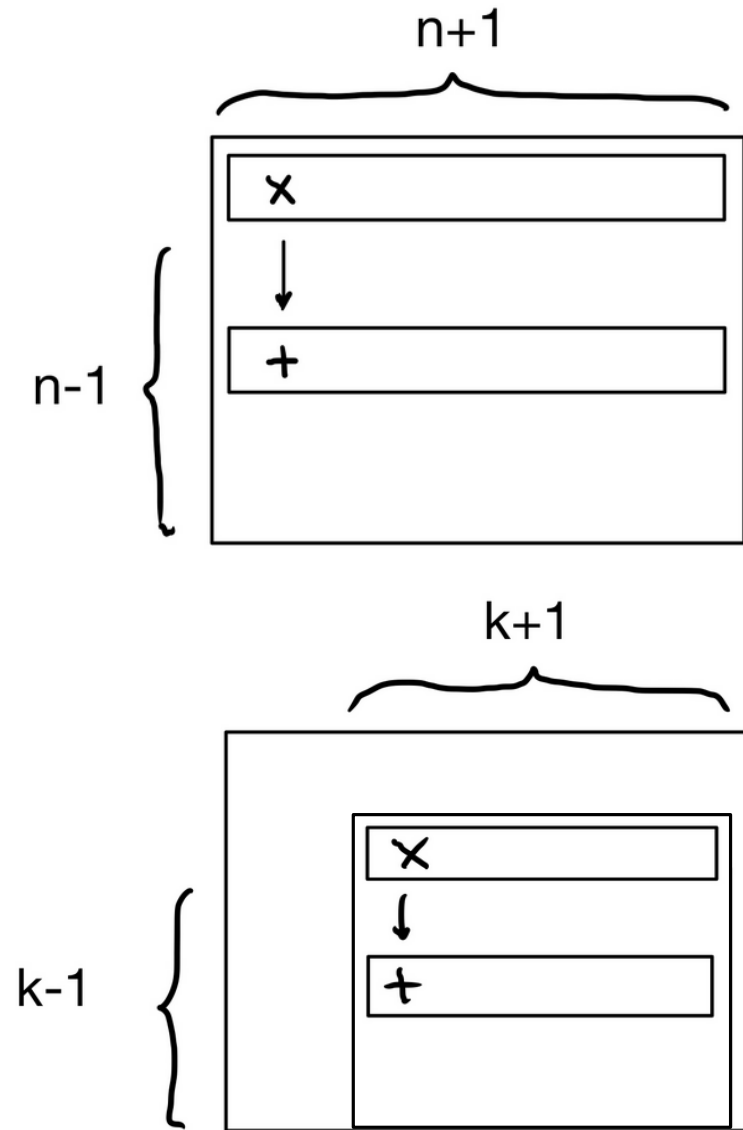
$$a'_{22}x_2 + a'_{23}x_3 = b'_2$$

$$a''_{33}x_3 = b''_3$$

This is an upper triangular system and can be solved in sequential order x_3 , x_2 and x_1

Forward elimination: Number of calculations

- For n equations in n unknowns, the augmented matrix is a $n \times (n+1)$ matrix.
- For each row i
 - Add a multiple of i^{th} row to all rows below it
$$\text{New Eq (2)} = \text{Old Eq (2)} - (a_{21} * \text{Old Eq (1)} / a_{11})$$
- For row 1, this process will require
 - There are $(n-1)$ rows below row 1, each of those has $(n+1)$ elements
 - Each element requires 1 multiplication, and 1 subtraction
 - Total calculations = $2 (n-1) (n+1) = 2n^2 - 2$
- When operating on i^{th} row, there are $k = n - i + 1$ unknowns, so there are $2k^2 - 2$ calcs required



Forward elimination: Number of calculations

- k ranges from n down to 1
- So, the total number of arithmetic opns required

$$\begin{aligned}\sum_{k=1}^n (2k^2 - 2) &= 2 \left(\sum_{k=1}^n k^2 - \sum_{k=1}^n 1 \right) \\ &= 2 \left(\frac{n(n+1)(2n+1)}{6} - n \right) \\ &= \frac{2}{3}n^3 + n^2 - \frac{5}{3}n\end{aligned}$$

Back-substitution: generalization

➤ Consider the system as $Lx=b$

$$\begin{bmatrix} l_{11} & 0 & \cdots & 0 \\ l_{21} & l_{22} & \ddots & 0 \\ \vdots & \vdots & \ddots & 0 \\ l_{n1} & l_{n2} & \cdots & l_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$

$$\det(L) = l_{11} * l_{22} * \cdots l_{nn} = \prod_{i=1}^n l_{ii}$$

For a non-singular L , all l_{ii} should be non-zero

Back-substitution: generalization

$$l_{11}x_1 = b_1 \rightarrow x_1 = \frac{b_1}{l_{11}}$$

$$l_{21}x_1 + l_{22}x_2 = b_2 \rightarrow x_2 = \frac{b_2 - l_{21}x_1}{l_{22}}$$

And so on

$$\begin{bmatrix} l_{11} & 0 & \cdots & 0 \\ l_{21} & l_{22} & \ddots & 0 \\ \vdots & \vdots & \ddots & 0 \\ l_{n1} & l_{n2} & \cdots & l_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$

$$l_{n1}x_1 + l_{n2}x_2 + \cdots + l_{n,n-1}x_{n-1} = b_n \rightarrow$$

$$x_n = \frac{b_n - l_{n1}x_1 - l_{n2}x_2 - \cdots - l_{n,n-1}x_{n-1}}{l_{nn}}$$

In general,
$$x_k = \frac{b_k - l_{k1}x_1 - l_{k2}x_2 - \cdots - l_{k,k-1}x_{k-1}}{l_{kk}}$$

Back-substitution: No. of calculations

$$x_1 = \frac{b_1}{l_{11}}$$

$$x_2 = \frac{b_2 - l_{21}x_1}{l_{22}}$$

$$x_k = \frac{b_k - l_{k1}x_1 - l_{k2}x_2 - \dots - l_{k,k-1}x_{k-1}}{l_{kk}}$$

To get $x_1 \rightarrow$ one division

To get $x_2 \rightarrow$ one division + one subtraction + one multiplication

To get $x_3 \rightarrow$ one division + two subtractions + two multiplications

....

To get $x_n \rightarrow$ one division + (n-1) subtractions + (n-1) multiplications

Total calculations to get $x_1, x_2, \dots, x_n \rightarrow$

n divisions + (1+2+...+(n-1)) subtractions + (1+2+...+(n-1)) multiplications

i.e., approximately n^2 calculations

Gauss elimination: with matrix format

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1$$

$$a_{21}x_1 + a_{22}x_2 + a_{23}x_3 = b_2$$

$$a_{31}x_1 + a_{32}x_2 + a_{33}x_3 = b_3$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & b_1 \\ a_{21} & a_{22} & a_{23} & b_2 \\ a_{31} & a_{32} & a_{33} & b_3 \end{bmatrix}$$

Augmented matrix

$$\begin{bmatrix} 1 & 1 & 2 & 3 \\ 2 & 3 & 1 & \vdots & 2 \\ 3 & -1 & -1 & 6 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 2 & 3 \\ 0 & 1 & -3 & \vdots & -4 \\ 0 & -4 & -7 & -3 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 2 & 3 \\ 0 & 1 & -3 & \vdots & -4 \\ 0 & 0 & -19 & -19 \end{bmatrix}$$

$$\rightarrow x = \begin{bmatrix} 2 \\ -1 \\ 1 \end{bmatrix}$$

Example: linearly dependent system

$$\begin{bmatrix} 1 & 2 & -1 & -2 & -1 \\ 2 & 1 & 1 & -1 & 4 \\ 1 & -1 & 2 & 1 & 5 \\ 1 & 3 & -2 & -3 & -3 \end{bmatrix} \Rightarrow \begin{bmatrix} 1 & 2 & -1 & -2 & -1 \\ 0 & -3 & 3 & 3 & 6 \\ 0 & -3 & 3 & 3 & 6 \\ 0 & 1 & -1 & -1 & -2 \end{bmatrix}$$

The **rank** of a matrix is defined as (a) the maximum number of linearly independent *row* vectors in the matrix or (b) the maximum number of linearly independent *column* vectors in the matrix.

Rank of augment matrix = rank of coefficient matrix = 2

Number of unknowns = 4

Example: pivot element is zero

$$A = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 1 & 2 \\ 4 & 5 & 6 \end{bmatrix}$$

$a_{11} = 0$, cannot be used for pivoting

Perform row interchange. Interchange rows (1) and (2)

$$A = \begin{bmatrix} 1 & 1 & 2 \\ 0 & 1 & 1 \\ 4 & 5 & 6 \end{bmatrix} \Rightarrow \begin{bmatrix} 1 & 1 & 2 \\ 0 & 1 & 1 \\ 0 & 1 & -2 \end{bmatrix} \Rightarrow \begin{bmatrix} 1 & 1 & 2 \\ 0 & 1 & 1 \\ 0 & 0 & -3 \end{bmatrix}$$

Coefficient Matrix is Upper Triangular

Rule: If pivot element = 0, then interchange that row with any of the later rows which will have a non-zero pivot element.

Example: pivot element is too small

- Sometimes, the candidate pivot element is quite small but not zero. Though, this does not lead to the “mathematical breakdown” of the algorithm, it does have a detrimental effect on the quality of the obtained solution when implemented on a finite precision computer.

$$\begin{bmatrix} 0.0001 & 0.5 \\ 0.4 & -0.3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 0.5 \\ 0.1 \end{bmatrix}$$

Our PC does 4-digit rounded decimal arithmetic

$$\left[\begin{array}{cc|c} 0.0001 & 0.5 & 0.5 \\ 0.4 & -0.3 & 0.1 \end{array} \right] \longrightarrow \left[\begin{array}{cc|c} 1 & 5000 & 5000 \\ 0 & -2000 & -2000 \end{array} \right]$$

Rounded value of -2000.3

Rounded value of -1999.9

Rule: Avoid not only “zero” for the pivot element but also small pivots.

$x_1 = 0$
 $x_2 = 1$ **Wrong Solution**

Pivoting

- Partial Pivoting: Make the element that is largest in magnitude from the remaining rows as the pivot element
- Threshold Pivoting: Perform interchange of rows only if the natural pivot element is significantly smaller than the competitor
- Complete Pivoting: Largest element in the entire remaining matrix is made the pivot element. This strategy therefore involves interchange of both rows and columns.