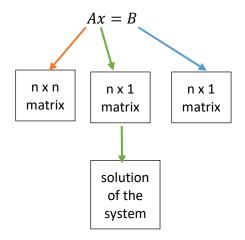
# **Chapter 3 Solution of Linear Systems**

In linear algebra, system of linear equations can be written as



N x N matrix 
$$\Rightarrow A = a_{ij} \Rightarrow \begin{bmatrix} a_{1\,1} & a_{1\,2} & a_{1\,3} & a_{1\,N} \\ a_{2\,1} & a_{2\,2} & a_{2\,3} & a_{2\,N} \\ ... & ... & ... & ... \\ a_{N\,1} & a_{N\,2} & a_{N\,3} & a_{N\,N} \end{bmatrix}$$

# 3.3 Upper Triangular Systems

An N x N matrix  $A = [a_{ij}]$  is called upper triangular matrix provided that  $a_{ij} = 0$  whenever

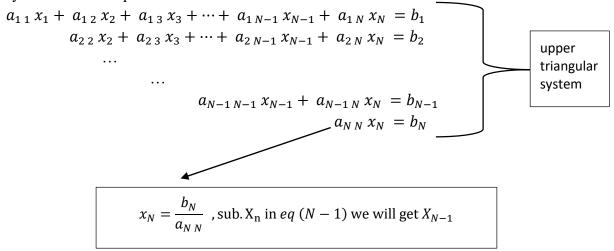
$$i > j \Rightarrow A = \begin{bmatrix} a_{1\,1} & a_{1\,2} & a_{1\,3} & a_{1\,N} \\ 0 & a_{2\,2} & a_{2\,3} & a_{2\,N} \\ 0 & 0 & a_{3\,3} & a_{3\,N} \\ 0 & 0 & 0 & a_{N\,N} \end{bmatrix}$$

An N x N matrix  $A = [a_{ij}]$  is called lower triangular matrix provided that  $a_{ij} = 0$  whenever

$$i < j \Rightarrow A = \begin{bmatrix} a_{1\,1} & 0 & 0 & 0 \\ a_{2\,1} & a_{2\,2} & 0 & 0 \\ a_{3\,1} & a_{3\,2} & a_{3\,3} & 0 \\ a_{N\,1} & a_{N\,2} & a_{N\,3} & a_{N\,N} \end{bmatrix}$$

The diagonal elements are  $a_{1\,1}, a_{2\,2}, \dots, a_{N\,N} \ (a_{k\,k}, k=1, \dots, N)$ 

If A is an upper triangular matrix, then the system Ax = B is said to be an upper triangular system of linear equations.



We can solve an upper triangular system by using back substitution.

Solve the following system of equations.

## Example 1:

$$\begin{array}{c}
 x_1 + 2x_2 + 3x_3 = 1 \\
 x_2 - 0x_3 = 0 \\
 x_3 = 5
 \end{array}
 \xrightarrow{\begin{array}{c}
 1 & 2 & 3 \\
 0 & 1 & 0 \\
 0 & 0 & 1
 \end{array}
 \begin{bmatrix}
 x_1 \\
 x_2 \\
 x_3
 \end{bmatrix}
 = \begin{bmatrix}
 1 \\
 0 \\
 5
 \end{bmatrix}$$

**Solution:** upper triangular system.

Last eq 
$$\Rightarrow x_3 = 5$$
  
 $2^{\text{nd}} \text{ eq} \Rightarrow x_2 - 0(5) = 0 \Rightarrow x_2 = 0$   
 $1^{\text{st}} \text{ eq} \Rightarrow x_1 - 2(0) + 3(5) = 1 \Rightarrow x_1 = -14$ 

The solution is  $(x_1, x_2, x_3) = (-14, 0, 5) \rightarrow \text{UNIQUE}$  solution (here  $a_{kk} \neq 0$  for all k = 1,2,3)

Example 2:

$$\begin{array}{c}
 x_1 + 0x_2 + 2x_3 &= 0 \\
 x_2 + 0x_3 &= 0 \\
 0x_3 &= 0
 \end{array}
 \xrightarrow{\begin{array}{c}
 1 & 0 & 2 \\
 0 & 1 & 0 \\
 0 & 0 & 0
 \end{array}
 \begin{bmatrix}
 x_1 \\
 x_2 \\
 x_3
 \end{bmatrix}
 = \begin{bmatrix}
 0 \\
 0 \\
 0
 \end{bmatrix}$$

$$\begin{array}{c}
 x_1 \\
 x_2 \\
 x_3
 \end{bmatrix}
 = \begin{bmatrix}
 0 \\
 0 \\
 0
 \end{bmatrix}$$

$$\begin{array}{c}
 x_1 \\
 x_2 \\
 x_3
 \end{bmatrix}
 = \begin{bmatrix}
 0 \\
 0 \\
 0
 \end{bmatrix}$$

#### **Solution:**

Last eq. is satisfied for any real number  $x_3$ , say  $x_3 = t$ 

$$2^{\rm nd} \, {\rm eq} \Rightarrow x_2 = 0$$

$$1^{st} eq \Rightarrow x_1 - 2t = 0 \Rightarrow x_1 = -2t$$

The solution is  $(x_1, x_2, x_3) = (-2t, 0, t)$ , where  $t \in \mathbb{R} \Rightarrow \text{INFINITE}$  solutions.  $(a_{kk} \neq 0 \text{ for all } k = 3)$ 

# Example 3:

$$\begin{array}{c} x_1 - x_2 + x_3 + x_4 = 1 \\ 0x_2 + x_3 + x_4 = 3 \\ x_3 - x_4 = 0 \\ x_4 = 1 \end{array} \rightarrow \begin{bmatrix} 1 & -1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \\ 0 \\ 1 \end{bmatrix}$$

# **Solution:**

Last eq. 
$$\Rightarrow x_4 = 1$$
  
 $3^{\text{rd}}$  eq.  $\Rightarrow x_3 - 1 = 0 \Rightarrow x_3 = 1$   
 $2^{\text{nd}}$  eq  $\Rightarrow 0x_2 + 1 + 1 = 3$ 

$$0 + 2 = 3$$
 not true for any value of  $x_2$ 

The system has NO solutions.  $(a_{22} = 0)$ 

0 + 2 = 3 not true for any value of  $x_2$ , so the system has no solution  $(a_{2\,2}=0)$ 

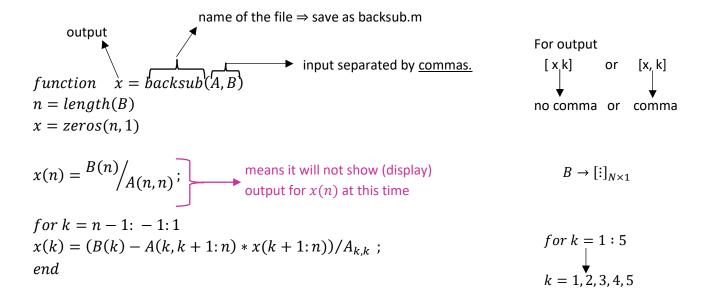
The upper triangular system has <u>a unique solution</u> if  $a_{kk} \neq 0$  for all k = 1, 2, ..., N

In this course, we will focus on the solutions of system for which we can find a unique solution.

Textbook: Appendix: Introduction to MATLAB. Book-marked.

Most of the programs are in the textbook.

Upper triangular system, the book has a program for back substitution on page 123.



$$a_{1\,1}x_1 + a_{1\,2}x_2 + a_{1\,3}x_3 + \dots + a_{1\,N-1}x_{N-1} + a_{1\,N}x_N = b_1 - (1)$$

$$a_{2\,2}x_2 + a_{2\,3}x_3 + \dots + a_{2\,N-1}x_{N-1} + a_{2\,N}x_N = b_2 - (2)$$

$$\vdots \qquad \vdots \qquad \vdots$$

$$a_{N-1\,N-1}x_{N-1} + a_{N-1\,N}x_N = b_{N-1} - (N-1)$$

$$a_{N\,N}x_N = b_N - (N)$$

$$x_N = \frac{b_N}{a_{N\,N}} \quad , \quad \text{sub. } x_N \text{ in } eq\ (N-1) \text{ we will get } x_{N-1}$$

In back substitution, we do N dimensions (1 at each step)

multiplications: 
$$0+1+2+\cdots+(N-1)=\frac{(N-1)N}{2}=\frac{N^2-N}{2}$$
 
$$\sum_{i=1}^{N}i=\frac{n(n+1)}{2}$$
 subtractions / additions:  $0+1+2+\cdots+(N-1)=\frac{(N-1)N}{2}=\frac{N^2-N}{2}$  Total operations are  $N+\frac{N^2-N}{2}+\frac{N^2-N}{2}=\frac{2N+N^2-N+N^2-N}{2}=\frac{2N^2}{2}=N^2$  Computational complexity of back substitution is of  $O(N^2)$ .