

COMP 2310 Final Exam

1. Prove $P1, P2, P3 \vdash C$

$$P1: (A \wedge B) \Rightarrow C \quad \sim A \vee \sim B \vee C$$

$$P2: (A \Rightarrow C) \Rightarrow D \quad B \Rightarrow E$$

$$P3: \sim B \vee E \quad E$$

$$C: B \Rightarrow (D \wedge E)$$

(Direct proof)

1. B Premise

2. $(A \wedge B) \Rightarrow C$ from 1

3. $(A \Rightarrow C) \Rightarrow D$ from 1

4. $\sim B \vee E$ from 1

5. $B \Rightarrow E$ 4, EI8

6. E 1, 5, I3

7. $\sim(A \wedge B) \vee C$ 2, EI8

8. $(\sim A \vee \sim B) \vee C$ 7, EI6

9. $\sim A \vee (\sim B \vee C)$ 8, EI2

10. $\sim A \vee (C \vee \sim B)$ 9, EI10

11. $(\sim A \vee C) \vee \sim B$ 10, EI2

12. $(A \Rightarrow C) \vee \sim B$ 11, EI8

13. $\sim B \vee (A \Rightarrow C)$ 12, EI10

14. $B \Rightarrow (A \Rightarrow C)$ 13, EI8

15. $B \Rightarrow D$ 14, I5

16. D 1, 15, I3

17. $D \wedge E$ 16, 6, I6

Thus $P1, P2, P3 \vdash C$.

□

2. Prove $P_1, P_2, P_3 \vdash C$

$$P_1: (\forall x)(\forall y)((P(x) \wedge Q(y)) \Rightarrow R(x,y))$$

$$P_2: (\exists x)(\forall y)((P(x) \wedge S(x,y)) \Rightarrow Q(y))$$

$$P_3: (\forall x)(\exists y)(P(x) \wedge \sim R(x,y) \wedge T(x,y))$$

$$C: (\exists x)(\exists y)(P(x) \wedge \sim S(x,y) \wedge T(x,y))$$

(proof by definition)

1. $(\forall x)(\forall y)((P(x) \wedge Q(y)) \Rightarrow R(x,y))$ from 1
2. $(\exists x)(\forall y)((P(x) \wedge S(x,y)) \Rightarrow Q(y))$ from 1
3. $(\forall y)((P(a) \wedge S(a,y)) \Rightarrow Q(y))$ 2, EI, a is a constant
4. $(\forall x)(\exists y)(P(x) \wedge \sim R(x,y) \wedge T(x,y))$ from 1
5. $(\exists y)(P(a) \wedge \sim R(a,y) \wedge T(a,y))$ 4, UI
6. $(P(a) \wedge \sim R(a,b)) \wedge T(a,b)$ 5, EI, b is a constant
7. $P(a) \wedge \sim R(a,b)$ 6, I2
8. $T(a,b) \wedge (P(a) \wedge \sim R(a,b))$ 6, I2
9. $T(a,b)$ 8, I2
10. $P(a)$ 7, I2
11. $\sim R(a,b) \wedge P(a)$ 7, EI
12. $\sim R(a,b)$ 11, I2
13. $(\forall y)((P(a) \wedge Q(y)) \Rightarrow R(a,y))$ 1, UI
14. $(P(a) \wedge Q(b)) \Rightarrow R(a,b)$ 13, UI
15. $\sim(P(a) \wedge Q(b))$ 12, 14, I4
16. $\sim P(a) \vee \sim Q(b)$ 15, E16
17. $\sim Q(b) \vee \sim P(a)$ 16, E10
18. $Q(b) \Rightarrow \sim P(a)$ 17, E18
19. $(P(a) \wedge S(a,b)) \Rightarrow Q(b)$ 3, UI
20. $(P(a) \wedge S(a,b)) \Rightarrow \sim P(a)$ 19, 18, I5
21. $\sim \sim P(a)$ 10, E15
22. $\sim(P(a) \wedge S(a,b))$ 21, 20, I4
23. $\sim P(a) \vee \sim S(a,b)$ 22, E16
24. $P(a) \wedge (\sim P(a) \vee \sim S(a,b))$ 10, 23, I6
25. $(P(a) \wedge \sim P(a)) \vee (P(a) \wedge \sim S(a,b))$ 24, E13
26. false $\vee (P(a) \wedge \sim S(a,b))$ 25, E1
27. $(P(a) \wedge \sim S(a,b)) \vee \text{false}$ 26, E10
28. $P(a) \wedge \sim S(a,b)$ 27, E6
29. $(P(a) \wedge \sim S(a,b)) \wedge T(a,b)$ 28, 9, I6
30. $(\exists y)(P(a) \wedge \sim S(a,y) \wedge T(a,y))$ 29, EQ
31. $(\exists x)(\exists y)(P(x) \wedge \sim S(x,y) \wedge T(x,y))$ 30, EQ

Hence $P_1, P_2, P_3 \vdash C$

□

3. Prove that $B - A = B \iff B \cap A = \emptyset$
(biconditional proof)

$$B - A = B \iff B \cap \bar{A} = B \quad (\text{theorem})$$

$$\iff (\forall x)(x \in B \cap \bar{A} \iff x \in B) \quad (\text{principle of extension})$$

$$\iff (\forall x)((x \in B \cap x \in \bar{A}) \iff x \in B) \quad (\text{definition of } \cap)$$

$$\iff (\forall x)((x \in B \cap x \notin A) \iff x \in B) \quad (\text{theorem})$$

$$\iff (\forall x)((x \in B \cap x \notin A \Rightarrow x \in B) \wedge (x \in B \Rightarrow (x \in B \cap x \notin A))) \quad (E20)$$

$$\iff (\forall x)((\sim(x \in B \cap x \notin A) \vee x \in B) \wedge (\sim x \in B \vee (x \in B \cap x \notin A))) \quad (E18)$$

$$\iff (\forall x)((x \notin B \vee x \in A) \vee x \in B) \wedge (\sim x \in B \vee (x \in B \cap x \notin A)) \quad (E16, E15)$$

$$\iff (\forall x)((x \notin B \vee (x \in B \vee x \in A)) \wedge ((\sim x \in B \vee x \in B) \wedge (\sim x \in B \vee x \notin A))) \quad (E17, E10, E14)$$

$$\iff (\forall x)((x \notin B \vee x \in B) \vee x \in A) \wedge ((x \in B \vee \sim x \in B) \wedge (\sim x \in B \vee x \in A)) \quad (E12, E10)$$

$$\iff (\forall x)((x \in B \vee \sim x \in B) \vee x \in A) \wedge (\text{true} \wedge (\sim x \in B \vee x \in A)) \quad (E10, E2)$$

$$\iff (\forall x)((\text{true} \vee x \in A) \wedge ((\sim x \in B \vee x \in A) \wedge \text{true})) \quad (E2, E9)$$

$$\iff (\forall x)((x \in A \vee \text{true}) \wedge (\sim x \in B \vee x \in A)) \quad (E10, E5)$$

$$\iff (\forall x)(\text{true} \wedge (\sim x \in B \vee x \in A)) \quad (E8)$$

$$\iff (\forall x)(\sim(x \in B \wedge x \in A) \wedge \text{true}) \quad (E16, E9)$$

$$\iff (\forall x)\sim(x \in B \wedge x \in A) \quad (E5)$$

$$\iff (\forall x)(\sim(x \in B \wedge x \in A) \vee \text{false}) \quad (E4)$$

$$\iff (\forall x)(\sim(x \in B \wedge x \in A) \vee x \in \emptyset) \quad (\text{theorem})$$

$$\iff (\forall x)(\sim(x \in B \cap A) \vee x \in \emptyset) \quad (\text{definition of } \cap)$$

$$\iff (\forall x)(x \in B \cap A \Rightarrow x \in \emptyset) \quad (E18)$$

$$\iff B \cap A = \emptyset \quad (\text{principle of extension})$$

Thus $B - A = B \iff B \cap A = \emptyset$.

□

4. let N be the set of all positive integers.

let R be a relation in $N \times N$ such that $((a,b), (x,y)) \in R \Leftrightarrow \frac{2a+1}{2^b} \leq \frac{2x+1}{2^y}$

Prove R is a partial order.

A partial order is reflexive, antisymmetric and transitive. Thus we break the proof into 3 steps:

a) R is reflexive.

We shall show that $(\forall (a,b) \in N \times N) ((a,b), (a,b)) \in R$ (definition of reflexive)
 $\Rightarrow (a,b) \in N \times N \Rightarrow ((a,b), (a,b)) \in R$ (UI)

(Direct proof)

Suppose $(a,b) \in N \times N$ (hypothesis)

$\Rightarrow a \in N \wedge b \in N$ (definition of \times)

We have $\frac{2a+1}{2^b} = \frac{2a+1}{2^b}$ (= is reflexive)

$\Rightarrow \frac{2a+1}{2^b} \leq \frac{2a+1}{2^b}$ (definition of \leq)

$\Rightarrow ((a,b), (a,b)) \in R$ (definition of R)

Thus R is reflexive (definition of reflexive)

b) R is antisymmetric

We shall prove $((a,b), (x,y)) \in R \wedge ((x,y), (a,b)) \in R \Rightarrow (a,b) = (x,y)$

(direct proof)

Suppose $((a,b), (x,y)) \in R \wedge ((x,y), (a,b)) \in R$ (hypothesis)

$\Rightarrow \begin{cases} ((a,b), (x,y)) \in R \\ ((x,y), (a,b)) \in R \end{cases}$ (E9, I2)

$\Rightarrow \begin{cases} \frac{2a+1}{2^b} \leq \frac{2x+1}{2^y} \dots (1) \\ \frac{2x+1}{2^y} \leq \frac{2a+1}{2^b} \dots (2) \end{cases}$ (definition of R)

We then have $\frac{2a+1}{2^b} \leq \frac{2x+1}{2^y} \wedge \frac{2x+1}{2^y} \leq \frac{2a+1}{2^b}$ ((1), (2), I6)

$\Rightarrow \left(\frac{2a+1}{2^b} < \frac{2x+1}{2^y} \vee \frac{2a+1}{2^b} = \frac{2x+1}{2^y} \right) \wedge \left(\frac{2x+1}{2^y} < \frac{2a+1}{2^b} \vee \frac{2x+1}{2^y} = \frac{2a+1}{2^b} \right)$ (definition of \leq)

$\Rightarrow \left(\frac{2a+1}{2^b} = \frac{2x+1}{2^y} \vee \frac{2a+1}{2^b} < \frac{2x+1}{2^y} \right) \wedge \left(\frac{2x+1}{2^y} = \frac{2a+1}{2^b} \vee \frac{2x+1}{2^y} < \frac{2a+1}{2^b} \right)$ (E10, = is reflexive)

$\Rightarrow \frac{2a+1}{2^b} = \frac{2x+1}{2^y} \vee \left(\frac{2a+1}{2^b} < \frac{2x+1}{2^y} \wedge \frac{2x+1}{2^y} < \frac{2a+1}{2^b} \right)$ (E14)

$\Rightarrow \frac{2a+1}{2^b} = \frac{2x+1}{2^y} \vee \text{false}$ ($\because \frac{2a+1}{2^b}$ can't be $<$ and $> \frac{2x+1}{2^y}$)

$\Rightarrow a=x \wedge b=y$ (hsa)

$\Rightarrow (a,b) = (x,y)$ (theorem)

Thus R is antisymmetric.

c) R is transitive

We shall prove $((a,b), (c,d)) \in R \wedge ((c,d), (e,f)) \in R \Rightarrow ((a,b), (e,f)) \in R$ (definition of transitive)
(direct proof)

Suppose $((a,b), (c,d)) \in R \wedge ((c,d), (e,f)) \in R$ (hypothesis)

$$\Rightarrow \begin{cases} ((a,b), (c,d)) \in R \\ ((c,d), (e,f)) \in R \end{cases} \quad (E9, I2)$$

$$\Rightarrow \begin{cases} \frac{2a+1}{2^b} \leq \frac{2c+1}{2^d} & \text{(definition of } R) \dots (3) \\ \frac{2c+1}{2^d} \leq \frac{2e+1}{2^f} & \dots (4) \end{cases}$$

We then have $\frac{2a+1}{2^b} \leq \frac{2c+1}{2^d} \wedge \frac{2c+1}{2^d} \leq \frac{2e+1}{2^f}$ $((3), (4), I6)$

$$\Rightarrow \frac{2a+1}{2^b} \leq \frac{2e+1}{2^f} \quad (\leq \text{ is transitive } \because \leq \text{ is a partial order})$$

$$\Rightarrow ((a,b), (e,f)) \in R \quad (\text{definition of } R)$$

Thus R is transitive.

We have shown that R is reflexive, antisymmetric and transitive.
Thus, R is a partial order.

□

5. Let $f: A \rightarrow B$.

Prove $(\exists g)(g: B \rightarrow A \wedge g \circ f = I_A) \Rightarrow f$ is onto

$\Rightarrow (\exists g)(g: B \rightarrow A \wedge g \circ f = I_A) \Rightarrow (\forall y \in B)(\exists x \in A) y = f(x)$ (definition of onto)

(direct proof)

Suppose $(\exists g)(g: B \rightarrow A \wedge g \circ f = I_A)$ (hypothesis)

$\Rightarrow h: B \rightarrow A \wedge h \circ f = I_A$ (EI)

$\Rightarrow h \circ f = I_A$ (E9, I2) ... (1)

We have $x = I_A(x)$ (definition of I_A)

$\Rightarrow x = h \circ f(x)$ ((1), sub=)

$\Rightarrow (\exists z \in B)(z = h(x) \wedge x = f(z))$ (definition of \circ)

$\Rightarrow a \in B \wedge (a = h(x) \wedge x = f(a))$ (EI)

$\Rightarrow a \in B \wedge (x = f(a) \wedge a = h(x))$ (E9)

$\Rightarrow (a \in B \wedge x = f(a)) \wedge a = h(x)$ (E11)

$\Rightarrow a \in B \wedge x = f(a)$ (I2)

$\Rightarrow (\exists a \in B) x = f(a)$ (E2)

Hence, f is onto.

□

6. Prove there does not exist a simple graph of order 4 and size 7.

Equivalently, we shall prove that there does not exist a simple graph where $|V|=4$ and $|E|=7$.

(proof by contradiction)

Suppose there exists a graph G where $|V|=4$ and $|E|=7$.

We know the maximum degree sequence of G is a complete graph, since a complete graph has all vertices paired with an edge or more formally, $(\forall v_1, v_2 \in V) \{v_1, v_2\} \in E \wedge v_1 \neq v_2$

The maximum degree sequence would be $\{3, 3, 3, 3\}$.

This is because every vertex has 3 other vertices to pair with, in an edge.

A vertex with a degree >3 implies $|V|>4$ since G is a simple graph.

However, $|V|=4$ in graph G .

Thus the largest degree sequence is $\{3, 3, 3, 3\}$.

If we add up all the degrees in the sequence, we get

$$\sum_{v \in V} \deg(v) = 3+3+3+3 = 12.$$

However, $\sum_{v \in V} \deg(v) = 2|E|$ (theorem)

$$\Rightarrow \sum_{v \in V} \deg(v) = 2(7) \quad (|E|=7)$$

$$\Rightarrow \sum_{v \in V} \deg(v) = 14.$$

We know $12 < 14$ (h.s.a)

This implies there are some self-loops or parallel edges in graph G .

However, G is simple, so there are no self-loops or parallel edges in G .

Hence, we have a contradiction!

Thus, there does not exist a simple graph G of order 4 and size 7.

□