

## Part III — Additive Combinatorics

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*A brief summary of important ideas and results in the course*

## 0 Introduction

Lecture 1

**Notation.**  $[N] = 1, 2, \dots, N$

**Theorem** (Szemerédi's theorem).  $\forall \delta > 0, \forall k, \exists N \forall A \subset [N]$  s.t.  $|A| \geq \delta N$  and  $A$  contains an arithmetic progression of length  $k$ .

This theorem looks like it might be nicely solvable or have a counter-example, but we need hard machinery to attack it. Developing such tools will be the goal of the course.

# 1 Discrete Fourier Analysis and Roth's Theorem

**Notation.** Let  $N \in \mathbb{N}$  and  $\omega = e^{2\pi i/N}$ .

Write  $\mathbb{Z}_N$  for the cyclic group of integers mod  $N$ .

Write  $\mathbb{E}_x f(x) = N^{-1} \sum_{x \in \mathbb{Z}_N} f(x)$  (same notation also used for other averages).

**Definition** (Discrete Fourier Transform). If  $f : \mathbb{Z}_N \rightarrow \mathbb{C}$ , we define this as

$$\hat{f}(r) = \mathbb{E}_x f(x) \omega^{-rx}$$

Roughly speaking, expectations are taken in "physical space" and sums in "frequency space". This gives rise to two inner products:

$$\begin{aligned} \langle f, g \rangle &= \mathbb{E}_x f(x) \overline{g(x)} \\ \langle \hat{f}, \hat{g} \rangle &= \sum_r \hat{f}(r) \overline{\hat{g}(r)} \end{aligned}$$

**Fact** (Parseval's Identity).  $\langle f, g \rangle = \langle \hat{f}, \hat{g} \rangle$

**Definition** (Convolution). The convolution of  $f$  and  $g$  is:

$$\begin{aligned} f * g(x) &= \mathbb{E}_{y+z=x} f(y)g(z) \\ \hat{f} * \hat{g}(x) &= \sum_{s+t=x} \hat{f}(s) \hat{g}(t) \end{aligned}$$

**Fact** (Convolution Law). For  $f, g : \mathbb{Z}_N \rightarrow \mathbb{C}$ ,  $\widehat{f * g}(r) = \hat{f}(r) \hat{g}(r)$

**Fact** (Inversion Formula). For  $f : \mathbb{Z}_N \rightarrow \mathbb{C}$ , we get  $f(x) = \sum_r \hat{f}(r) \omega^{rx}$

**Definition** (Dilation Notation). If  $a$  is invertible mod  $N$ , then we write  $f_a(x) = f(a^{-1}x)$

**Fact** (Dilation Rule).  $\hat{f}_a(r) = \hat{f}(ar)$

**Definition** (Indicators). If  $A \subset \mathbb{Z}_N$ , write

$$A(x) = \begin{cases} 1, & x \in A \\ 0, & x \notin A \end{cases}$$

Note that if  $|A| = \alpha N$ , then  $\hat{A}(0) = \mathbb{E}_x A(x) = \alpha$

**Definition** (Norms). We define the norms  $\|f\|_p = (\mathbb{E}_x |f(x)|^p)^{1/p}$  and  $\|\hat{f}\|_p = \left( \sum_r |\hat{f}(r)|^p \right)^{1/p}$

So, if  $A \subset \mathbb{Z}_N$ , then  $\|A\|_2^2 = \langle A, A \rangle = \alpha$

Also, by Parseval, we get  $\sum_r |\hat{A}(r)|^2 = \alpha$

**Theorem** (Roth). For every  $\delta > 0$ ,  $\exists N$  s.t. every  $A \subset [N]$  of size at least  $\delta N$  contains an arithmetic progression of at least 3.

Lecture 2

*Basic Ideas.* We use a density increment argument. If  $A$  has density  $\alpha$  and has no 3 term A.P., then we can find a reasonably long A.P.  $P$  s.t.  $|A \cap P|/|P| \gg \alpha$ . Then, we're done or pass to  $P$  and start again with a better density. Iterate, and as  $\alpha \leq 1$ , we get an A.P..

Consider  $A \subset \mathbb{Z}_N$  and define  $B = C = A \cap [N/3, 2N/3]$ . Observe that  $(x, y, z)$  an A.P. in  $A \times B \times C$  in  $\mathbb{Z}_N$ , then it also is in  $[N]$ .

Let  $N$  be odd and  $\alpha$  the density of  $A$ .

If  $|B| < \frac{\alpha N}{5}$ , then either  $|A \cap [1, N/3]|$  or  $|A \cap (2N/3, N]|$  is at least  $\frac{2\alpha N}{5}$ , so we pass to one of these and increment  $A$ 's density.

Hence assume  $|B| = |C| \geq \frac{\alpha N}{5}$ .

We define the *3AP Density* of  $(A, B, C)$ , the probability that a random  $(x, y, z)$  with  $x + z = 2y$  lies in  $A \times B \times C$  as

$$\mathbb{E}_{x+z=2y} A(x)B(y)C(z)$$

Then we get, after some manipulation and calling  $\beta = \gamma = \|\hat{B}\|_2 = \|\hat{C}\|_2$  the density of  $B$  (and  $C$ ):

$$\mathbb{E}_{x+z=2y} A(x)B(y)C(z) = \alpha\beta\gamma + \sum_{r \neq 0} \hat{A}(r)\hat{B}(-2r)\hat{C}(r)$$

But

$$\left| \sum_{r \neq 0} \hat{A}(r)\hat{B}(-2r)\hat{C}(r) \right| \leq \max_{r \neq 0} |\hat{A}(r)| \|\hat{B}\|_2 \|\hat{C}\|_2$$

So, given  $\max_{r \neq 0} |\hat{A}(r)| \leq \frac{\alpha^2}{10}$  and  $\frac{\alpha^3}{50} > \frac{1}{N}$ , we get 3APs, so assume we have an  $r$  such that  $|\hat{A}(r)| \geq \frac{\alpha^2}{10}$

**Lemma.** Let  $\epsilon > 0, r \in \mathbb{Z}_N$ . Then  $[N]$  can be partitioned into APs of length at least  $\frac{\epsilon\sqrt{N}}{8\pi}$  on which the function  $x \mapsto \omega^{rx}$  varies by at most  $\epsilon$ .

*Proof.* Let  $m = \lfloor \sqrt{N} \rfloor$  and look at  $1, \omega^r, \dots, \omega^{mr}$ , two of which are sufficiently close, terms  $u$  and  $v$ . Write  $t = v - u$  and use this to get  $|\omega^{btr} - \omega^{atr}| \leq (b - a)\frac{2\pi}{m}$ , then partition  $[N]$  into congruence classes mod  $t$  and partition each congruence class.  $\square$

**Definition (Balanced Function).** **Balanced Function**  $f$  of  $A$  is  $f(x) = A(x) - \alpha$

Then  $\mathbb{E}_x f(x) = 0$  and for  $r \neq 0$ ,  $\hat{f}(r) = \hat{A}(r)$ . From before, we can pick  $r$  s.t.  $|\hat{f}(r)| \geq \alpha^2/10$ . If we let  $P_i$  be partitions given by the lemma, then we get:

$$\frac{\alpha^2}{10} \leq N^{-1} \left| \sum_x f(x)\omega^{-rx} \right| \leq \frac{\alpha^2}{20} + N^{-1} \sum_i \left| \sum_{x \in P_i} f(x) \right|$$

So,  $N^{-1} \sum_i \left| \sum_{x \in P_i} f(x) \right| \geq \frac{\alpha^2}{20} N$ .

Also,  $\sum_i \sum_{x \in P_i} f(x) = 0$ . Using these results, we can conclude that there is an  $i$  such that

$$\sum_{x \in P_i} f(x) \geq \frac{\alpha^2}{40} |P_i|$$

That is,  $|A \cap P_i| \geq (\alpha + \alpha^2/40) |P_i|$   $\square$

**Theorem** (Behrend, 1947). For every  $N, \exists$  a subset  $A \subset [N]$  of size  $N/\exp(c\sqrt{\log N})$

Lecture 3

*Proof.* Let  $m$  and  $d$  be positive integers, write in this proof  $[N] = \{0, 1, \dots, N-1\}$  and consider the grid  $[m]^d$ .

Note that in  $\mathbb{R}^d$ , no sphere contains points  $x, y, z$  with  $x + z = 2y$  (i.e. a 3AP). But  $\sum x_i^2$  can take at most  $m^2 d$  values in  $[m]^d$ , so there's a sphere that goes through at least  $\frac{m^d}{m^2 d}$  points.

Then, write a map  $\phi$  that sends points in  $[m]^d$  to a base  $2m$  number in the natural way, which gives  $\phi(x) + \phi(z) = 2\phi(y)$  iff  $x + z = 2y$ . Hence, we've constructed a set of at least  $m^{d-2}/d$  numbers, all less than  $(2m)^d$  with no 3AP. So, we let  $N = (2m)^d$  and, as  $d$  is arbitrary, we choose it to get the most points -  $d = \sqrt{\log N}$  is a pretty good choice, which gives us the bound.  $\square$

## 2 Bohr Sets and Bogolyubov's Method

**Definition** (Bohr Set). Let  $K \subset \mathbb{Z}_N$  and  $\epsilon > 0$ . Then, the **Bohr set** is defined as:

$$B(K, \epsilon) = \{x \in \mathbb{Z}_N : |1 - \omega^{rx}| \leq \epsilon \forall r \in K\}$$

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